## Polyphase Decomposition

## The Decomposition

- Consider an arbitrary sequence $\{x[n]\}$ with a z-transform $X(z)$ given by

$$
X(z)=\sum_{n=-\infty}^{\infty} x[n] z^{-n}
$$

- We can rewrite $X(z)$ as

$$
X(z)=\sum_{k=0}^{M-1} z^{-k} X_{k}\left(z^{M}\right)
$$

where

$$
\begin{aligned}
& X_{k}(z)=\sum_{n=-\infty}^{\infty} x_{k}[n] z^{-n}=\sum_{n=-\infty}^{\infty} x[M n+k] z^{-n} \\
& 0 \leq k \leq M-1 \\
& \text { Copvighut ozno, }, \text { s. . Mitra }
\end{aligned}
$$

## Polyphase Decomposition

- The subsequences $\left\{x_{k}[n]\right\}$ are called the polyphase components of the parent sequence $\{x[n]\}$
- The functions $X_{k}(z)$, given by the $z$-transforms of $\left\{x_{k}[n]\right\}$, are called the polyphase components of $X(z)$

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## Polyphase Decomposition

- The relation between the subsequences $\left\{x_{k}[n]\right\}$ and the original sequence $\{x[n]\}$ are given by

$$
x_{k}[n]=x[M n+k], \quad 0 \leq k \leq M-1
$$

- In matrix form we can write

$$
X(z)=\left[\begin{array}{llll}
1 & z^{-1} & \cdots & z^{-(M-1)}
\end{array}\right]\left[\begin{array}{c}
X_{0}\left(z^{M}\right) \\
X_{1}\left(z^{M}\right) \\
\vdots \\
X_{M-1}\left(z^{M}\right)
\end{array}\right]
$$

## Polyphase Decomposition

- The polyphase decomposition of an FIR transfer function can be carried out by inspection
- For example, consider a length-9 FIR transfer function:

$$
H(z)=\sum_{n=0}^{8} h[n] z^{-n}
$$

## Polyphase Decomposition

- Its 4-branch polyphase decomposition is given by

$$
H(z)=E_{0}\left(z^{4}\right)+z^{-1} E_{1}\left(z^{4}\right)+z^{-2} E_{2}\left(z^{4}\right)+z^{-3} E_{3}\left(z^{4}\right)
$$

where
$E_{0}(z)=h[0]+h[4] z^{-1}+h[8] z^{-2}$
$E_{1}(z)=h[1]+h[5] z^{-1}$
$E_{2}(z)=h[2]+h[6] z^{-1}$
$E_{3}(z)=h[3]+h[7] z^{-1}$
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## Polyphase Decomposition

- The polyphase decomposition of an IIR transfer function $H(z)=P(z) / D(z)$ is not that straight forward
- One way to arrive at an $M$-branch polyphase decomposition of $H(z)$ is to express it in the form $P^{\prime}(z) / D^{\prime}\left(z^{M}\right)$ by multiplying $P(z)$ and $D(z)$ with an appropriately chosen polynomial and then apply an $M$-branch polyphase decomposition to $P^{\prime}(z)$
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## Polyphase Decomposition

- Note: The above approach increases the overall order and complexity of $H(z)$
- However, when used in certain multirate structures, the approach may result in a more computationally efficient structure
- An alternative more attractive approach is discussed in the following example


## Polyphase Decomposition

- Therefore $H(\mathrm{z})$ can be expressed as

$$
H(z)=E_{0}\left(z^{2}\right)+z^{-1} E_{1}\left(z^{2}\right)
$$

where

$$
\begin{aligned}
& E_{0}(z)=\frac{1}{2}\left(\frac{0.105573+z^{-1}}{1+0.105573 z^{-1}}\right) \\
& E_{1}(z)=\frac{1}{2}\left(\frac{0.52786+z^{-1}}{1+0.52786 z^{-1}}\right)
\end{aligned}
$$

## Polyphase Decomposition

- Note: In the above polyphase decomposition, branch transfer functions $E_{i}(z)$ are stable allpass functions
- Moreover, the decomposition has not increased the order of the overall transfer function $H(z)$


## FIR Filter Structures Based on Polyphase Decomposition

- We shall demonstrate later that a parallel realization of an FIR transfer function $H(z)$ based on the polyphase decomposition can often result in computationally efficient multirate structures
- Consider the $M$-branch Type I polyphase decomposition of $H(z)$ :

$$
H(z)=\sum_{k=0}^{M-1} z^{-k} E_{k}\left(z^{M}\right)
$$

## FIR Filter Structures Based on Polyphase Decomposition

- The transpose of the Type I polyphase FIR filter structure is indicated below



## FIR Filter Structures Based on Polyphase Decomposition

- A direct realization of $H(z)$ based on the Type II polyphase decomposition is shown below



## Computationally Efficient Decimators

- Using the cascade equivalence \#1 we arrive at the computationally efficient decimator structure shown below on the right



## Computationally Efficient Decimators

- It is thus necessary to compute $v[n]$ at

$$
n=\ldots,-2 M,-M, 0, M, 2 M, \ldots
$$

- Computational requirements are therefore $N$ multiplications and ( $N-1$ ) additions per output sample being computed
- However, as $n$ increases, stored signals in the delay registers change


## Computationally Efficient Decimators and Interpolators

- However, here the arithmetic units are operative at all instants of the output sampling period which is $1 / M$ times that of the input sampling period
- Similar savings are also obtained in the case of the interpolator structure developed using the polyphase decomposition


## Computationally Efficient Decimators

- To illustrate the computational efficiency of the modified decimator structure, assume $H(z)$ to be a length $-N$ structure and the input sampling period to be $T=1$
- Now the decimator output $y[n]$ in the original structure is obtained by downsampling the filter output $v[n]$ by a factor of M


## Computationally Efficient Decimators

- Hence, all computations need to be completed in one sampling period, and for the following ( $M-1$ ) sampling periods the arithmetic units remain idle
- The modified decimator structure also requires $N$ multiplications and ( $N-1$ ) additions per output sample being computed


## Computationally Efficient Interpolators

- Figures below show the computationally efficient interpolator structures



## Computationally Efficient Decimators and Interpolators

- More efficient interpolator and decimator structures can be realized by exploiting the symmetry of filter coefficients in the case of linear-phase filters $H(z)$
- Consider for example the realization of a factor-of-3 ( $M=3$ ) decimator using a length-12 Type 1 linear-phase FIR lowpass filter


## Computationally Efficient Decimators and Interpolators

- Note that $E_{1}(z)$ still has a symmetric impulse response, whereas $E_{0}(z)$ is the mirror image of $E_{2}(z)$
- These relations can be made use of in developing a computationally efficient realization using only 6 multipliers and 11 two-input adders as shown on the next slide

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## Rational Sampling Rate Converter

- The complexity of the design of the fractional sampling rate converter depends on the ratio of the sampling rates between the input and the output digital signals
- For example, in digital audio applications, the three different sampling frequencies employed are $44.1 \mathrm{kHz}, 32 \mathrm{kHz}$, and 48 kHz


## Computationally Efficient Decimators and Interpolators

- The corresponding transfer function is $H(z)=h[0]+h[1] z^{-1}+h[2] z^{-2}+h[3] z^{-3}+h[4] z^{-4}+h[5] z^{-5}$
$+h[5] z^{-6}+h[4] z^{-7}+h[3] z^{-8}+h[2] z^{-9}+h[1] z^{-10}+h[0] z^{-11}$
- A conventional polyphase decomposition of $H(z)$ yields the following subfilters:
$E_{0}(z)=h[0]+h[3] z^{-1}+h[5] z^{-2}+h[2] z^{-3}$
$E_{1}(z)=h[1]+h[4] z^{-1}+h[4] z^{-2}+h[1] z^{-3}$
$E_{2}(z)=h[2]+h[5] z^{-1}+h[3] z^{-2}+h[0] z^{-3}$
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## Computationally Efficient Decimators and Interpolators

- Factor-of-3 decimator with a linear-phase decimation filter

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## Rational Sampling Rate Converter

- As a consequence there are three different values for the sampling rate conversion factor:
- 2:3 (or 3:2), 147:160 (or 160:147), and 320:441 (or 441:320)
- Likewise, in digital video applications, the sampling rates of composite video signals are 14.3181818 MHz and 17.734475 MHz


## Rational Sampling Rate Converter

- The sampling rates of the digital component video signal are 13.5 MHz and 6.75 MHz for the luminance and the color-difference signals, respectively, for the NTSC and PAL systems
- Here, again there are different sampling rate conversion factors


## Rational Sampling Rate Converter

- The general structure for a rational sampling rate converter shown below

can be made computationally efficient by making use of one of the structures based on the polyphase decompositions


## Rational Sampling Rate Converter

with the factor-of- $M$ down-sampler moved to all $L$ branches as shown below


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## Rational Sampling Rate Converter

- Consider the $k$-th branch of the structure shown in the previous slide

$$
\rightarrow E_{k}(\varepsilon)-\mid t-E^{-\pi-z} \longrightarrow \sqrt{M}
$$

- Using the identity $\mu M-\lambda L=1$ we can write

$$
z^{-k}=z^{-k(\mu M-\lambda L)}=z^{-k \mu M}-z^{k \lambda L}
$$

- Hence, we can replace the block of $k$ delays with a block of $k \mu M$ unit delays and a block of $k \lambda L$ unit advances as shown next


## Rational Sampling Rate Converter

- This branch can be further redrawn by invoking the noble identites as shown below
- We next interchange the positions of the upsampler and the down-sampler



## Rational Sampling Rate Converter

- Its equivalent realizable form is as shown below



## Rational Sampling Rate Converter <br> 

- Finally by combining all $k$ branches we arrive at a computationally efficient rational sampling rate converter


## Rational Sampling Rate Converter

- Example - The basic form of a rational sampling rate converter with an interpolation factor $2 / 3$ needed in the conversion of a digital audio signal of 48kHz rate to one of $32-\mathrm{kHz}$ rate is shown below


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## Rational Sampling Rate Converter

- For this design we have $L=2$ and $M=3$
- The identity $3 \mu-2 \lambda=1$ is thus satisfied with $\mu=\lambda=1$
- Hence, the general structure of this converter is as indicated below


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## Rational Sampling Rate Converter

- By realizing the sub-filters $E_{0}(z)$ and $E_{1}(z)$ in Type I polyphase forms and then applying the cascade equivalence we arrive at the final computationally efficient structure shown in the next slide
- Here, all filters operate at the $16-\mathrm{kHz}$ rate

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## A Useful Identity

- The cascade multirate structure shown below appears in a number of applications

$$
x[n] \longrightarrow \uparrow L \longrightarrow H(z) \longrightarrow \bigsqcup^{\prime} \longrightarrow y[n]
$$

- Equivalent time-invariant digital filter obtained by expressing $H(z)$ in its $L$-term Type I polyphase form $\sum_{k=0}^{L-1} z^{-k} E_{k}\left(z^{L}\right)$ is shown below

$$
x[n] \longrightarrow E_{0}(z) \longrightarrow y[n]
$$

