

No calculators or notes of any kind are allowed.

This exam consists of 5 problems, each of equal weight. All answers must be justified with appropriate reasoning and calculations.

Notation:

$$\begin{split} \mathbb{N} &= \{0,1,2,3,\ldots\}\\ \mathbb{R}_{\geq 0} &= \{x \in \mathbb{R} \,|\, x \geq 0\}\\ a \text{ divides } b \text{ is denoted } a | b \end{split}$$

Question 1. Let P and Q be propositions

(a) Give one example of a sentence which is a proposition. Give one example of a sentence which is not a proposition. These can be mathematical or any other type of sentences.

Solution: Some examples of propositions: Two plus two equals four. Two plus two equals five. Right now the sky is blue. $\forall x, 3 + x = 7$. Some examples which are not propositions: Does two plus two equal 4? I always tell lies. 3 + x = 7 (does not have truth value unless x is specified or there is a quantifier).

(b) If P is false and Q is true, what is the truth value of P ⇒ Q. This often confuses people when they first learn about "⇒". Explain why this is the right way to define "⇒" with the help of an example (in mathematical or everyday language) of propositions P and Q.

Solution 1: P="2 + 2 = 5", Q = "1 + 1 = 2". We can obtain Q from P by multiplying both sides of P by zero and then adding 1 to both sides twice.

Solution 2: P="It is Monday", Q ="I will wear a hat". Then $P \implies Q$ is the same as the statement "If it is Monday then I will wear a hat". Now if today is Tuesday and I am wearing a hat am I a truth teller or a liar. In my if ..then.. statement I did not claim anything about Tuesday. So no matter if I wear a hat or not on Tuesday I was telling the truth.

(c) Determine if $(P \Longrightarrow Q) \lor (Q \Longrightarrow P)$ is a tautology or not by using a truth table. Be sure to state your conclusion.

Solution:

$T ext{ } F ext{ } F ext{ } T ext{ } T$
F T T F T
F F T T T

The proposition $(P \implies Q) \lor (Q \implies P)$ is thus a tautology as it is true in all cases, regardless of the truth value of the individual propositions P and Q.

Question 2. For each statement below, indicate if it is TRUE(T) or FALSE(F). No justification is needed. Write your answers in your exam booklet in a table like this

	a	b	с	d	e	f	g	h	i	j
T/F	F	Т	F	F	Т	Т	F	F	F	Т

(a) $f : \{1, 2, 3\} \to \{1, 3\}$ defined by f(x) = x is a function.

F: f is not defined at all elements in the domain. In particular f(2) is not defined

(b) $f : \mathbb{R}_{\geq 0} \to \mathbb{R}_{\geq 0}$ defined by $f(x) = x^2$ is a bijection.

T: The function is a bijection because $y = x^2$ has a unique solution for all $y \in \mathbb{R}_{\geq 0}$.



Also we can see that x^2 has a well-defined inverse \sqrt{x} for $x \ge 0$.

(c) $f : \mathbb{N} \to \mathbb{N}$ defined by $f(x) = x^2$ is a bijection.

F: Not surjective. For example there is no element that maps to 3. That is $f(x) = x^2 = 3$ does not have a natural number solution.

(d) The cardinalities of \mathbb{N} and \mathbb{Z} are different because \mathbb{Z} is essentially double \mathbb{N} .

F: As discussed in class these have the same cardinality. For example here is a bijection

$$f(n) = \begin{cases} -\frac{n}{2} & \text{if } n \text{ is even} \\ \frac{n+1}{2} & \text{if } n \text{ is odd} \end{cases}$$

(e) The cardinalities of \mathbb{Q} and \mathbb{R} are different even though every real number can be approximated to arbitrary accuracy by rational numbers.

T: We discussed Cantor's diagonal argument in class that proves \mathbb{Q} and \mathbb{R} have different cardinalities. The ideas of the proof was to show that there can be no surjection.

(f) Let A be a finite set with cardinality |A| = n.

 $|\mathcal{P}(A \times A)| =$ "the cardinality of the power set of the cartesian product $A \times A$ " = $(2^n)^n$?

T: $|\mathcal{P}(A \times A)| = 2^{|A \times A|} = 2^{(n^2)} = 2^{(n \cdot n)} = (2^n)^n$

- (g) Let P(x) and Q(y) be propositions depending on variables x and y respectively. The negation of ∀x∃y(P(x) ⇒ Q(y)) is ∃y∀x(P(x) ∧ ¬Q). F: ∀xR(x) negates to ∃x¬R(x).
- (h) Define a relation ~ on ℝ by x ~ y if and only if x² ≤ y². The relation ~ defines a partial order on ℝ.
 F: If x = -1 and y = 1 then x² ≤ y² and y² ≤ x², but x ≠ y. So the relation is not antisymmetric.
- (i) If A and B are sets such that A × B = B × A then A = B
 F: This was from the homework. If A is any set and B = Ø then A × B = B × A = Ø.
- (j) Let A and B be finite sets. If |B| > |A| then there exists an injection f : A → B.
 T: See the homework. This is difficult to prove in general, but for finite sets it is easier.
- Question 3. (a) State the inclusion-exclusion principle for 3 sets. Solution: $|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|$.
 - (b) How many 7 digit numbers in N start or end with the digits 21. For example 2156973, 2982721. Solution: There are 10⁵ numbers that start with 21 and 9×10⁴ numbers that end in 21 because numbers can't start with a zero (only one point was taken off in case this was missed). There are 10³ numbers both starting and ending with 21. Using inclusion-exclusion for 2 sets, the answer is 100000 + 90000 - 1000 = 189000
 - (c) Use induction to prove that for all $n \in \mathbb{N}$, $n \ge 1$,

$$\sum_{k=1}^{n} k^2 = \frac{n(n+1)(2n+1)}{6}$$

Solution: To prove the formula using induction, we'll start by verifying the base case when n = 1:

When n = 1:

$$\sum_{k=1}^{1} k^2 = 1^2 = 1$$

And

$$\frac{1(1+1)(2(1)+1)}{6} = \frac{1(2)(3)}{6} = 1$$

The formula holds for n = 1.

Next, we'll assume that the formula holds for some arbitrary positive integer k, and then prove that it also holds for k + 1.

Assuming the formula holds for k:

$$\sum_{i=1}^{k} i^2 = \frac{k(k+1)(2k+1)}{6}$$

Now, let's consider the sum for k + 1:

$$\sum_{i=1}^{k+1} i^2 = (k+1)^2 + \sum_{i=1}^{k} i^2$$

Now use the induction assumption.

$$\begin{split} \sum_{i=1}^{k+1} i^2 &= (k+1)^2 + \frac{k(k+1)(2k+1)}{6} \\ &= \frac{6(k+1)^2 + k(k+1)(2k+1)}{6} \\ &= \frac{(k+1)(6(k+1) + k(2k+1))}{6} \\ &= \frac{(k+1)(2k^2 + 7k + 6)}{6} \\ &= \frac{(k+1)(k+2)(2k+3)}{6} \\ &= \frac{(k+1)((k+1) + 1)(2(k+1) + 1)}{6} \end{split}$$

We have the expression in the same form as the original formula, but with k + 1 substituted for n. Thus, by induction, we have shown that the formula holds for all $n \ge 1$. Note: It's probably easier to expand the above expression and compare with the expansion of what we are trying to show with n = k + 1, but factoring is more fun!

Therefore, the equation

$$\sum_{k=1}^{n} k^2 = \frac{n(n+1)(2n+1)}{6}$$

holds for all $n \ge 1$.

(d) Consider the permutation $\rho = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 5 & 6 & 2 & 1 & 8 & 3 & 4 & 7 \end{pmatrix}$. Write ρ^2 as a product of disjoint cycles.

solution: $\rho^2 = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 8 & 3 & 6 & 5 & 7 & 2 & 1 & 4 \end{pmatrix} = (18457)(236)$

Question 4. Consider the weighted graph $G = (V, \mathcal{E}, w)$ where the vertex set is $V = \{A, B, C, D, E\}$, the edge set is $\mathcal{E} = \{\{A, C\}, \{A, D\}, \{A, E\}, \{B, E\}, \{C, E\}, \{D, E\}\}$ and the weights are $w(\{A, C\}) = 4, w(\{A, D\}) = 10, w(\{A, E\}) = 8, w(\{B, E\}) = 2, w(\{C, E\}) = 9, w(\{D, E\}) = 6.$

(a) Make a large clear sketch of the graph. It is suggested that you draw it as a regular pentagon.



(b) Find a minimal spanning tree of G using a greedy algorithm. Your answer can be a sketch. Briefly explain each step in your selection of edges.

Solution:

To find a minimum spanning tree (MST) using Prim's algorithm, we: (1) Initialize a tree with a single vertex, chosen arbitrarily from the graph. (2) Grow the tree by one edge: Of the edges that connect the tree to vertices not yet in the tree, find the minimum-weight edge, and transfer it to the tree. (3) Repeat step 2 until all vertices are in the tree.

Let's go step by step:

Step 1: Start with vertex A as the initial tree.

Step 2: Look for the minimum weight edge that connects the current tree to a vertex outside the tree. Add that vertex and the edge to the MST. Minimum weight edge is A,C. MST: {AC} (just writing the edges and not the vertices)

Step 3: Minimum weight edge is A,E. MST: {AC, AE}

Step 4: . Minimum weight edge is B,E. MST: {AC, AE, BE}

Step 5: Minimum weight edge is D,E. MST: {AC, AE, BE, DE}

Step 6: All vertices have been included in the MST, so the algorithm is complete.

The final MST starting at vertex A is: MST: {AC, AE, BE, DE}



Or, redrawing to look more like a tree.



(c) Use a greedy algorithm to find a (vertex) colouring of G. Your answer can be a sketch.



(d) What is the chromatic number of G. Justify your answer.

We can see that vertices A, C, and E form a complete subgraph K3 (a triangle), as they are pairwise adjacent. Thus, we have an instance of K3 within graph G.

Therefore, since graph G contains a subgraph isomorphic to K3, we can conclude that the chromatic number of graph G is at least 3 (lower bound). Additionally, since we were able to color graph G using three colors, we have shown that the chromatic number is at most 3 (upper bound).

Hence, we can state that the chromatic number of graph G is exactly 3.

Question 5. (a) Find all solutions to the linear Diophantine equation: 60x + 33y = 9.

Method A: Here is the standard but long solution.

To find all solutions to the Diophantine equation 60x + 33y = 9, we will use the extended Euclidean algorithm. The first step is to find the greatest common divisor (GCD) of 60 and 33.

Step 1: Find the GCD of 60 and 33.

Using the Euclidean algorithm:

 $60 = 1 \times 33 + 27$ $33 = 1 \times 27 + 6$ $27 = 4 \times 6 + 3$ $6 = 2 \times 3 + 0$

So the GCD of 60 and 33 is 3. Since 3|9 we know there is a solution to the Diophantine equation.

Step 2: Express the GCD (3) as a linear combination of 60 and 33.

Using the Euclidean algorithm backward:

$$3 = 27 - 4 \times 6$$

= 27 - 4 × (33 - 27)
= 5 × 27 - 4 × 33
= 5 × (60 - 33) - 4 × 33
= 5 × 60 - 9 × 33

Therefore $15 \times 60 - 27 \times 33 = 9$.

Thus, we have found one solution to the equation 60x + 33y = 3, which is $x_0 = 15$ and $y_0 = -27$.

Step 3: Find all solutions by adding all solution of 60x + 33y = 0 which, by reducing to the equivalent equation 20x + 11y = 0 by dividing by the GCD, is seen to be x = 11n, y = -20n.

All solution of the Diophantine equation are

$$\begin{aligned} x &= 15 + 11n\\ y &= -27 - 20n \end{aligned}$$

for $n \in \mathbb{Z}$.

Method B:

Dividing the equation by 3 we get 20x + 11y = 3 which cannot be reduced further since 11 is a prime. To get the units digit to be 3, we see y has to end in a 3 or 7. We see that $x_0 = 4_{,0} = -7$ is a solution. The solutions to 20x + 11y = 0, is x = 11n, y = -20n. So the general solution to the Diophantine equation is x = 15 + 11n, y = -27 - 20n for $n \in \mathbb{Z}$.

- (b) Depending how you solve this problem it may be useful to recall that
 - a|b is an order relation.
 - Let p be a prime and $a_i \in \mathbb{Z}$. If $p|a_1a_2\cdots a_n$ then $p|a_1 \vee p|a_2 \vee \cdots \vee p|a_n$. A special case is that if p|ab then p divides a or p divides b.

Let p be a prime and $n \in \mathbb{N}$. Prove that the set of divisors of p^n is $\{1, p, p^2, \dots, p^{n-1}, p^n\}$. Solution:

Let D be the let of divisors of p^n . We need to show (A) $D \subseteq \{1, p, p^2, \dots, p^{n-1}, p^n\}$ and, (B) $D \supseteq \{1, p, p^2, \dots, p^{n-1}, p^n\}$.

(A) For each i = 1, ..., n we can write $p^n = p^i p^{n-i}$. Since p^i divides the right-hand side, it also divides the left. So $p^i \in D$ for i = 1, ..., n, and thus we have shown (A).

(B) Let $q \in D$. By the fundamental theorem of arithmetic, q can be written as a product of primes. That is $q = p_1 \cdot p_2 \cdots p_m$. Now for each $k = 1 \dots m$, $p_k | q$. Since $q | p^n$, by transitivity of divisibility order relation, we have $p_k | p^n$ and by applying the given theorem we have $p_k | p$. So $p_k = 1$ or p. Thus $q = p^r$, and $r \leq n$ as $q \leq p^n$. So we have shown (B) which completes the proof.

(B) can also be proved by invoking the uniqueness of prime factorization: As the factorization of p^n is $p^n = p \cdot p \cdots p$, there can be no other prime factor of p^n except p. Then use the argument in (B) above for any q that divides p^n .

Many people tried as inductive proof, which is possible but it requires the same argument as in step (B), so is of no advantage.