# Random graphs and network statistics 

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## Appendix A

## Probability

Here are some miscellaneous facts from probability theory that are used in the text.

## A. 1 Inequalities

Proposition A.1.1 (Markov's inequality). For any random number $X \geq$ 0 and any $a>0$,

$$
\mathbb{P}(X \geq a) \leq a^{-1} \mathbb{E} X
$$

Proof. First, note that $\mathbb{P}(X \geq a)=\mathbb{E} 1(X \geq a)$ where $1(A)$ in general denotes the indicator of the event $A$. Hence by the linearity of expectation,

$$
a \mathbb{P}(X \geq a)=a \mathbb{E} 1(X \geq a)=\mathbb{E} a 1(X \geq a)
$$

Next, the inequalities

$$
a 1(X \geq a) \leq X 1(X \geq a) \leq X
$$

which are valid for any realization of $X$, and the monotonicity of the expectation imply that

$$
\mathbb{E} a 1(X \geq a) \leq \mathbb{E} X
$$

Hence $a \mathbb{P}(X \geq a) \leq \mathbb{E} X$, and the claim follows.
Proposition A.1.2 (Chebyshev's inequality). For any random number $X$ with a finite mean $\mu=\mathbb{E} X$ and any $a>0$,

$$
\mathbb{P}(|X-\mu| \geq a) \leq a^{-2} \operatorname{Var}(X)
$$

Proof. By applying Markov's inequality for $Y=(X-\mu)^{2}$, we find that

$$
\begin{aligned}
\mathbb{P}(|X-\mu| \geq a) & =\mathbb{P}\left((X-\mu)^{2} \geq a^{2}\right) \\
& \leq\left(a^{2}\right)^{-1} \mathbb{E}(X-\mu)^{2}=a^{-2} \operatorname{Var}(X)
\end{aligned}
$$

The following inequality is due to the Finnish-born Wassily Hoeffding.
Proposition A.1.3. Let $S_{n}=\sum_{i=1}^{n} X_{i}$ where the summands are independent and bounded by $a_{i} \leq X_{i} \leq b_{i}$. Then for any $t>0$,

$$
\begin{aligned}
& \mathbb{P}\left(S_{n} \geq \mathbb{E} S_{n}+t\right) \leq e^{-\frac{2 t^{2}}{\sum_{i}\left(b_{i}-a_{i}\right)^{2}}} \\
& \mathbb{P}\left(S_{n} \leq \mathbb{E} S_{n}-t\right) \leq e^{-\frac{2 t^{2}}{\sum_{i}\left(b_{i}-a_{i}\right)^{2}}}
\end{aligned}
$$

and

$$
\mathbb{P}\left(\left|S_{n}-\mathbb{E} S_{n}\right| \geq t\right) \leq 2 e^{-\frac{2 t^{2}}{\sum_{i}\left(b_{i}-a_{i}\right)^{2}}}
$$

Proof. A well-written proof of the first inequality, based on an extremality property related to convex stochastic orders, is available in the original research article Hoeffding [5]. The second inequality follows by applying the first inequality to $\tilde{S}_{n}=-S_{n}$ and the third inequality follows from the first two by the union bound.

## A. 2 Weak convergence of probability measures

Let $\mu, \mu_{1}, \mu_{2}, \ldots$ be probability distributions on $\mathbb{R}$. We say that $\mu_{n} \rightarrow \mu$ weakly if $\int \phi(x) \mu_{n}(d x) \rightarrow \int \phi(x) \mu(d x)$ for every bounded continuous function $\phi: \mathbb{R} \rightarrow \mathbb{R}$. We say that $\mu_{n} \rightarrow \mu$ weakly and with $k$-th moments, if in addition $\mu_{n}$ and $\mu$ have finite $k$-th moments and $\int|x|^{k} \mu_{n}(d x) \rightarrow \int|x|^{k} \mu(d x)$. The sequence $\left(\mu_{n}\right)$ is called uniformly integrable if $\sup _{n} \int|x| \mu_{n}(d x) 1(|x|>$ $K) \rightarrow 0$ as $K \rightarrow \infty$. Let $X, X_{1}, X_{2}, \ldots$ be real-valued random variables. We say that $X_{n} \rightarrow X$ weakly (resp. with weakly with $k$-th moments) if the corresponding probability distributions converge weakly (resp. weakly with $k$-th moments). We say that ( $X_{n}$ ) is uniformly integrable if the collection of corresponding probability distributions is uniformly integrable.

Lemma A.2.1. Let $X_{n}$ and $X$ be random numbers such that $X_{n} \rightarrow X$ weakly with 1st moments. Then the sequence $\left(X_{n}\right)$ is uniformly integrable.

Proof. Given $\epsilon>0$, by Lebesgue's dominated convergence we may choose $K>0$ such that $E X 1(X>K) \leq \epsilon / 3$. Then let $\phi_{K}$ be a continuous bounded function such that $\phi_{K}(x)=x$ for $x \leq K$ and $\phi_{K}=0$ for $x \geq K+1$. Then

$$
x 1(x \leq K) \leq \phi_{K}(x) \leq x 1(x \leq K+1)
$$

so that

$$
\begin{aligned}
\mathbb{E} X_{n} 1\left(X_{n}>K+1\right) & =\mathbb{E} X_{n}-\mathbb{E} X_{n} 1\left(X_{n} \leq K+1\right) \\
& \leq \mathbb{E} X_{n}-\mathbb{E} \phi_{K}\left(X_{n}\right) \\
& =\mathbb{E} X_{n}-\mathbb{E} \phi_{K}(X)+\mathbb{E} \phi_{K}(X)-\mathbb{E} \phi_{K}\left(X_{n}\right) \\
& \leq \mathbb{E} X_{n}-\mathbb{E} X 1(X \leq K)+\mathbb{E} \phi_{K}(X)-\mathbb{E} \phi_{K}\left(X_{n}\right) \\
& =\mathbb{E} X 1(X>K)+\mathbb{E} X_{n}-\mathbb{E} X+\mathbb{E} \phi_{K}(X)-\mathbb{E} \phi_{K}\left(X_{n}\right) \\
& \leq \epsilon / 3+\left|\mathbb{E} X_{n}-\mathbb{E} X\right|+\left|\mathbb{E} \phi_{K}\left(X_{n}\right)-\mathbb{E} \phi_{K}(X)\right| .
\end{aligned}
$$

Then we may choose $n_{0}$ so large that $\left|\mathbb{E} X_{n}-\mathbb{E} X\right| \leq \epsilon / 3$ and $\left|\mathbb{E} \phi_{K}\left(X_{n}\right)-\mathbb{E} \phi_{K}(X)\right| \leq$ $\epsilon / 3$ for all $n>n_{0}$. Hence $\mathbb{E} X_{n} 1\left(X_{n}>K+1\right) \leq \epsilon$ for all $n>n_{0}$. Furthermore, for every $1 \leq m \leq n_{0}$ we may choose, again by Lebesgue's dominated convergence, $K_{m}$ such that $\mathbb{E} X_{m} 1\left(X_{m}>K_{m}\right) \leq \epsilon$. Now if we choose $L=\max \left\{K+1, K_{1}, \ldots, K_{n_{0}}\right\}$, it follows that $\sup _{n} \mathbb{E} X_{n} 1\left(X_{n}>L\right) \leq \epsilon$.

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