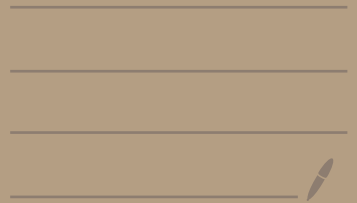
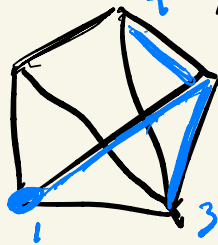
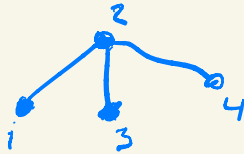

Missing lemmas &
graph polynomials



Missing lemmas about embedding trees

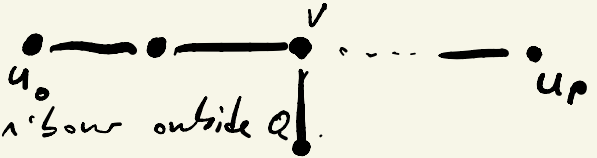
Lm: A graph with $\delta(G) \geq m-1$ contains all trees with m nodes as a subgraph.

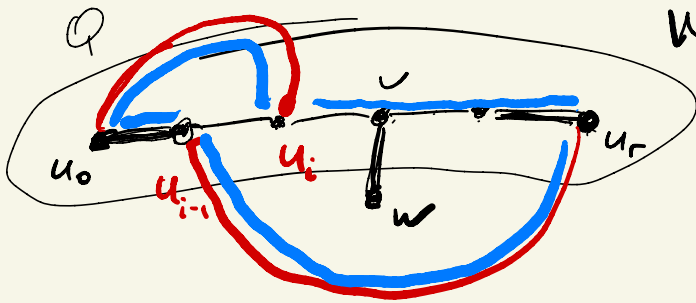
Proof: Embed arbitrarily, keeping the embedded graph connected



Lm: A connected graph on $\geq m$ nodes with min degree $> \frac{m}{2}$ contains a path on m nodes as a subgraph.

Pf: Assume Q longest path in G ,
 $|Q| < m$.
 $\frac{m}{2} + 1$
 v has a neighbour outside Q .





WTS there is a path containing all vertices of Q and $w \in N(v) - Q$.

Enough to show there is a cycle whose vertex set is Q .

Q longest path $\Rightarrow N(u_0) \subseteq Q$
 $N(u_r) \subseteq Q$

Pigeonhole principle:

$$\left\{ \underset{\uparrow}{i} : u_0 u_{i+1} \in E \right\} \cap \left\{ \underset{\uparrow}{i} : u_r u_i \in E \right\}$$

$$\{1 \dots r-2\} \qquad \{1 \dots r-2\}$$

There is a "blue" cycle

$$u_0 u_i Q u_r u_{i-1} Q u_0.$$

Cut this cycle to a path with v as an endpoint and extend the path to w

Longer path. \square

Notice: Not true if ^{we do} not assume connected:



Ramsey numbers of graphs:

Recall $R(n, m) =$ smallest N s.t
any 2-partition of
 $E(K_N)$ has a blue
 K_n or a red K_m

$R(G, H) =$ smallest N s.t if
I 2-colour the edges of

Note

$$R(G, H) \leq R(|G|, |H|).$$

K_N I always get
a blue G or a
red H .

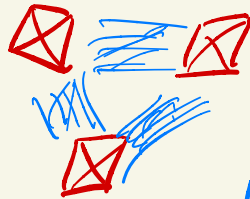
Thm: $R(K_n, T) = (m-1)(n-1) + 1$

\uparrow
m vertex tree

regardless
of the
tree.

Proof

\geq :



cliques
of size
 $m-1$

blue subgraph
is Turan
graph

$K_{m-1, \dots, m-1}$
 $\underbrace{\hspace{2cm}}$
 $n-1$ parts

so has no
 K_n .

\leq :

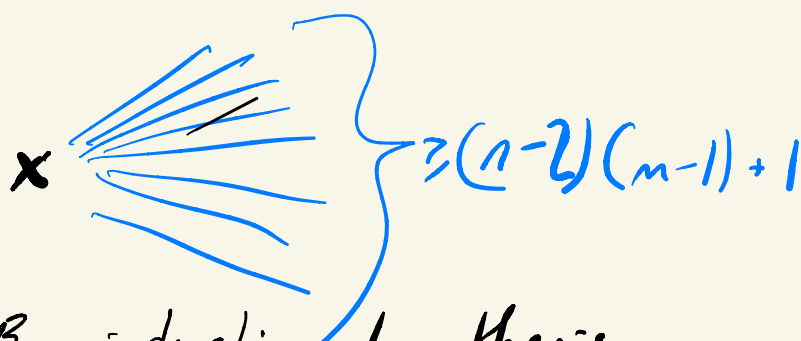
Proof by induction
on n .

(vacuous when $n=1$)
(trivial when $n=2$)

Assume blue-red colouring
of $E(K_{(m-1)(n-1)+1})$.

IF red
subgraph has min degree $\geq m-1$,
then red T by lemma.

So assume some node x has $> (m-1)(n-1) - (m-1)$
blue neighbours $(m-1)(n-2)$



By induction hypothesis,
among $N_{\text{blue}}(x)$ is either
blue K_{n-1} or red T .



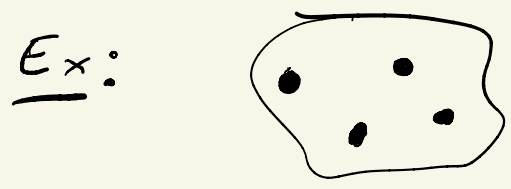
so blue K_n in

$$\{x\} \cup N_{\text{blue}}(x).$$



Chromatic polynomial

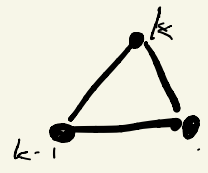
Def: $\chi_G(k) = \#\{k\text{-colourings of } G\}$



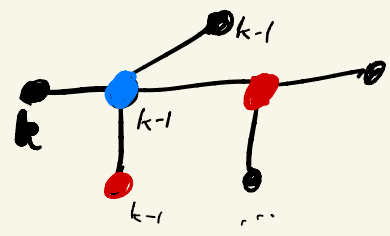
$$\chi_{K_n} = k^n$$



$$\chi_e = k(k-1)$$

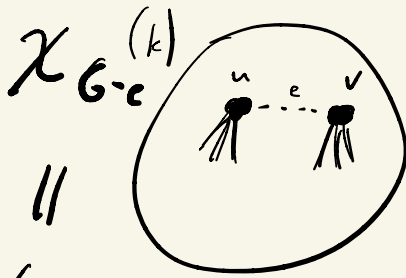


$$\chi_{K_n} = \overbrace{k(k-1) \cdots (k-n+1)} \\ = (k)_n = \frac{k!}{(k-n)!}$$



$$\chi_T \\ \parallel \\ k(k-1)^{n-1} \\ (n = \# \text{ nodes})$$

if $n \leq k$



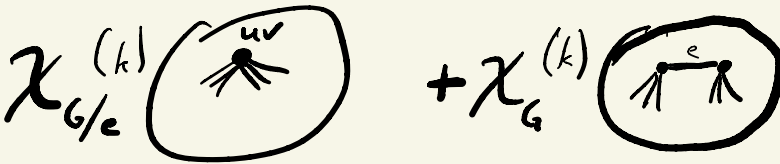
colourings of $G-e$ where u, v get same colour

+

_____ || _____

different colours

||



for any k .

- $\chi_{K_n}^{(k)} = k^n$ ^{degree n}

- $\chi_G^{(k)} = \chi_{G-e}^{(k)} - \chi_{G/e}^{(k)}$ for any edge $e \in E(G)$. ^{degree $n-1$}

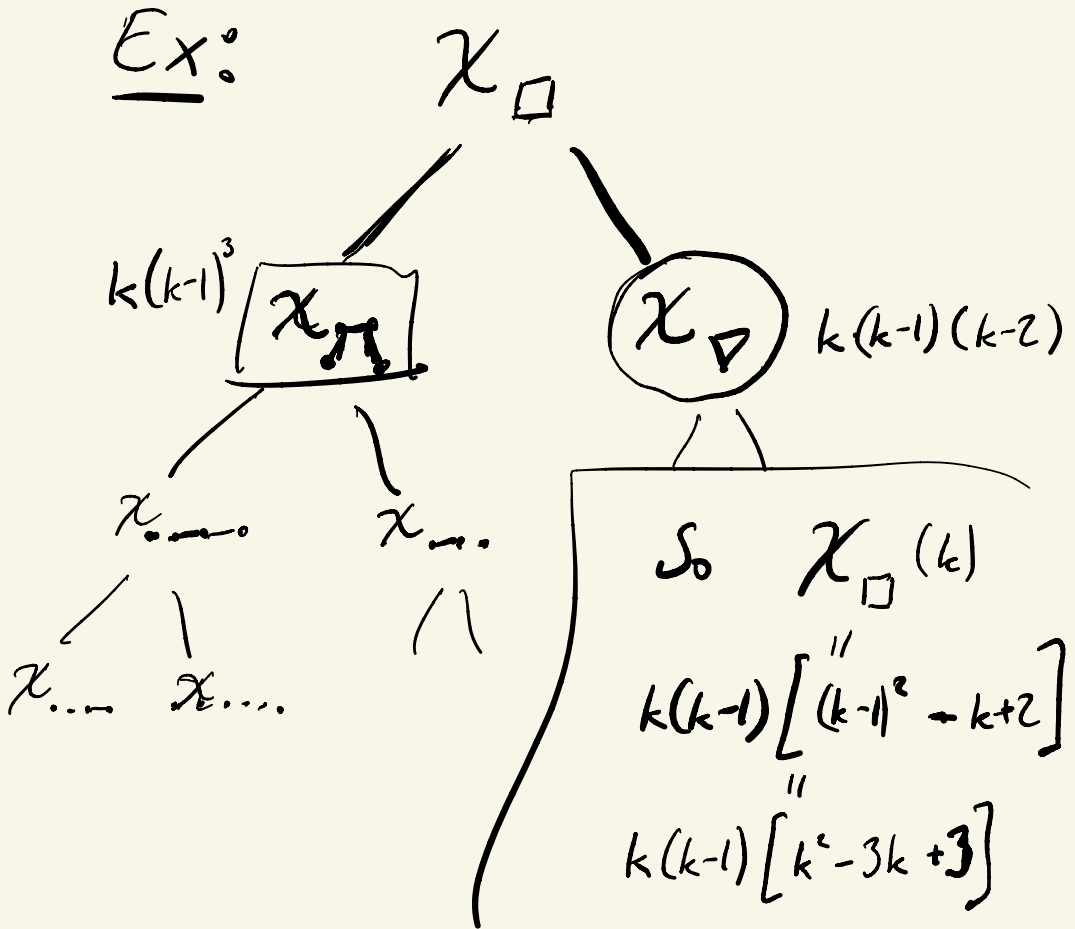
so by induction over # edges

- χ_G polynomial of degree $n = |G|$.
- leading coefficient 1.

- integer coefficients
- alternating coefficients:

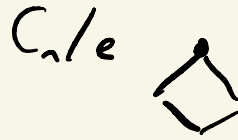
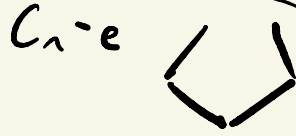
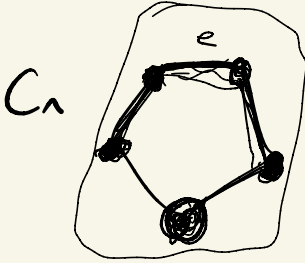
$$\chi_G(k) = k^n - a_1 k^{n-1} + a_2 k^{n-2} - \dots$$

where $a_i \in \mathbb{N}$
 $a_n = 0$



Ex $\chi_{C_n}(k) = \chi_{P_{n-1}}(k) - \chi_{C_{n-1}}(k)$

$= k(k-1)^{n-1} - \chi_{C_{n-1}}(k)$



Base case

$\chi_{C_3} = \chi_{K_3}$

$= k(k-1)(k-2)$



$\chi_{C_2}(k)$ $k(k-1)$

$\chi_{C_1}(k)$ k

Induction \Rightarrow

$\chi_{C_n}(k) = (k-1)^n + (k-1)(-1)^n$

$$\chi_G(k) = \text{\#k-colourings of } G \quad \text{if } k \in \mathbb{N}.$$

has n roots (with multiplicity) in \mathbb{C} .

$$\chi_G(0) = \chi_G(1) = \dots = \chi_G(\underbrace{\chi(G)-1}_{\text{chrom number}}) = 0$$

What are the other $n - \chi(G)$ roots?

Answer: $A = \{\text{algebraic numbers } i_n\}$

There is a real number r s.t. the set of possible roots of chromatic polynomials of graphs is

$$A \subseteq (-\infty, r) \cup \{0, 1\}$$

$\varphi =$
golden ratio.

$$\mathbb{N} \left(\frac{32}{27}, \varphi \right)$$

$$\bar{\chi}_G(k) \quad \chi_G(k) = k^n - a_1 k^{n-1} + a_2 k^{n-2} - \dots + a_{n-1} k$$

$$\therefore (-1)^n \bar{\chi}_G(-k) = k^n + a_1 k^{n-1} + a_2 k^{n-2} + \dots + a_{n-1} k$$

\forall
 $\chi_G(k)$

\uparrow
 \mathbb{N}

if $k \in \mathbb{N}$

Thm: $\bar{\chi}_G(k) = \# (\mathcal{P}, \gamma)$ where

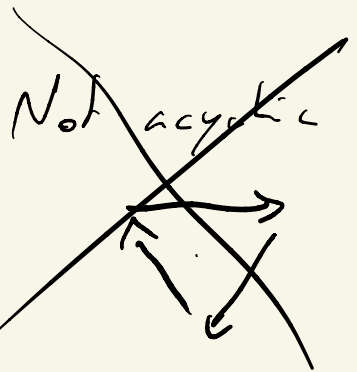
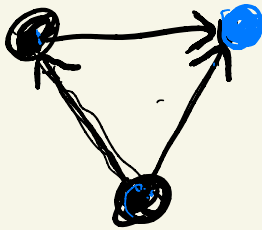
- \mathcal{P} , orientation of G acyclic
- $\gamma: V(G) \rightarrow \{1, \dots, k\}$ such that

$$u \xrightarrow{\mathcal{P}} v \Rightarrow \gamma(u) \geq \gamma(v)$$

" \mathcal{P} and γ are weakly compatible"

" \mathcal{P} and γ strongly compatible" if γ is also a proper graph colouring.

Acyclic orientation



Thm:

$(-1)^n \chi_G(-k)$ is the number of pairs (\mathcal{S}, γ) where \mathcal{S} acyclic orientation of G "weakly compatible" s.t.

- $\gamma: V(G) \rightarrow \{1, \dots, k\}$
- $u \xrightarrow{\mathcal{S}} v \Rightarrow \gamma(u) \leq \gamma(v).$

Pf: $\bar{\chi}_G = (-1)^n \chi_G(-k)$

$\Psi_G =$ numbers of weakly compatible pairs (\mathcal{S}, γ) .

Recall $\chi_G(x) = \chi_{G-e}(x) + \chi_{G|e}(x)$ if e ordinary edge

$\chi_G^{(x)} = k^n$ if G has no edges

$\chi_G^{(x)} = 0$ if G has loop

So $\bar{\chi}_G(k) = k^n$ if G has no edges

$\bar{\chi}_G(k) = 0$ if G has loops

$\bar{\chi}_G(k) = \bar{\chi}_{G-e}(k) + \bar{\chi}_{G/e}(k)$ if e ordinary edge

Enough to show

no edges to orient, every map $V \rightarrow \{1, \dots, k\}$ ok $\Psi_{K_n}(k) = k^n$

$\Psi_G(k) = 0$ if G has loops

$\Psi_G(k) = \Psi_{G-e}(k) + \Psi_{G/e}(k)$ if e ordinary edge!

no acyclic orientation

(s, γ) -pairs with $\gamma(u) \neq \gamma(v)$

$= \Psi_{G-e}$

+
(s, γ) -pairs with $\gamma(u) = \gamma(v)$ $u \xrightarrow{s} v$

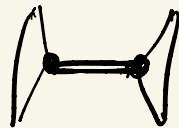
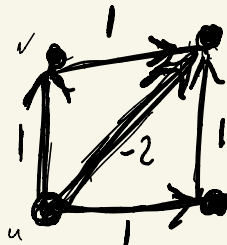
+
(s, γ) -pairs with $\gamma(u) = \gamma(v)$ $u \xleftarrow{s} v$ $= \Psi_{G/e}$



A Flow on ^{directed} graph G is a map \mathbb{Z}_n $f: E(G) \rightarrow \mathbb{Z}_n$ s.t

$$\sum_{e \in N_-(v)} f(e) = \sum_{e \in N_+(v)} f(e)$$

{ nowhere zero \mathbb{Z}_n flows }
on G

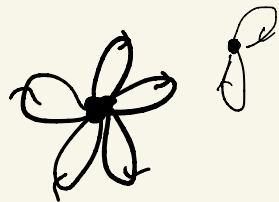


||
 $\varphi_G(\mathbb{Z}_n)$

Note: φ_G does not depend on the orientation of the graph.

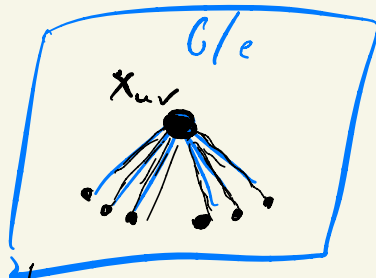
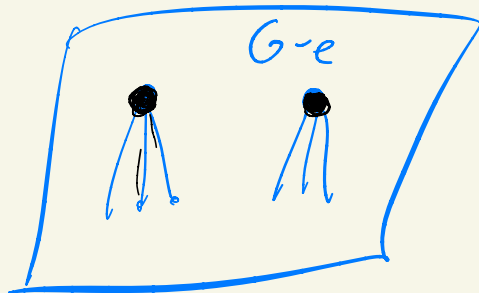
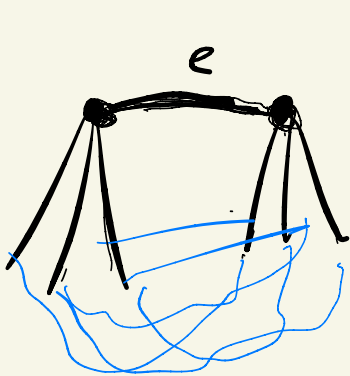
● If G has a bridge, then $\varphi_G(\mathbb{Z}_n) = 0$.

● If G has r loops then $\varphi_G(\mathbb{Z}_n) = (n-1)^r$



● If e ordinary (no loop, no bridge)

$$\varphi_G(\mathbb{Z}_n) = \varphi_{G-e}(\mathbb{Z}_n) + \varphi_{G/e}(\mathbb{Z}_n)$$



nowhere zero

A flow on G/e
gives a flow on G
and an almost-flow on $G-e$

(by assigning the
appropriate value
to $f(e)$).

This flow on G is nowhere zero

(\Rightarrow)

The "flow" on $G-e$ was
not conserved at u & v ,
i.e. the induced almost-flow
on $G-e$ is a flow.

so
$$\varphi_G(\mathbb{Z}_n) = \varphi_{G-e}(\mathbb{Z}_n) + \varphi_{G/e}(\mathbb{Z}_n).$$



So we get a graph polynomial

$$\varphi_G(n) := \varphi_G(\mathbb{Z}_n) \quad \text{in } n.$$

of degree $\chi(G) = |E| - |V| + |K|$

↑
set of
connected
components.

leading term $n^{\chi(G)}$

alternating integer
coefficients

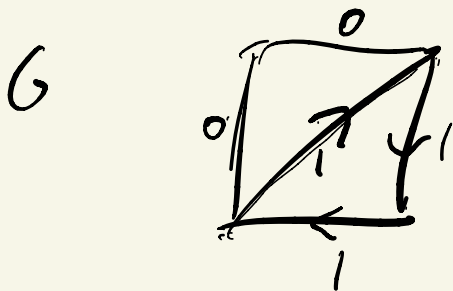
$$\varphi_G = \varphi_{G/e} - \varphi_{G-e}$$

So again

$$(-1)^{\chi(G)} \varphi_G(-n) \geq \varphi_G(n) \quad \text{if } n \in \mathbb{N}.$$

What does $(-1)^{\chi(G)} \varphi_G(-n)$ count?

Thm: $(-1)^{\chi(G)} \varphi_G(-n)$ is the number of
pairs (f, \mathcal{O}) where $f: E(G) \rightarrow \mathbb{Z}_n$
flow and \mathcal{O} is a
totally cyclic orientation of $G/\text{supp } f$



A totally cyclic orientation is one where every edge is included in a directed cycle.

In particular, $|\Psi_G(-1)|$ counts the number of totally cyclic orientations of G .

$|\chi_G(-1)|$ counts the number of acyclic orientations of G .

Tutte polynomial

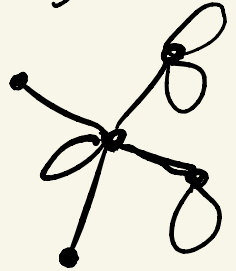
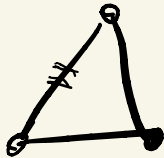
Def: $T_G(x,y)$ is

- $x^b y^l$ if G has b bridges and l loops and no other edges.

$$T_{G/e}(x,y) + T_{G/e}(x,y)$$

if e is an ordinary edge

Ex: $G = K_3$.



$$\begin{aligned} T_{\Delta} &= T_{\text{---}} + T_{\text{---}} \\ &= x^2 + T_{\text{---}} + T_{\text{---}} \\ &= x^2 + y + x \end{aligned}$$

Thm: $T_G(x,y) = \sum_{A \subseteq E} (x-1)^{r(E)-r(A)} (y-1)^{n(A)}$

where $r(A) = |V(G)| - \underset{\substack{\uparrow \\ \text{Connected} \\ \text{components}}}{K(G[A])}}{|E(A)|}$

"rank"

$n(A) = |A| - r(A)$

"nullity"

$r(\Delta) = 2$

Ex: $T_{\Delta} =$

$x^2 + x + y$

$$\left[\begin{array}{l} \Delta \quad (x-1)^0 (y-1)^1 \\ \text{---} \quad 3 (x-1)^0 \cdot (y-1)^0 \\ \text{---} \quad 3 (x-1)^1 (y-1)^0 \\ \text{---} \quad (x-1)^2 (y-1)^0 \end{array} \right] = \boxed{\begin{array}{l} (?) \\ y-1 \\ + \\ 3 \\ + \\ 3(x-1) \\ + \\ (x-1)^2 \end{array}}$$

Thm: $\varphi_G(x) = (-1)^{|E|-|V|+K(G)} T_G(0, 1-x)$

$$\chi_G(x) = (-x)^{-K(G)} T_G(x-1, 0)$$

$$T_G(2, 1) = \# \text{ forests}$$

$$T_G(1, 1) = \# \text{ spanning forests}$$

$$T_G(1, 2) = \# \text{ spanning subgraphs}$$

$$T_G(-2, 0) = \# \text{ acyclic orientation}$$

⋮