

Network Security: TLS 1.3 handshake

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Outline

- TLS 1.3 full handshake: 1-RTT
- Security properties, identity protection

Please refer to the Information Security course for an introduction to TLS

TLS 1.3 full handshake

<u>Client</u> <u>Server</u>

```
ClientHello
+ key share*
+ signature algorithms*
                                  1. Parameter
                                                  2. DHE or ECDHE
+ supported groups*
                                                  key exchange
                                  negotiation
+ server name*
+ certificate authorities*
                                                                     ServerHello
                                                                    + key share*
                                                         {EncryptedExtensions}
                                                         {CertificateRequest*}
                                                                  {Certificate*}
                                          3. Server
       4. Client
                                                           {CertificateVerify*}
                                          authentication
       authentication
                                                                       {Finished}
       (typically omitted)
                                                             [ApplicationData*]
{Certificate*}
                          5. Key
{CertificateVerify*}
                          confirmation
{Finished}
[Application data]
                                                             [Application data]
                            6. Protected session data
```

TLS 1.3 full handshake

```
1. C \rightarrow S:
                   N<sub>c</sub>, supported_versions, supported_groups, signature_algorithms,
                   cipher_suites, server_name, certificate_authorities, gx
2. S \rightarrow C: N_s, version, cipher_suite, g^y
                   EncryptedExtensions
                   Cert<sub>s</sub>, Sign<sub>s</sub>(TH)
                                                                       encrypted with K<sub>shts</sub>
                   HMAC_{Kfks}(TH)
3. C \rightarrow S: Cert<sub>c</sub>, Sign<sub>c</sub>(TH)
                                                                       encrypted with K<sub>chts</sub>
                   HMAC_{Kfkc}(TH)
```

 N_c , N_s = client and server random

Cert_c, Cert_s = certificate chain

TH = transcript hash i.e. hash of all previous messagas

Exchange keys K_{chts}, K_{shts}, K_{fkc}, K_{fks} session keys K_{cats}, K_{sats} derived from g^{xy} and TH

TLS 1.3 algorithms

- Small number of modern cipher suites
- AEAD ciphers: encryption and authentication always together
- Perfect forward secrecy required
 - Only ephemeral key exchanges: DHE or ECDHE
 - Old RSA handshake is not supported

1-RTT handshake

Client Server Client does not know which ClientHello groups the server supports but + key share* + signature algorithms* makes a guess + supported groups* + server name* ServerHello + certificate authorities* + key_share* {EncryptedExtensions} CertificateRequest*} {Certificate*} CertificateVerify*} {Finished} ([ApplicationData*]) {Certificate*} {CertificateVeri Finished}

[Application data]

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[Application data]

1-RTT handshake

- TLS 1.3 handshake causes only one round-trip delay
 - Client can send HTTP request (application data) right after client Finished
 - TLS 1.2 and most other key-exchange protocols require two RTT
 - Important for page load times in web browsing
- However, TCP + TLS 1.3 together cause 2-RTT latency
 - QUIC avoids this because it runs over UDP
- Sometimes TLS 1.3 hanshake takes two RTT:
 - If server does not support the group of key_share in ClientHello, server sends HelloRetryRequest to ask for a different curve
 - DTLS server under DoS attack can send a Cookie in HelloRetryRequest

Key derivation

one or both, as available

Inputs to key derivation:

- 1. PSK (external PSK or resumption PSK)
- 2. DHE/ECDHE secret
- 3. Transcript of handshake messages, up to the point where the key is derived

Keys:

- client_early_traffic_secret → used to derive AEAD keys for early data in 0-RTT (...)
- client/server_handshake_traffic_secret → used to derive AEAD keys for handshake messages {...} and Finished HMAC keys
- client/server_application_traffic_secret_N → used to derive AEAD encryption keys for post-handshake application data and messages [...]
- resumption_master_secret and ticket_nonce → derive resumption PSK
- exporter_master_secret → used to create keys for the application layer

Post-handshake client authentication

- Server can request client authentication any time, either during or after the TLS handshake
- Post-handshake client authentication allows time for user action, such as inserting a smartcard
 - Application can give user more access rights after the authentication

References

- TLS 1.3, <u>RFC 8446</u>
- The New Illustrated TLS Connection, https://tls13.ulfheim.net/

Exercises

- Use a network sniffer (e.g. Netmon, Ethereal) to look at TLS handshakes. Can you spot a full handshake and session resumption? Can you see the plaintext SNI?
- Compare TLS 1.3 and TLS 1.2 handshakes in network trace: Can you see the difference is round-trips, identity protection?
- How would you modify the TLS 1.3 handshake to improve identity protection? Learn about PEAP. How does PEAP protect the client identity?
- Consider removing message fields from the handshake. How does each message field contribute to security?
- Why have the supported and mandatory-to-implement cipher suites in TLS changed over time?
- Why did most web servers for a long time prefer the RSA handshake?
- One reason why the RSA handshake it is no longer supported in TLS 1.3 is that it does not provide PFS. Is it possible to implement PFS without Diffie-Hellman?
- Finds applications that could benefit significantly from the 0-RTT handshake. Is there any cost to deploying it?
- What problems arise if you want to set up multiple secure (HTTPS) web sites behind a NAT or on virtual servers that share one IP address? How to TLS 1.3 and TLS 1.2 solve this issue?
- If an online service (e.g. webmail) uses TLS with server-only authentication to protect passwords, is the system vulnerable to offline password cracking?