## CS-E4500 Advanced Course in Algorithms (5 cr)

1. Randomized polynomial identity testing. Let $F$ be a field with at least $q$ elements.
(a) Let $f, \tilde{f} \in F[x]$ be polynomials of degree at most $d$. Show that if $f \neq \tilde{f}$ then a uniform random $\xi \in F$ satisfies $f(\xi) \neq \tilde{f}(\xi)$ with probability at least $1-d / q$.
(b) Let $a, b, c \in F[x]$ be three polynomials, each of degree at most $d$ and each given as a sequence of coefficients. Present a randomized test that verifies $c=a b$ and uses $O(d)$ operations in $F$. If $c=a b$ the test must accept with probability 1 ; if $c \neq a b$ the test must reject with probability at least $1-d / q$.

Hints: For part (a), recall what we know about low-degree polynomials. For part (b), reduce to part (a) and carefully justify that your algorithm uses $O(d)$ operations in $F$.
2. Testing a matrix product. Let $A, B, C$ be three $n \times n$ matrices with entries in a field $F$. Present a randomized algorithm that tests whether $C=A B$ using $O\left(n^{2}\right)$ operations in $F$. When $C=A B$, your algorithm must always assert that $C=A B$. When $C \neq A B$, your algorithm must assert that $C \neq A B$ with probability at least $1 / 2$.
Hints: Select a uniform random $x \in\{0,1\}^{n} \subseteq F^{n}$. Study the probability that $C x \neq$ $A(B x)$ when $C \neq A B$.
3. Evaluation algorithm for the \#CNFSAT proof polynomial $P_{\mathscr{C}}$. Let $\mathscr{C}$ be a collection of clauses $C_{1}, C_{2}, \ldots, C_{m}$ over $n$ variables $x_{1}, x_{2}, \ldots, x_{n}$ taking values in $\{0,1\}$. Present detailed pseudocode for an algorithm that, given as input $\mathscr{C}$, a prime $q$ with $2^{n / 2+2} m n \leq q \leq 2^{n / 2+3} m n$, and $\xi \in \mathbb{F}_{q}$, computes the value $P_{\mathscr{C}}(\xi) \in \mathbb{F}_{q}$ in time $O\left(2^{n / 2}(m n)^{c}\right)$ for some constant $c>0$. Carefully justify the running time of your algorithm. You may use the near-linear-time toolbox for univariate polynomials and algorithms for modular arithmetic in $\mathbb{F}_{q}$ as subroutines without detailed pseudocode, but make sure that you specify with care the input to each subroutine.
Hints: The polynomial $P_{\mathscr{C}} \in \mathbb{F}_{q}[x]$ is defined in the lecture slides. Observe that your algorithm needs to work for an arbitrary $\xi \in \mathbb{F}_{q}$, not only for $\xi \in\{0,1\}$. Also observe that the given input is $\mathscr{C}, q$, and $\xi$. In particular, the polynomials $a_{1}, a_{2}, \ldots, a_{n / 2}$ need to be constructed inside your algorithm.
4. Delegating matrix multiplication. Suppose you have two $n \times n$ matrices, $X$ and $Y$, with entries in a finite field $F$ with at least four elements. You want to delegate the task of computing the product matrix $X Y$ to your three friends Alice, Bob, and Charlie so that none of your three friends individually gains any information about the matrices $X$ and $Y$ other than the size parameter $n$. Describe a protocol that employs Alice, Bob, and Charlie to help you so that you obtain the product matrix $X Y$ without you yourself putting in more work than $O\left(n^{2}\right)$ operations in $F$. You can assume you have a subroutine that returns independent uniform random elements of $F$.

Hints: Recall Shamir's secret sharing. Extend each matrix $X, Y$ to a matrix whose entries are polynomials of degree at most one with coefficients in $F$, where the constant of each polynomial is the original matrix entry. Have Alice, Bob, and Charlie each multiply a pair of $n \times n$ matrices $X^{(\mathrm{A})}, Y^{(\mathrm{A})}, X^{(\mathrm{B})}, Y^{(\mathrm{B})}$, and $X^{(\mathrm{C})}, Y^{(\mathrm{C})}$ with entries
in $F$. Recover the product matrix $X Y$ by interpolation from the products $X^{(\mathrm{A})} Y^{(\mathrm{A})}$, $X^{(\mathrm{B})} Y^{(\mathrm{B})}$, and $X^{(\mathrm{C})} Y^{(\mathrm{C})}$ that Alice, Bob, and Charlie supply to you. Carefully justify that each of your friends on his or her own does not gain any information about $X$ and $Y$ other than the size parameter $n$.

Deadline and submission instructions. This problem set is due no later than Sunday 10 March 2019, 20:00 (8pm), Finnish time. Please submit your solutions as a single PDF file via e-mail to the lecturer (petteri.kaski (atsymbol)aalto.fi). Please use the precise title

CS-E4500 Problem Set 6: [your-student-number]
with "[your-student-number]" replaced by your student number. For example, assuming that my student number is 123456 , I would carefully title my e-mail

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