

Computational Algebraic Geometry

Elimination theory

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Definition

Given $I = \langle f_1, \dots, f_s \rangle \subset k[x_1, \dots, x_n]$ the l -th **elimination ideal** I_l is the ideal of $k[x_{l+1}, \dots, x_n]$ defined by

$$I_l = I \cap k[x_{l+1}, \dots, x_n].$$

Theorem (The Elimination Theorem)

Let $I \subset k[x_1, \dots, x_n]$ be an ideal and let G be a Groebner basis of I wrt to lex order where $x_1 > x_2 > \dots > x_n$. Then, for every $0 \leq l \leq n$, the set

$$G_l = G \cap k[x_{l+1}, \dots, x_n]$$

is a Groebner basis of the l -th elimination ideal I_l .

Theorem (The Extension Theorem)

Let $I = \langle f_1, \dots, f_s \rangle \subseteq \mathbb{C}[x_1, \dots, x_n]$ and let I_1 be the first elimination ideal of I . For each $1 \leq i \leq s$, write f_i in the form

$$f_i = g_i(x_2, \dots, x_n)x_1^{N_i} + \text{terms in which } x_1 \text{ has degree} < N_i,$$

where $N_i \geq 0$ and $g_i \in \mathbb{C}[x_2, \dots, x_n]$ is nonzero. Suppose that we have a *partial solution* $(a_2, \dots, a_n) \in V(I_1)$. If $(a_2, \dots, a_n) \notin V(g_1, \dots, g_s)$, then there exists $a_1 \in \mathbb{C}$ such that $(a_1, a_2, \dots, a_n) \in V(I)$.

Let $I = \langle xy - 1 \rangle \subseteq \mathbb{C}[x, y]$. Fix the lex order with $x > y$.

- 1 What is the first elimination ideal I_1 ?
- 2 What is the set of partial solutions $\mathbb{V}(I_1)$?
- 3 Which partial solutions in $\mathbb{V}(I_1)$ extend to a complete solution in $\mathbb{V}(I)$?

Today:

- The geometry of elimination
- Implicitization

The geometry of elimination

Variety of the elimination ideal

Let π_I be the projection map

$$\begin{aligned}\pi_I : \quad \mathbb{C}^n &\rightarrow \mathbb{C}^{n-I}, \\ (a_1, \dots, a_n) &\mapsto (a_{I+1}, \dots, a_n).\end{aligned}$$

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Lemma

Let $V = \mathbb{V}(f_1, \dots, f_s) \subseteq \mathbb{C}^n$ and let $I_I = \langle f_1, \dots, f_s \rangle \cap \mathbb{C}[x_{I+1}, \dots, x_n]$ be the I -th elimination ideal of $\langle f_1, \dots, f_s \rangle$. Then, in \mathbb{C}^{n-I} we have

$$\pi_I(V) \subseteq \mathbb{V}(I_I).$$

Proof: let $f \in I_e$. let $(a_1, \dots, a_n) \in V$. Then

$$f(a_1, \dots, a_n) = 0, \text{ because } f \in \langle f_1, \dots, f_s \rangle.$$

Since f includes only the variable x_{e+1}, \dots, x_n , then

$$f(a_1, \dots, a_n) = f(\pi_e(a_1, \dots, a_n)) = 0.$$

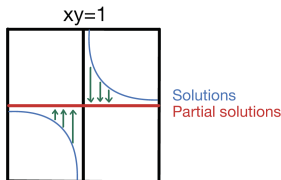
Hence f vanishes on every point of $\pi_e(V)$. \square

Variety of the elimination ideal

- Recall: The points of $\mathbb{V}(I_I)$ are called **partial solutions**.
- By Lemma, we can write $\pi_I(V)$ as

$$\pi_I(V) = \{(a_{l+1}, \dots, a_n) \in \mathbb{V}(I_I) : \exists a_1, \dots, a_l \in \mathbb{C} \text{ with } (a_1, \dots, a_l, a_{l+1}, \dots, a_n) \in V\}.$$

- $\pi_I(V)$ consists precisely of the **partial solutions that extend to complete solutions**.



$$\pi_1(V) = \{a \in \mathbb{C} : a \neq 0\} \Rightarrow \pi_1(V) \text{ is not an affine variety!}$$

The Geometric Extension Theorem

Theorem (The Geometric Extension Theorem)

Let $V = \mathbb{V}(f_1, \dots, f_s) \subseteq \mathbb{C}^n$, let g_i be as in the Extension Theorem. If I_1 is the first elimination ideal of $\langle f_1, \dots, f_s \rangle$, then we have the equality in \mathbb{C}^{n-1}

$$\mathbb{V}(I_1) = \pi_1(V) \cup (\mathbb{V}(g_1, \dots, g_s) \cap \mathbb{V}(I_1)).$$

- $\pi_1(V)$ fills up $\mathbb{V}(I_1)$ besides possibly a part that lies in $(\mathbb{V}(g_1, \dots, g_s) \cap \mathbb{V}(I_1))$

Proof: This theorem follows from the Extension Theorem and the previous lemma. \square

The Geometric Extension Theorem

It is not clear how big the missing part is and it can be unnaturally large:

- $(y - z)x^2 + xy - 1$ and $(y - z)x^2 + xz - 1$ generate the same ideal as $yx - 1$ and $zx - 1$
- $I_1 = \langle y - z \rangle$
- The partial solutions are $\{(a, a) : a \in \mathbb{C}\}$

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- $I_1 = \langle y - z \rangle$
- The partial solutions are $\{(a, a) : a \in \mathbb{C}\}$
- The first set of generators: $g_1 = g_2 = (y - z)$ and hence the Geometric Extension Theorem says nothing about the size of $\pi_1(V)$
- The second set of generators: $g_1 = y$ and $g_2 = z$ and hence all partial solutions besides $(0, 0)$ extend to a complete solution

The Closure Theorem

Theorem (The Closure Theorem)

Let $V = \mathbb{V}(f_1, \dots, f_s) \subseteq \mathbb{C}^n$ and let I_l be the l -th elimination ideal of $\langle f_1, \dots, f_s \rangle$. Then

- 1 $\mathbb{V}(I_l)$ is the smallest affine variety containing $\pi_l(V) \subseteq \mathbb{C}^{n-l}$.
- 2 When $V \neq \emptyset$, there is an affine variety $W \subsetneq \mathbb{V}(I_l)$ such that $\mathbb{V}(I_l) - W \subseteq \pi_l(V)$.

Proof: 1) We will postpone the proof of the first part.
 2) We prove it for the special case $l=1$.

By the Geometric Extension Theorem

$$V(I_1) = \pi_1(V) \cup \underbrace{(V(g_1, \dots, g_s) \cap V(I_1))}_W.$$

Hence $V(I_1) - W \subseteq \pi_1(V)$ and we are done if $W \neq V(I_1)$.

If $W = V(I_1)$, we need to modify eqns defining V so that W becomes smaller.

OBSERVATION: If $W = V(I_1)$, then
 $V = V(f_1, \dots, f_s, g_1, \dots, g_s)$.

Proof of obs: " \supseteq " is clear. " \subseteq " let $(a_1, \dots, a_n) \in V$.

Then $(a_2, \dots, a_n) \in \pi_1(V) \subseteq V(I_1) = W$. Hence g_i 's vanish on (a_2, \dots, a_n) and $(a_1, \dots, a_n) \in V(f_1, \dots, f_s, g_1, \dots, g_s)$.

Let $\tilde{I} = \langle f_1, \dots, f_s, g_1, \dots, g_s \rangle$. Then I and \tilde{I} might be different and hence I_1 and \tilde{I}_1 might be different. However, $V(I_1) = V(\tilde{I}_1)$ because they are both the smallest varieties containing $\pi_1(V)$.

Next we define a better basis for \tilde{I} . let

$\tilde{f}_i = f_i - g_i x_1^{N_i}$. Then $\tilde{I} = \langle \tilde{f}_1, \dots, \tilde{f}_s, g_1, \dots, g_s \rangle$.
 The geom. Ext. Then for $V = V(\tilde{f}_1, \dots, \tilde{f}_s, g_1, \dots, g_s)$
 gives

$$V(I_1) = V(\tilde{I}_1) = \pi_1(V) \cup \tilde{W}$$

where \tilde{W} consists of the partial solutions where the leading coefficients of $\tilde{f}_1, \dots, \tilde{f}_s, g_1, \dots, g_s$ vanish.

If $\tilde{W} \subsetneq W$, then we are done. It is not guaranteed that \tilde{W} is strictly smaller. In this case we can repeat the process. Each time the degrees of x_1 drop (or remain zero), so that eventually the generators will have degree 0 in x_1 . This means that V can be defined by the vanishing of pol's in $\mathbb{C}[x_2, \dots, x_n]$ and every partial solution (a_2, \dots, a_n) extends to a solution (a_1, \dots, a_n) for any $a_1 \in \mathbb{C}$. In this case we can choose $W = \emptyset$. \square

The Closure Theorem

Theorem (The Closure Theorem)

Let $V = \mathbb{V}(f_1, \dots, f_s) \subseteq \mathbb{C}^n$ and let I_l be the l -th elimination ideal of $\langle f_1, \dots, f_s \rangle$. Then

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- 2 When $V \neq \emptyset$, there is an affine variety $W \subsetneq \mathbb{V}(I_l)$ such that $\mathbb{V}(I_l) - W \subseteq \pi_l(V)$.

Corollary

Let $V \subseteq \mathbb{C}^n$ be an affine variety. Then there are affine varieties $Z_i \subset W_i \subseteq \mathbb{C}^{n-1}$ for $i \leq 1 \leq p$ such that

$$\pi_l(V) = \bigcup_{i=1}^p (W_i - Z_i).$$

Sets of this form are called **constructible**.

The Closure Theorem

Corollary

Let $V = \mathbb{V}(f_1, \dots, f_s) \subseteq \mathbb{C}^n$, and assume that for some i , f_i can be written as

$$f_i = c_i x_1^N + \text{terms in which } x_1 \text{ has degree } < N,$$

where $c \in \mathbb{C}$ is nonzero and $N > 0$. If I_1 is the first elimination ideal, then

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The Extension and Closure Theorems hold over any
algebraically closed field.

Implicitization

Rational parametrizations

Definition

A **rational function** in t_1, \dots, t_m with coefficients in k is a quotient f/g of two polynomials $f, g \in k[t_1, \dots, t_m]$, where g is not the zero polynomial. The set of all rational functions is denoted $k(t_1, \dots, t_m)$.

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- **rational parametric representation** of V consists of $r_1, \dots, r_n \in k(t_1, \dots, t_m)$ such that

$$x_1 = r_1(t_1, \dots, t_m), \dots, x_n = r_n(t_1, \dots, t_m)$$

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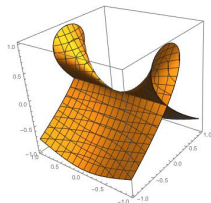
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lie in V

- require that V is the **smallest variety** containing these points
- if r_1, \dots, r_n are polynomials, then **polynomial parametric representation**

Parametric vs implicit form

- original defining equations $f_1 = \dots = f_s = 0$ are called an **implicit representation**
- it is easy to **draw** a parametric description of a curve on a computer



- plotted not using $x^2 - y^2z^2 + z^3 = 0$ but

$$x = t(u^2 - t^2), y = u, z = u^2 - t^2.$$

- if we want to know **whether the point $(1, 2, -1)$ is on the above surface**, then implicit representation is useful:

$$1^2 - 2^2(-1)^2 + (-1)^3 = 1 - 4 - 1 = -4$$

Parametric vs implicit form

Desirability of having both representations leads to the questions

- **(Parametrization)** Does every affine variety have a rational parametric description?
- **(Implicitization)** Given a parametric representation of an affine variety, can we find the defining equations (i.e. can we find an implicit representation)?
- The answer to the first question is no. Those that can be parametrized are called **unirational**.
- It is difficult to tell whether a given variety is unirational or not.
- We will learn that the answer to the second question is always yes.

Implicitization

- The parametrization **might not fill up** all of the variety.
- Implicitization asks for the defining equations of the smallest variety containing the parametrization.
- How to find **missing points**?

Twisted cubic

- the twisted cubic has parametrization

$$x = t, y = t^2, z = t^3$$

- the tangent vector to the curve at a point is $(1, 2t, 3t^2)$
- the tangent line is parametrized

$$(t, t^2, t^3) + u(1, 2t, 3t^2) = (t + u, t^2 + 2tu, t^3 + 3t^2u)$$

- a parametrization of the entire surface is

$$x = t + u, y = t^2 + 2tu, z = t^3 + 3t^2u$$

- the tangent surface lies on the variety V defined by

$$-4x^3z + 3x^2y^2 - 4y^3 + 6xyz - z^2 = 0$$

- Is V the smallest variety containing the tangent surface?
- If yes, does the tangent surface fill up V completely?

Polynomial parametrization

- We consider the polynomial parametrization

$$x_1 = f_1(t_1, \dots, t_m), \dots, x_n = f_n(t_1, \dots, t_m),$$

where $f_1, \dots, f_n \in k[t_1, \dots, t_m]$.

- We can think of it as the map F

$$\begin{aligned} k^m &\rightarrow k^n, \\ (t_1, \dots, t_m) &\mapsto (f_1(t_1, \dots, t_m), \dots, f_n(t_1, \dots, t_m)) \end{aligned}$$

- Then $F(k^m) \subseteq k^n$ is equal to the subset of k^n parametrized by the polynomials f_1, \dots, f_n .
- The solution to the implicitization problem finds the smallest algebraic variety containing $F(k^m)$

Polynomial parametrization

Next we want to **connect the implicitization and elimination**. The polynomial equations

$$x_1 = f_1(t_1, \dots, t_m), \dots, x_n = f_n(t_1, \dots, t_m),$$

define a variety

$$V = \mathbb{V}(x_1 - f_1, \dots, x_n - f_n) \subseteq k^{n+m}.$$

The points of V can be written as

$$(t_1, \dots, t_m, f_1(t_1, \dots, t_m), \dots, f_n(t_1, \dots, t_m)).$$

Hence V is the graph of the map F .

Polynomial parametrization

We also have the inclusion $i : k^m \rightarrow k^{n+m}$ defined by

$$i(t_1, \dots, t_m) \mapsto (t_1, \dots, t_m, f_1(t_1, \dots, t_m), \dots, f_n(t_1, \dots, t_m))$$

and the projection $\pi_m : k^{n+m} \rightarrow k^n$ defined by

$$\pi_m(t_1, \dots, t_m, x_1, \dots, x_n) = (x_1, \dots, x_n).$$

A commutative diagram with three nodes. The top node is k^{n+m} . The bottom-left node is k^m . The bottom-right node is k^n . An arrow labeled i (in blue) points from k^m to k^{n+m} . An arrow labeled π_m (in blue) points from k^{n+m} to k^n . A horizontal arrow labeled F (in blue) points from k^m to k^n .

$$i(k^m) = V \text{ and } F(k^m) = \pi_m(i(k^m)) = \pi_m(V)$$

Polynomial Implicitization

Theorem (Polynomial Implicitization)

If k is an infinite field, let $F : k^m \rightarrow k^n$ be the map defined by the polynomial parametrization

$$x_1 = f_1(t_1, \dots, t_m), \dots, x_n = f_n(t_1, \dots, t_m).$$

Let $I = \langle x_1 - f_1, \dots, x_n - f_n \rangle \subseteq k[t_1, \dots, t_m, x_1, \dots, x_n]$ and let $I_m = I \cap k[x_1, \dots, x_n]$ be the m -th elimination ideal. Then $\mathbb{V}(I_m)$ is the smallest variety in k^n containing $F(k^m)$.

Proof: let $V = V(I) \subseteq k^{n+m}$, the graph of F . If $k = \mathbb{C}$, then $F(\mathbb{C}^m) = \pi_m(V)$ and by the Closure Theorem $V(I_m)$ is the smallest variety containing $\pi_m(V)$.

Suppose k is an infinite subfield of \mathbb{C} . let $V_k(I_m)$ be the variety in k^n and $V_{\mathbb{C}}(I_m)$ be the variety in \mathbb{C}^n . We have $F(k^m) = \pi_m(V_k) \subseteq V_k(I_m)$.

let $Z_k = V_k(g_1, \dots, g_s) \subseteq k^n$ be any variety containing $F(k^m)$. We must show that $V_k(I_m) \subseteq Z_k$. Since $g_i = 0$ on Z_k , also $g_i = 0$ on $F(k^m)$. Hence $g_i \circ F$ vanishes on k^m , i.e. it is the zero function on k^m .

Since k is an infinite field, then $g_i \circ F$ is a zero function in $k[t_1, \dots, t_m]$. Hence also $g_i \circ F$ vanishes on \mathbb{C}^m and thus g_i 's vanish on $F(\mathbb{C}^m)$. Hence $Z_{\mathbb{C}} = V_{\mathbb{C}}(g_1, \dots, g_s)$ is a variety containing $F(\mathbb{C}^m)$. Since the theorem true for \mathbb{C} , it follows that $V_{\mathbb{C}}(I_m) \subseteq Z_{\mathbb{C}}$ in \mathbb{C}^n . It follows $V_k(I_m) \subseteq Z_k$. \square

Implicitization algorithm

- Let

$$x_1 = f_1(t_1, \dots, t_m), \dots, x_n = f_n(t_1, \dots, t_m).$$

for polynomials $f_1, \dots, f_n \in k[t_1, \dots, t_m]$.

- Consider the ideal $I = \langle x_1 - f_1, \dots, x_n - f_n \rangle$.
- Compute the Groebner basis of I with respect to a lexicographic order where every t_i is greater than every x_j .
- The elements of the Groebner basis not involving t_1, \dots, t_m define the smallest variety in k^n containing the parametrization.

Implicitization algorithm

Tangent surface of the twisted cubic:

- $I = \langle x - t - u, y - t^2 - 2tu, z - t^3 - 3t^2u \rangle \subseteq \mathbb{R}[t, u, x, y, z]$
- Fix the lex order with $t > u > x > y > z$
- A Groebner basis is given by

$$g_1 = t + u - x,$$

$$g_2 = u^2 - x^2 + y,$$

$$g_3 = ux^2 - uy - x^3 + (3/2)xy - (1/2)z,$$

$$g_4 = uxy - uz - x^2y - xz + 2y^2,$$

$$g_5 = uxz - uy^2 + x^2z - (1/2)xy^2 - (1/2)yz,$$

$$g_6 = uy^2 - uz^2 - 2x^2yz + (1/2)xy^3 - xz^2 + (5/2)y^2z,$$

$$g_7 = x^3z - (3/4)x^2y^2 - (3/2)xyz + y^3 + (1/4)z^2.$$

- Since g_7 is the only Groebner basis element consisting of variables x, y, z only, then $\mathbb{V}(g_7)$ solves the implicitization problem.

Tangent surface of the twisted cubic

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Quiz: Which partial solutions $(x, y, z) \in \mathbb{V}(I_2) = \mathbb{V}(g_7) \subseteq \mathbb{C}^3$ extend to a solution of $\mathbb{V}(I) \subseteq \mathbb{C}^5$?

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- $I_1 = \langle g_2, \dots, g_7 \rangle$ is the first elimination ideal of I_2 and the coefficient of u^2 in g_2 is 1 \Rightarrow all partial solutions in $\mathbb{V}(I_2)$ extend to a solution in $\mathbb{V}(I_1)$

Tangent surface of the twisted cubic

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Quiz: Which partial solutions $(x, y, z) \in \mathbb{V}(I_2) = \mathbb{V}(g_7) \subseteq \mathbb{C}^3$ extend to a solution of $\mathbb{V}(I) \subseteq \mathbb{C}^5$?

- $I_1 = \langle g_2, \dots, g_7 \rangle$ is the first elimination ideal of I_2 and the coefficient of u^2 in g_2 is 1 \Rightarrow all partial solutions in $\mathbb{V}(I_2)$ extend to a solution in $\mathbb{V}(I_1)$
- the coefficient of t in g_1 is 1 \Rightarrow all partial solutions in $\mathbb{V}(I_1)$ extend to a solution in $\mathbb{V}(I)$

Tangent surface of the twisted cubic

- Since all partial solutions in $\mathbb{V}(I_2) = \mathbb{V}(g_2)$ extend to a complete solution in $\mathbb{V}(I)$, then the tangent surface of the twisted cubic fills $\mathbb{V}(g_2)$.

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- It can be shown that in this example every **real partial solution** $(x, y, z) \in \mathbb{V}(I_2)$ extends to a **real complete solution** in $\mathbb{V}(I)$.

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- It can be shown that in this example every **real partial solution** $(x, y, z) \in \mathbb{V}(I_2)$ extends to a **real complete solution** in $\mathbb{V}(I)$.
- In general there is no easy answer to the question whether a parametrization fills the minimal variety containing it and each case has to be analyzed separately.

Rational parametrization

$$x = \frac{u^2}{v}, y = \frac{v^2}{u}, z = u$$

- (x, y, z) always lies on the surface $x^2y = z^3$
- **clearing the denominators** in the above parametrization gives the ideal

$$I = \langle vx - u^2, uy - v^2, z - u \rangle$$

- the second elimination ideal is $I_2 = \langle z(x^2y - z^3) \rangle$
- $\mathbb{V}(I_2) = \mathbb{V}(x^2y - z^3) \cup \mathbb{V}(z)$
- hence $\mathbb{V}(I_2)$ is **not** the smallest variety containing the parametrization

Rational parametrization

- We consider a **rational parametrization**

$$x_1 = \frac{f_1(t_1, \dots, t_m)}{g_1(t_1, \dots, t_m)}, \dots, x_n = \frac{f_n(t_1, \dots, t_m)}{g_n(t_1, \dots, t_m)},$$

where $f_1, g_1, \dots, f_n, g_n \in k[t_1, \dots, t_m]$.

- The map $F : k^m \rightarrow k^n$ given by

$$F(t_1, \dots, t_m) = \left(\frac{f_1(t_1, \dots, t_m)}{g_1(t_1, \dots, t_m)}, \dots, \frac{f_n(t_1, \dots, t_m)}{g_n(t_1, \dots, t_m)} \right)$$

might not be defined everywhere because of the denominators.

- Let $W = \mathbb{V}(g_1 g_2 \cdots g_n) \subset k^m$.
- Then $F : k^m - W \rightarrow k^n$ is well-defined.

Rational parametrization

$$\begin{array}{ccc} & k^{u+m} & \\ i \nearrow & & \searrow \pi_m \\ k^m - W & \xrightarrow{F} & k^u \end{array}$$

- $i(k^m - W) \subseteq \mathbb{V}(I)$ where $I = \langle g_1 x_1 - f_1, \dots, g_n x_n - f_n \rangle$
- $\mathbb{V}(I)$ might not be the smallest variety containing $i(k^m - W)$
- add an extra variable to control the denominators:

$$J = \langle g_1 x_1 - f_1, \dots, g_n x_n - f_n, 1 - gy \rangle \subseteq k[y, t_1, \dots, t_m, x_1, \dots, x_n],$$

where $g = g_1 \cdot g_2 \cdots g_n$

Rational parametrization

$$\begin{array}{ccc} & k^{n+m+1} & \\ j \nearrow & & \searrow \pi_{m+1} \\ k^m - W & \xrightarrow{F} & k^n \end{array}$$

The map $j : k^m - W \rightarrow k^{n+m+1}$ is defined by

$$j(t_1, \dots, t_m) = \left(\frac{1}{g(t_1, \dots, t_m)}, t_1, \dots, t_m, \frac{f_1(t_1, \dots, t_m)}{g_1(t_1, \dots, t_m)}, \dots, \frac{f_n(t_1, \dots, t_m)}{g_n(t_1, \dots, t_m)} \right).$$

We have $j(k^m - W) = \mathbb{V}(J)$ and

$$F(k^m - W) = \pi_{m+1}(j(k^m - W)) = \pi_{m+1}(\mathbb{V}(J)).$$

Rational Implicitization

Theorem (Rational Implicitization)

If k is an infinite field, let $F : k^m - W \rightarrow k^n$ be the function defined by the rational parametrization

$$x_1 = \frac{f_1(t_1, \dots, t_m)}{g_1(t_1, \dots, t_m)}, \dots, x_n = \frac{f_n(t_1, \dots, t_m)}{g_n(t_1, \dots, t_m)}.$$

Let

$J = \langle g_1 x_1 - f_1, \dots, g_n x_n - f_n, 1 - gy \rangle \subseteq k[y, t_1, \dots, t_m, x_1, \dots, x_n]$, where $g = g_1 \cdot g_2 \cdots g_n$, and let $J_{m+1} = J \cap k[x_1, \dots, x_n]$ be the $(m+1)$ -st elimination ideal. Then $\mathbb{V}(J_{m+1})$ is the smallest variety in k^n containing $F(k^m - W)$.

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This theorem gives an **implicitization algorithm** for rational parametrizations similarly to the polynomial parametrizations case.

Implicitization algorithm

- $x = \frac{u^2}{v}, y = \frac{v^2}{u}, z = u$
- $J = \langle vx - u^2, uy - v^2, z - u, 1 - uvw \rangle$
- the third elimination ideal is $J_3 = \langle x^2y - z^3 \rangle$

Guiding questions

We started the **Groebner bases** chapter with four guiding questions:

- The ideal description problem: Does every ideal $I \subset k[x_1, \dots, x_n]$ have a finite generating set?
- The ideal membership problem: Given $f \in k[x_1, \dots, x_n]$ and ideal $I = \langle f_1, \dots, f_s \rangle$, determine if $f \in I$.
- The problem of solving polynomial equations: Find all common solutions in k^n of a system of polynomial equations

$$f_1(x_1, \dots, x_n) = \dots = f_n(x_1, \dots, x_n) = 0.$$

- The implicitization problem: If V is given by a rational parametric representation, find a system of polynomial equations that defines V .

Solving systems of polynomial equations

Not all systems of polynomial equations have nice solutions:

$$xy = 4, y^2 = x^3 - 1$$

The Groebner basis for the lex order with $x > y$ is given by

$$g_1 = 16x - y^2 - y^4, g_2 = y^5 + y^3 - 64$$

The polynomial g_2 has **no rational roots**. One can obtain numerically

$$y = 2.21363, -1.78719 \pm 1.3984i, 0.680372 \pm 2.26969i.$$

These solutions can be substituted into g_1 to find the values of x . We will return to the topic when we talk about **numerical algebraic geometry**.

Today:

- The geometry of elimination
- Implicitization

Next time:

- Hilbert's Nullstellensatz
- Rational ideals
- Ideal-variety correspondence