- **Weeks 1–2:** **informal introduction**
  - network = path

- **Week 3:** **graph theory**
- **Weeks 4–7:** **models of computing**
  - what can be computed (efficiently)?
- **Weeks 8–11:** **lower bounds**
  - what cannot be computed (efficiently)?
- **Week 12:** **recap**

Week 6

– CONGEST model: bandwidth limitations

CONGEST model

- **LOCAL model: arbitrarily large messages**
- **CONGEST model: $O(\log n)$-bit messages**

CONGEST model

- **Any of these can be encoded in $O(\log n)$-bit messages:**
  - node identifier
  - number of nodes
  - number of edges
  - distance between two nodes …

CONGEST model

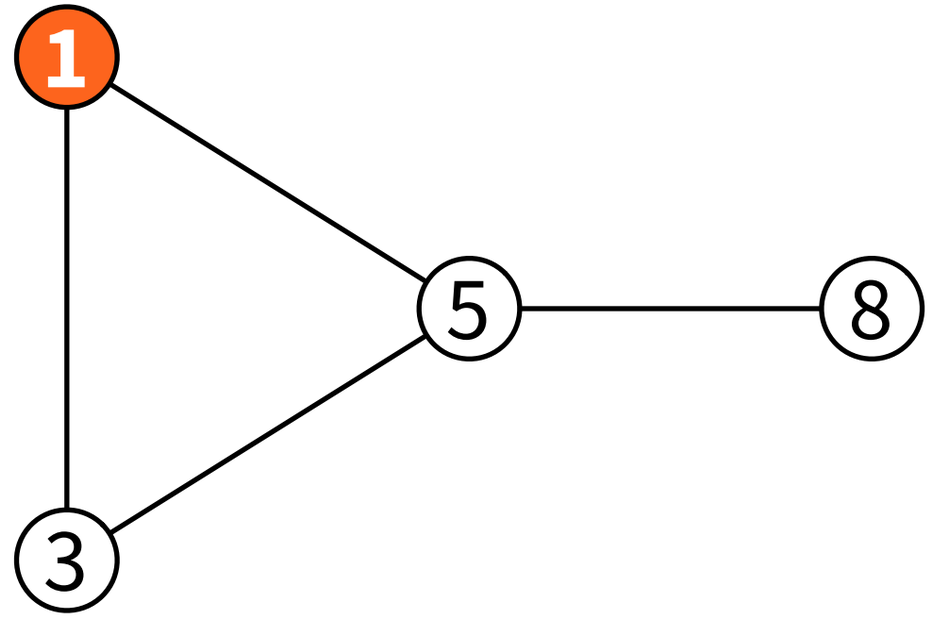
- **Many algorithms that we have seen only send small messages**
  - can be used directly in the CONGEST model
- **Exception: algorithm Gather**
  - may need to send $O(n^2)$-bit messages

CONGEST model

- **$O(n)$ time trivial in the LOCAL model**
  - brute force approach: Gather + solve locally
- **$O(n)$ time non-trivial in the CONGEST model**
- **Today: how to find all-pairs shortest paths in $O(n)$ time**

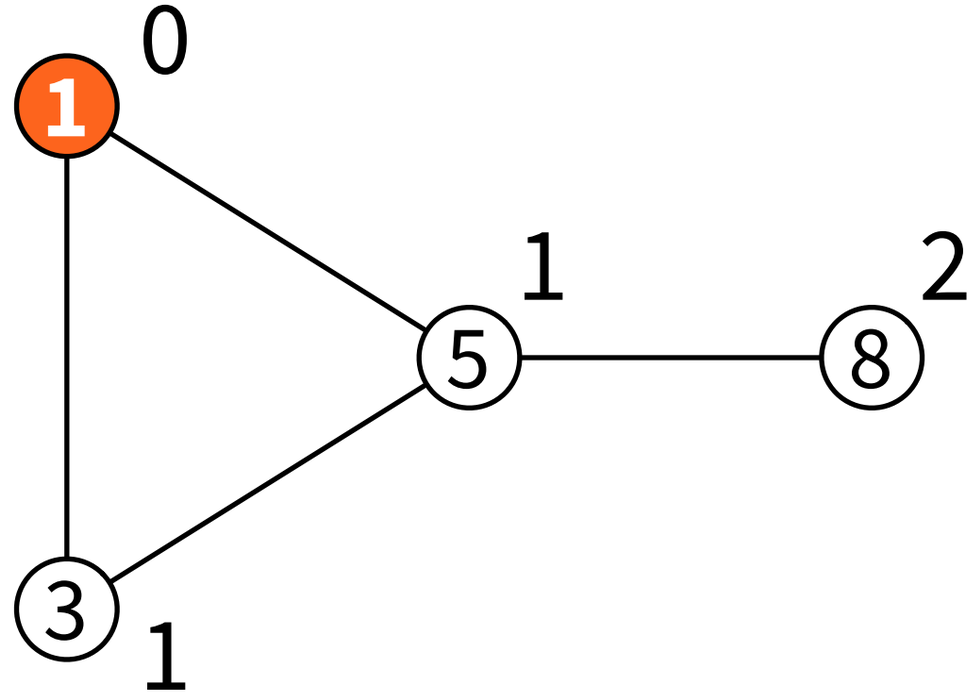
Single-source shortest paths

Input:



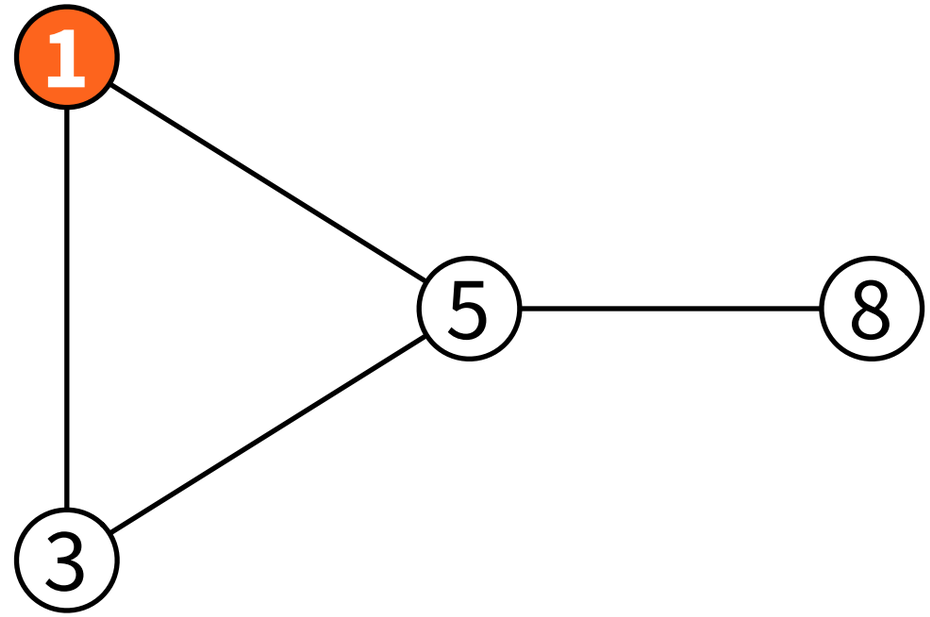
Single-source shortest paths

Output:



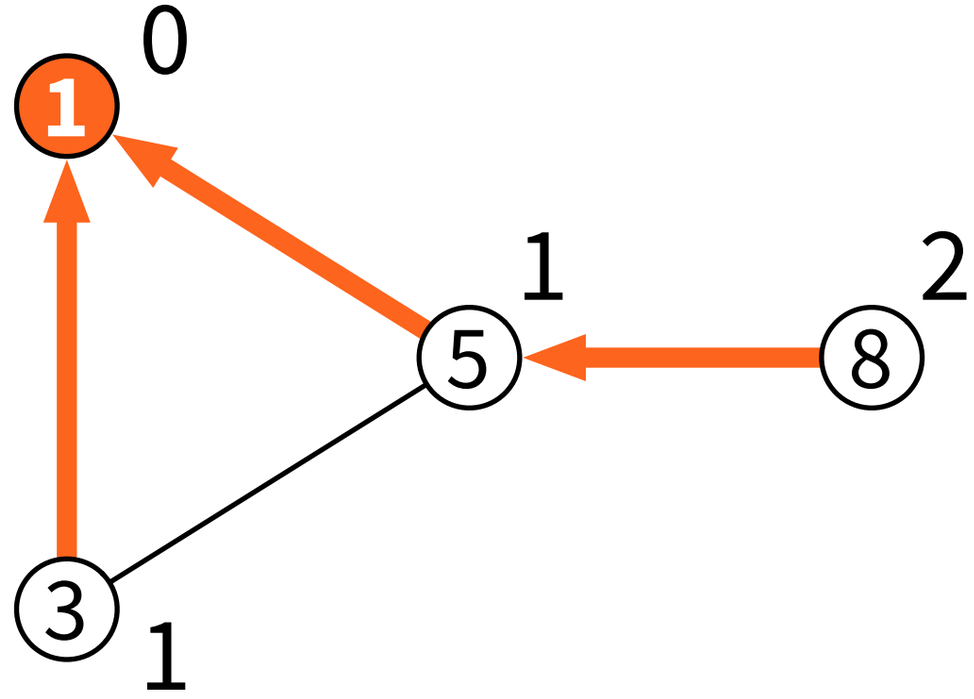
BFS tree

Input:



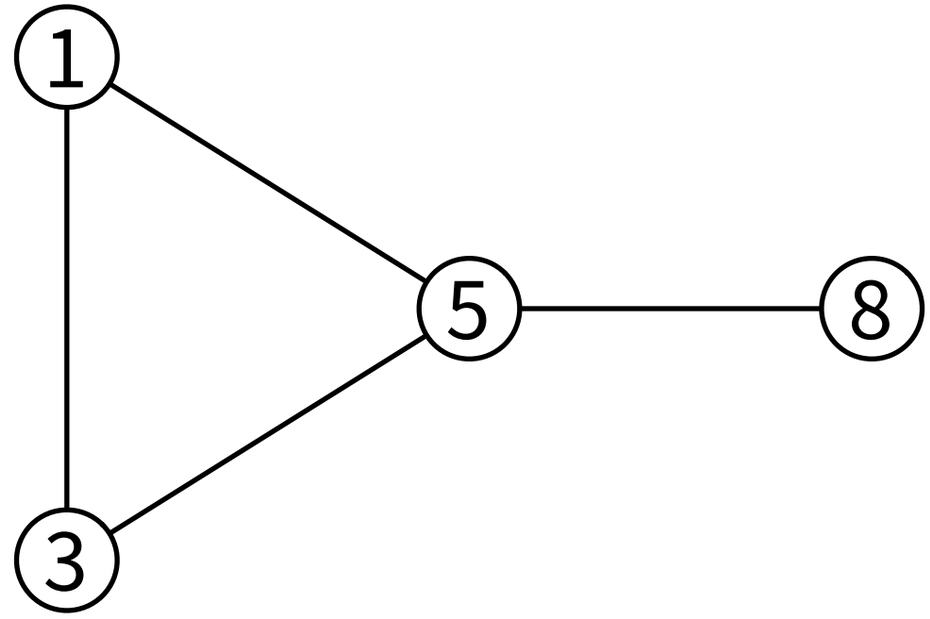
BFS tree

Output:



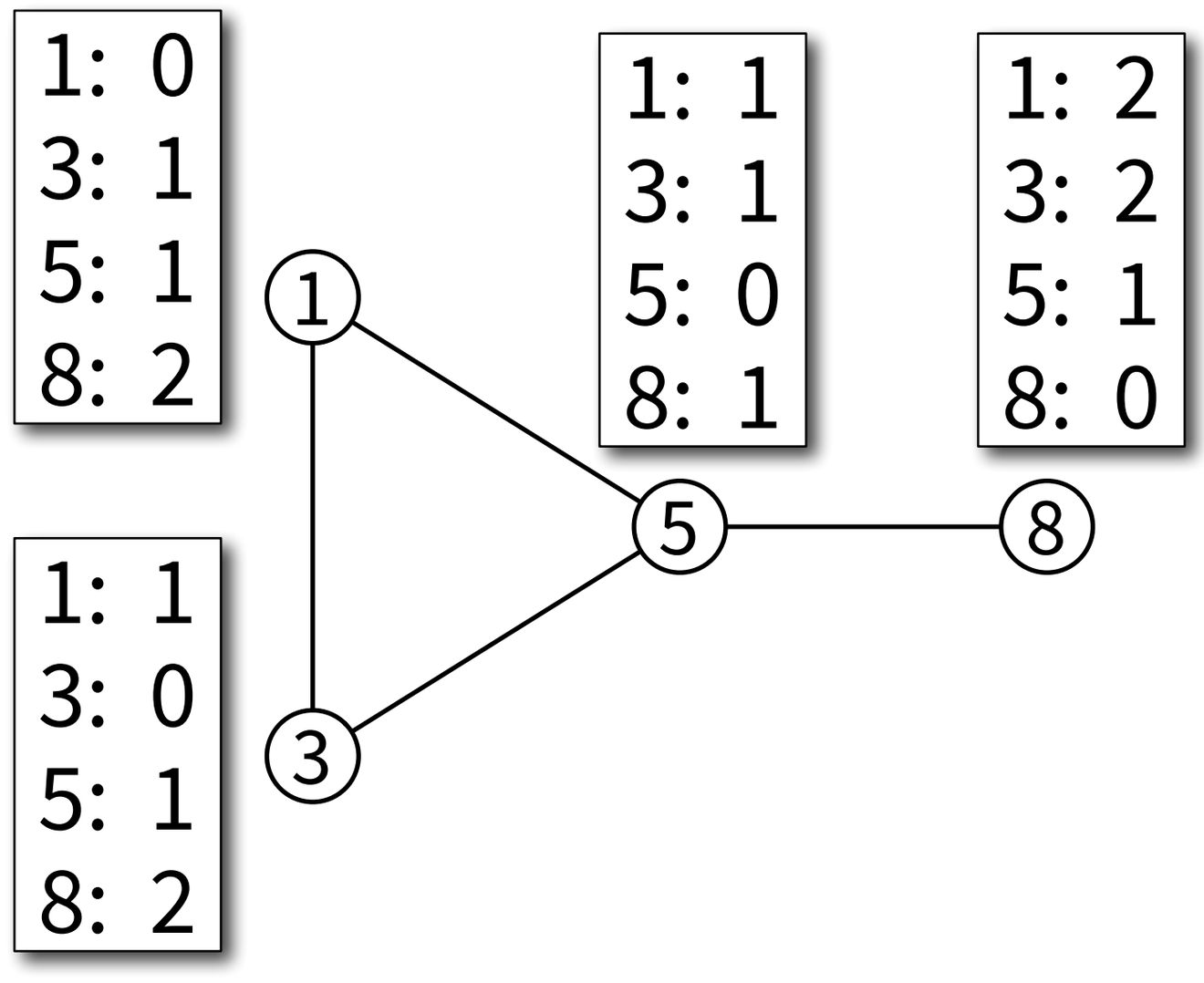
All-pairs shortest paths

Input:



All-pairs shortest paths

Output:

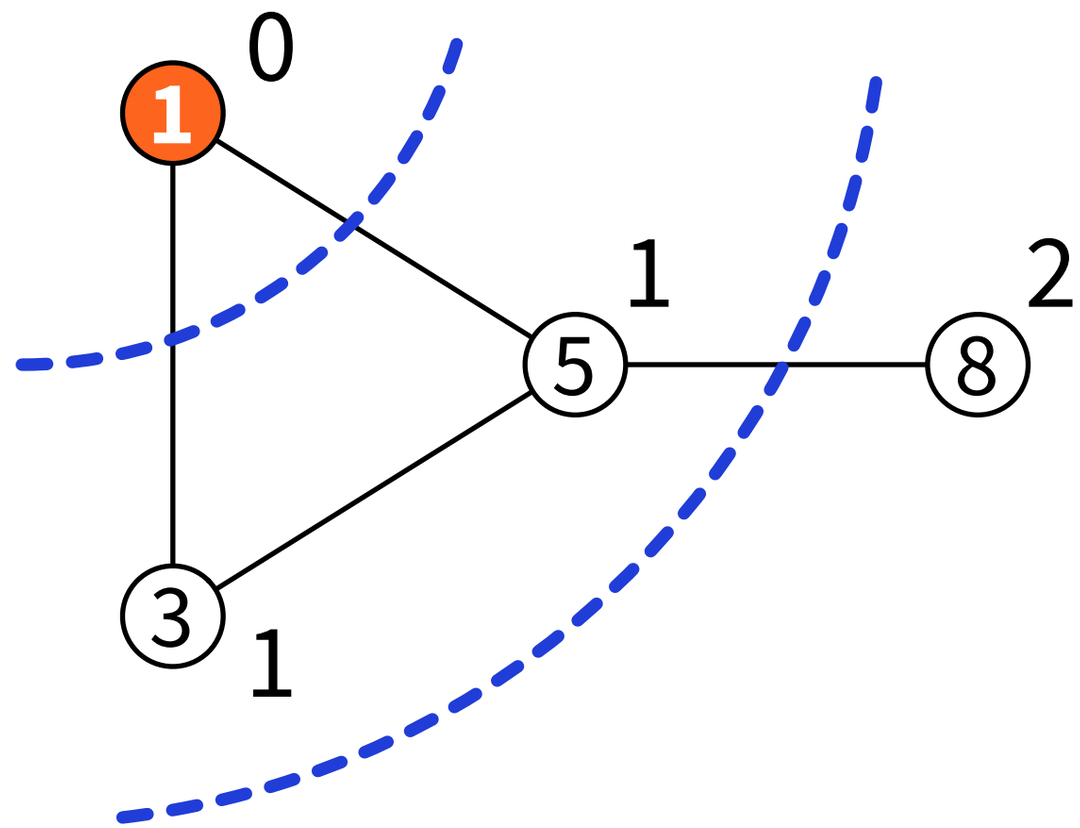


Algorithm Wave

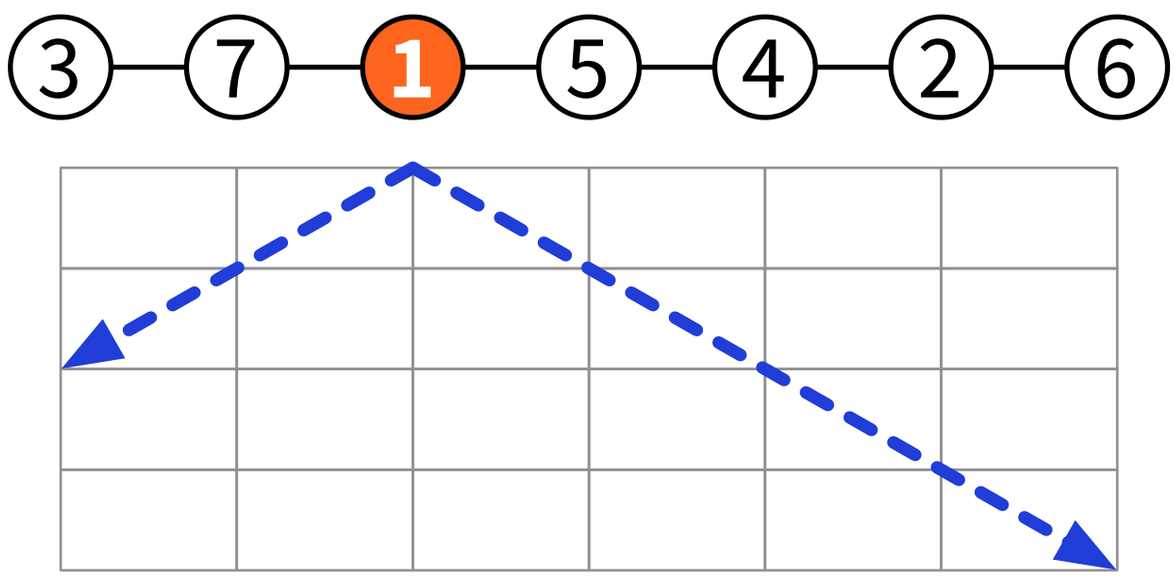
- **Solves single-source shortest paths in time $O(\mathrm{diam}(G))$**
- **Leader creates a ‘wave’, other nodes propagate it**
- **Wave first received in round t: distance to leader is t**

Algorithm Wave

Output:



Algorithm Wave

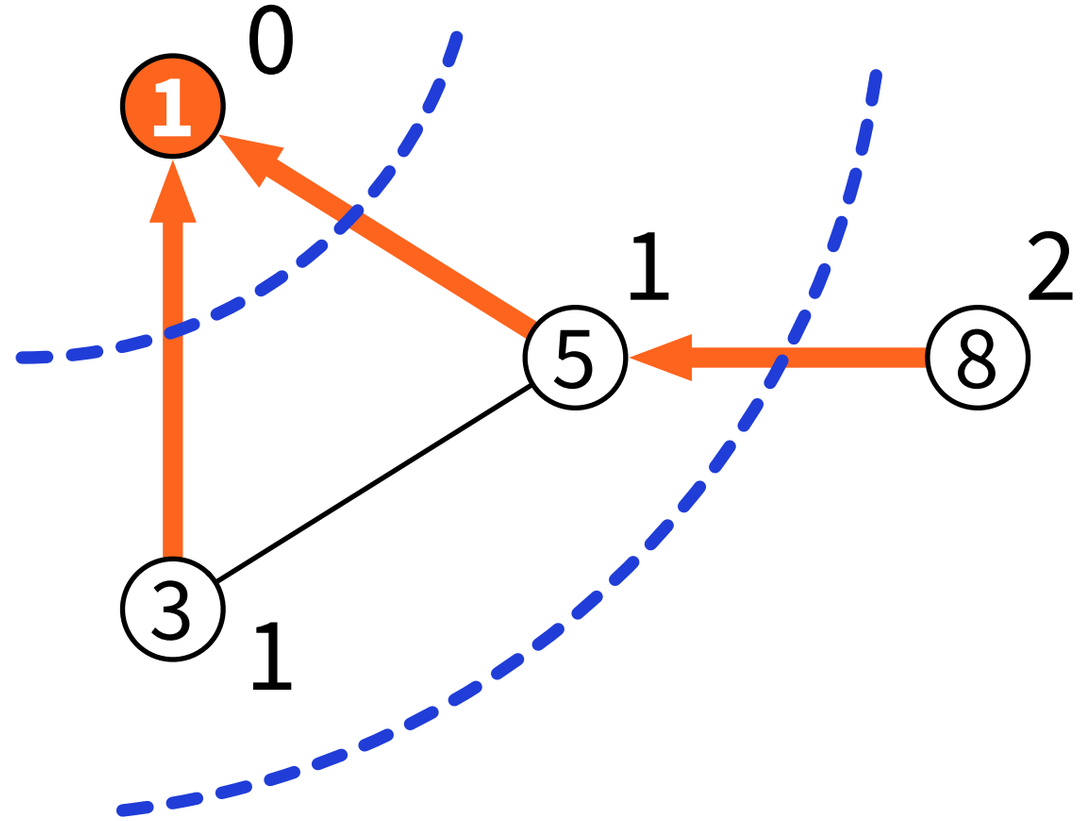


Algorithm BFS

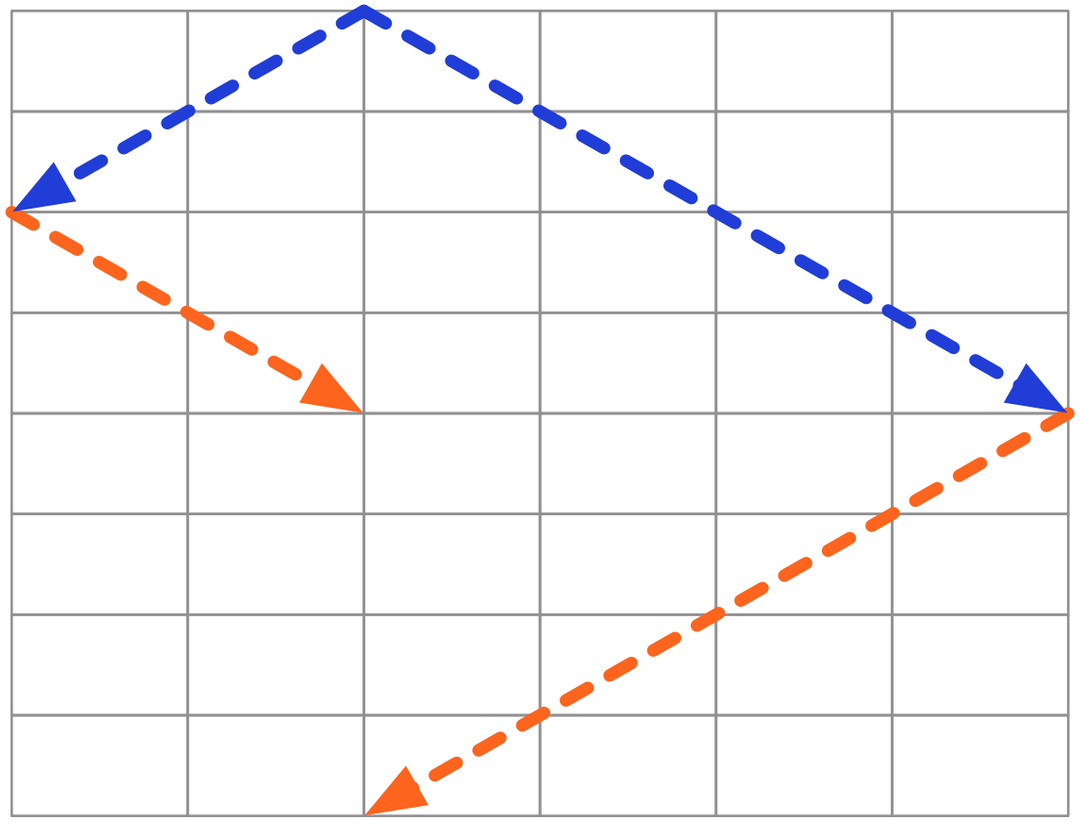
- **Wave + handshakes**
- **Tree construction:**
  - “proposal” + “accept”
  - everyone knows their parent & children
- **Acknowledgements back from leaf nodes**

Algorithm BFS

Output:



Algorithm BFS



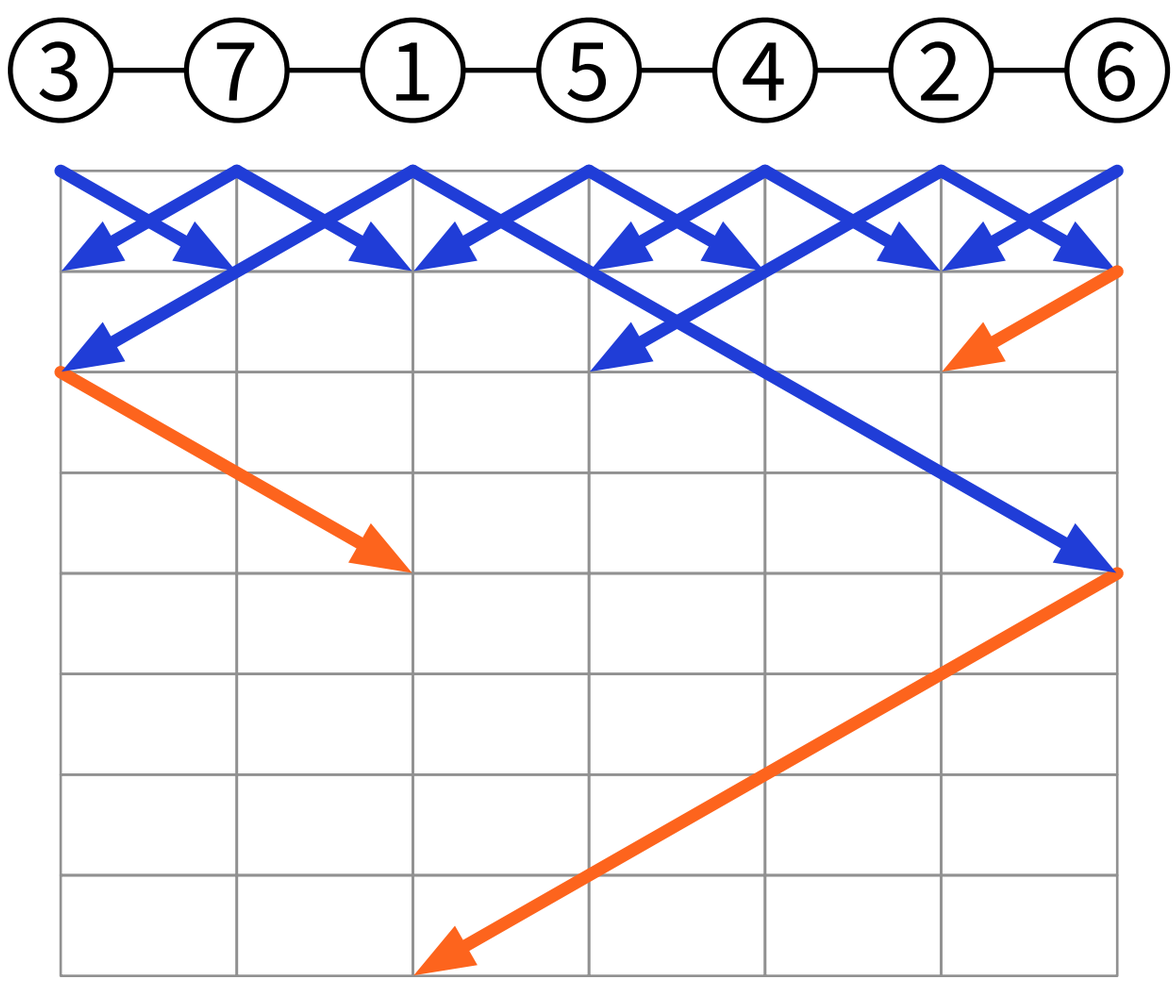
propose,
accept

ack

Algorithm Leader

- **Each node creates a separate BFS process**
- **When two BFS processes “collide”, the one with the smaller root “wins”**
  - each node only needs to send messages related to ***one BFS process***
- **One tree wins everyone else → leader**

Algorithm Leader



propose, accept

ack

Recap until now

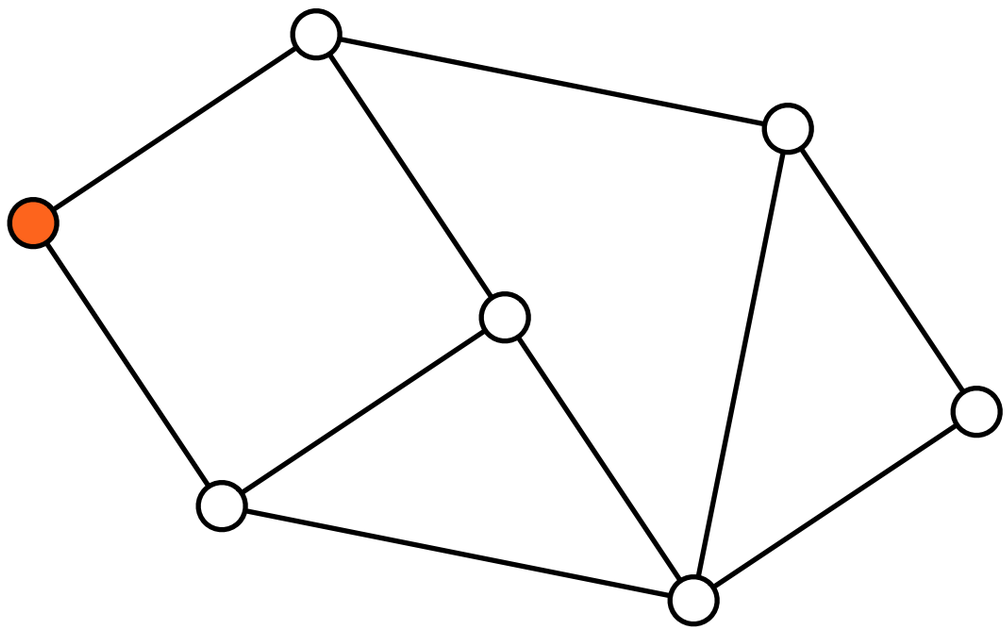
- **SSSP: Wave algorithm**
- **BFS tree: Wave algorithm + acceptance**
- **Leader election: Many BFS in parallel**
- **All these problems can be solved in $O(\mathrm{diam}(G))$ rounds in the CONGEST model**

Algorithm APSP

- **Basic idea: run Wave from each node**
- **Challenge: congestion**
  - all waves parallel → too many bits per edge
  - all waves sequentially → takes too long
- **Solution: pipelining**

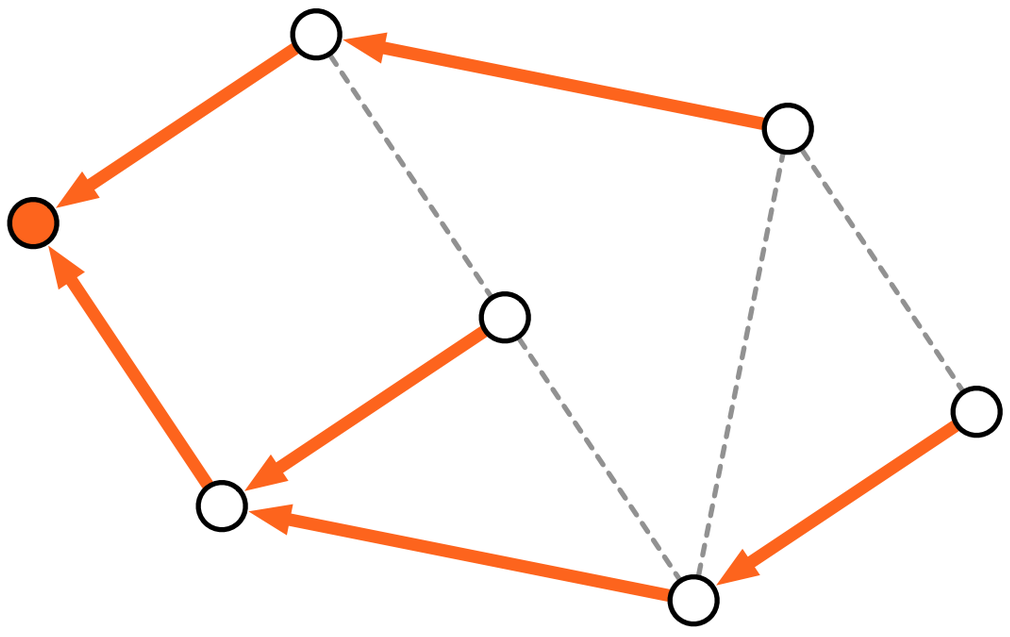
Algorithm APSP

- **Elect leader**



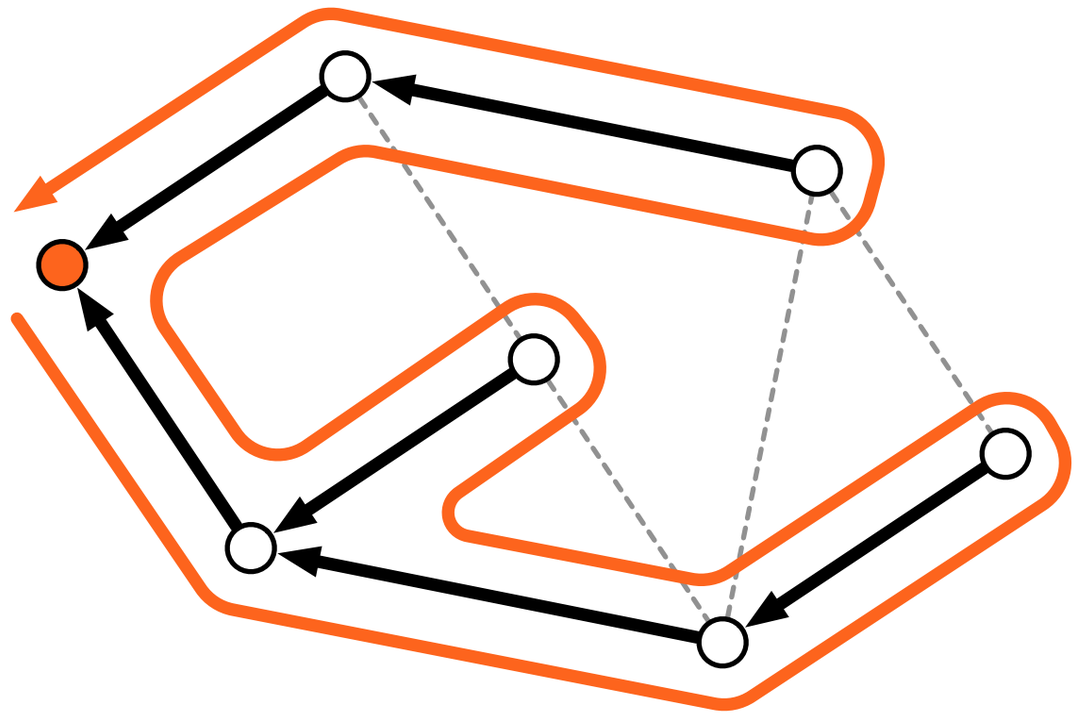
Algorithm APSP

- **Elect leader, construct BFS tree**



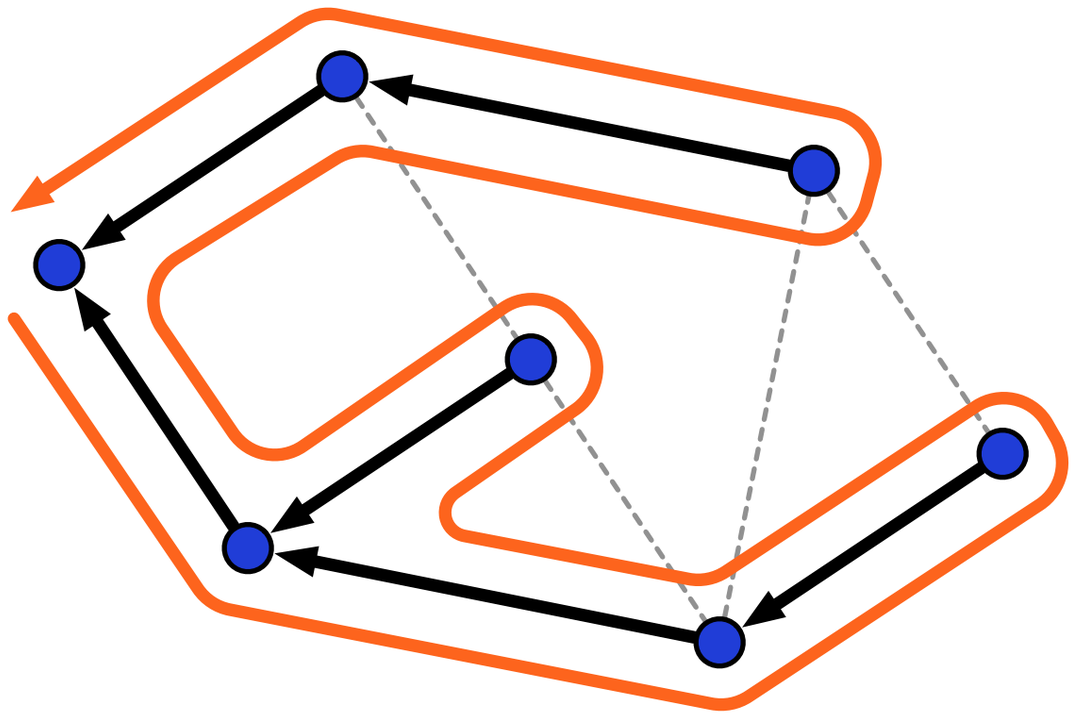
Algorithm APSP

- **Move token along BFS tree slowly (every 2 rounds)**



Algorithm APSP

- **Create *wave* every time we visit a new node**



Algorithm APSP

- **Algorithm animation:**
  http://users.ics.aalto.fi/suomela/apsp/

- $|E|$ in a BFS tree: ***n* - 1**
- Token traverses **2** times each edge of the BFS tree
- Total number of rounds: **2(2(*n* - 1)) + *O*(diam(*G*))**

Pipelining

- ***n* operations, each operation takes time *t***
- **Parallel: *t* rounds, bad congestion**
- **Sequential: *nt* rounds, no congestion**
- **Pipelining: *n* + *t* rounds, no congestion**

Summary

- **LOCAL model: unlimited bandwidth**
- **CONGEST model: $O(\log n)$ bandwidth**
- **$O(n)$ or $O(\mathrm{diam}(G))$ time is no longer trivial**
- **Example: all-pairs shortest paths in time $O(n)$, pipelining helps**

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