CS-E4075 - Special Course in Machine Learning, Data Science and Artificial Intelligence D: Signed graphs: spectral theory and applications

# Spectral clustering 

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Algorithms for $k$-means will do well on these data.

But how about this?


The results:



Let's talk about spectral clustering.



## Derivation

Let us try to cluster the graph on the right. There are two "obvious" clusters, but can we use a clustering algorithm to discover them?
We will try to represent the graph vertices in a way that is suitable for an algorithm such as $k$-means.


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We will try to represent the graph vertices in a way that is suitable for an algorithm such as $k$-means.

## We will consider the adjacency matrix:

$$
W=\left(\begin{array}{llllllllll}
0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
1 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\
1 & 1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 \\
\hline
\end{array}\right)
$$


$w_{i j}$ is the weight of the edge connecting $v_{i}$ and $v_{j}$.

One way to make sure we can use $k$-means is by:

- being able to compute distances between vertices: $d\left(v_{i}, v_{j}\right)$,
- being able to compute the mean of a cluster of vertices: $\mu_{j}$.

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- being able to compute the mean of a cluster of vertices: $\mu_{j}$.

One way to accomplish this is to assign a real number $y_{i}$ to each vertex $v_{i}$, so that

- $d\left(v_{i}, v_{j}\right)=\left|y_{i}-y_{j}\right|$ and
- $\mu_{j}=\frac{1}{\left|C_{j}\right|} \sum_{i \in C_{j}} y_{j}$.

Thus, our goal is to find a mapping of each vertex $v_{i} \mapsto y_{i} \in \mathbb{R}$ so that (intuitively) similar vertices are in the same cluster and different vertices are in different clusters.

In order to ensure that connected vertices are close, we can try to choose the $y_{i}$ 's so that the following is small:

$$
\operatorname{cost}(y)=\sum_{i=1}^{n} \sum_{j=1}^{n} w_{i j}\left(y_{i}-y_{j}\right)^{2}
$$



$$
\operatorname{cost}(y)=\sum_{i=1}^{n} \sum_{j=1}^{n} w_{i j}\left(y_{i}-y_{j}\right)^{2}=2 \sum_{i=1}^{n} y_{i}^{2} \sum_{j=1}^{n} w_{i j}-2 \sum_{i=1}^{n} \sum_{j=1}^{n} w_{i j} y_{i} y_{j}
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Note: $\sum_{j=1}^{n} w_{i j}=d_{i}$ is the degree of vertex $v_{i}$, and $D_{i i}=\sum_{j=1}^{n} w_{i j}$.

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Note: $\sum_{j=1}^{n} w_{i j}=d_{i}$ is the degree of vertex $v_{i}$, and $D_{i i}=\sum_{j=1}^{n} w_{i j}$.
Let $y=\left(y_{1}, \ldots, y_{n}\right)^{T}$. Note that $\operatorname{cost}(y)=2 y^{T} D y-2 y^{T} W y=2 y^{\top} L y$.

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- Choosing $y=\left(y_{1}, \ldots, y_{n}\right)^{T}$ to minimize $\sum_{i=1}^{n} \sum_{j=1}^{n} w_{i j}\left(y_{i}-y_{j}\right)^{2}$ seems like a good idea.

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- We want to cluster the vertices of a graph using $k$-means.
- We assign a real number $y_{i}$ to every vertex so that we can compute distances and means.
- Choosing $y=\left(y_{1}, \ldots, y_{n}\right)^{T}$ to minimize $\sum_{i=1}^{n} \sum_{j=1}^{n} w_{i j}\left(y_{i}-y_{j}\right)^{2}$ seems like a good idea.
- We define $L=D-W$ and show that minimizing $y^{\top} L y$ is equivalent.


## Some properties of the Laplacian

1. $L$ is symmetric and positive semidefinite (all eigenvalues are real and $\geq 0$ ).

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Remember we want to minimize $y^{\top} L y$.

- By property 1 , the minimum is at least $0 . y=(0,0 \ldots, 0)^{T}$ is a trivial solution, so we impose the constraint $y^{\top} y=1$.


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Remember we want to minimize $y^{\top} L y$.

- By property 1 , the minimum is at least $0 . y=(0,0 \ldots, 0)^{T}$ is a trivial solution, so we impose the constraint $y^{\top} y=1$.
- If the graph is connected, the vector $\mathbf{1}=(1,1, \ldots, 1)^{T}$ is a solution with $\mathbf{1}^{T} L \mathbf{1}=0$. We impose the constraint $y^{\top} \mathbf{1}=0$.


## Objective

$$
\begin{aligned}
\min _{y} & y^{\top} L y \\
\text { subject to } & y^{T} y=1 \\
& y^{\top} \mathbf{1}=0
\end{aligned}
$$

Since the vector $1=(1,1, \ldots, 1)^{T}$ is an eigenvector corresponding to the smallest eigenvalue, the above is solved (in a connected graph) by the eigenvector corresponding to the second smallest eigenvalue.

Spectral clustering protoalgorithm for 2-way connected graph partitioning. Input: Graph $G=(V, E)$ with adjacency matrix $W$.

1. Compute the Laplacian $L=D-W$.
2. Compute $y$, the eigenvector of $L$ corresponding to the second smallest eigenvalue.
3. Run $k$-means treating the entries of $y$ as one-dimensional data points.

Let's try to cluster our graph.


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$$
W=\left(\begin{array}{llllllllll}
0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
1 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\
1 & 1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0
\end{array}\right)
$$

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W=\left(\begin{array}{cccccccccc}
0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
1 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\
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0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0
\end{array}\right), L=\left(\begin{array}{cccccccccccccccc}
4 & -1 & -1 & -1 & -1 & 0 & 0 & 0 & 0 \\
-1 & 4 & -1 & -1 & -1 & 0 & 0 & 0 & 0 & 0 \\
-1 & -1 & 4 & -1 & -1 & 0 & 0 & 0 & 0 & 0 \\
-1 & -1 & -1 & 4 & -1 & 0 & 0 & 0 & 0 & 0 \\
-1 & -1 & -1 & -1 & 5 & -1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & -1 & 5 & -1 & -1 & -1 & -1 \\
0 & 0 & 0 & 0 & 0 & -1 & 4 & -1 & -1 & -1 \\
0 & 0 & 0 & 0 & 0 & -1 & -1 & 4 & -1 & -1 \\
0 & 0 & 0 & 0 & 0 & -1 & -1 & -1 & 4 & -1 \\
0 & 0 & 0 & 0 & 0 & -1 & -1 & -1 & -1 & 4
\end{array}\right)
$$

## Let's try to cluster our graph.


$\boldsymbol{W}=\left(\begin{array}{cccccccccc}0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0\end{array}\right), L=\left(\begin{array}{ccccccccccccc}4 & -1 & -1 & -1 & -1 & 0 & 0 & 0 & 0 & 0 \\ -1 & 4 & -1 & -1 & -1 & 0 & 0 & 0 & 0 & 0 \\ -1 & -1 & 4 & -1 & -1 & 0 & 0 & 0 & 0 & 0 \\ -1 & -1 & -1 & 4 & -1 & 0 & 0 & 0 & 0 & 0 \\ -1 & -1 & -1 & -1 & 5 & -1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & -1 & 5 & -1 & -1 & -1 & -1 \\ 0 & 0 & 0 & 0 & 0 & -1 & 4 & -1 & -1 & -1 \\ 0 & 0 & 0 & 0 & 0 & -1 & -1 & 4 & -1 & -1 \\ 0 & 0 & 0 & 0 & 0 & -1 & -1 & -1 & 4 & -1 \\ 0 & 0 & 0 & 0 & 0 & -1 & -1 & -1 & -1 & 4\end{array}\right)$

Eigenvalues: ( $0, \sim 0.3,5,5,5,5,5,5,5, \sim 6.7$ ).
Second smallest eigenvector: $(0.33,0.33,0.33,0.33,0.23,-0.23,-0.33,-0.33,-0.33,-0.33)^{T}$.

Let's try to cluster our graph.

$W=\left(\begin{array}{cccccccccc}0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0\end{array}\right), L=\left(\begin{array}{ccccccccccccc}4 & -1 & -1 & -1 & -1 & 0 & 0 & 0 & 0 & 0 \\ -1 & 4 & -1 & -1 & -1 & 0 & 0 & 0 & 0 & 0 \\ -1 & -1 & 4 & -1 & -1 & 0 & 0 & 0 & 0 & 0 \\ -1 & -1 & -1 & 4 & -1 & 0 & 0 & 0 & 0 & 0 \\ -1 & -1 & -1 & -1 & 5 & -1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & -1 & 5 & -1 & -1 & -1 & -1 \\ 0 & 0 & 0 & 0 & 0 & -1 & 4 & -1 & -1 & -1 \\ 0 & 0 & 0 & 0 & 0 & -1 & -1 & 4 & -1 & -1 \\ 0 & 0 & 0 & 0 & 0 & -1 & -1 & -1 & 4 & -1 \\ 0 & 0 & 0 & 0 & 0 & -1 & -1 & -1 & -1 & 4\end{array}\right)$

Eigenvalues: $(0, \sim 0.3,5,5,5,5,5,5,5, \sim 6.7)$.
Second smallest eigenvector: $(0.33,0.33,0.33,0.33,0.23,-0.23,-0.33,-0.33,-0.33,-0.33)^{T}$.

$\longleftarrow$ We can cluster this with $k$-means.

## Another example

Fully connected weighted graph.
Eigenvector:
$(0.32,0.34,0.28,0.34,0.29,-0.32,-0.27,-0.31,-0.36,-0.31)^{T}$



## Optimizing graph cuts

## Clustering with cuts

Given a graph $G=(V, E)$, and a vertex subset $S \subseteq V$, with adjacency matrix $A$,

$$
\operatorname{cut}(S, \bar{S})=E(S, \bar{S})=\sum_{i \in S, j \in \bar{S}} A_{i j} .
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Given a graph $G=(V, E)$, and a vertex subset $S \subseteq V$, with adjacency matrix $A$,

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## Optimizing RatioCut

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\text { Define a vector } x \text { as }
$$

$\operatorname{RatioCut}(S, \bar{S})=\left(\frac{1}{|S|}+\frac{1}{|\bar{S}|}\right) \operatorname{cut}(S, \bar{S})$.

$$
x_{i}= \begin{cases}\sqrt{|\bar{S}| /|S|} & \text { if } v_{i} \in S \\ -\sqrt{|S| /|\bar{S}|} & \text { if } v_{i} \in \bar{S}\end{cases}
$$

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$$

We have

- $x^{T} L x=|V| \cdot \operatorname{RatioCut}(S)$,
- $x^{\top} 1=0$,
- $\|x\|_{2}=\sqrt{|V|}$.


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Objective:

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$$
\begin{aligned}
\min _{x} & x^{\top} L x \\
\text { subject to } & x^{\top} 1=0 \\
& \|x\|_{2}=\sqrt{|V|} \\
& x_{i} \text { as defined above. }
\end{aligned}
$$

## Clustering with cuts

## Optimizing RatioCut

$k$ clusters:
$\operatorname{RatioCut}\left(S_{1}, \ldots, S_{k}\right)=\sum_{i=1}^{k} \frac{\operatorname{cut}\left(S_{i}, \bar{S}_{i}\right)}{\left|S_{i}\right|}$.

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x_{i j}= \begin{cases}\frac{1}{\sqrt{\left|S_{i}\right|}} & \text { if } v_{j} \in S_{i} \\ 0 & \text { if } v_{j} \in \bar{S}_{i}\end{cases}
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We have

- $x_{i}^{T} L x_{i}=\frac{\operatorname{cut}\left(S_{i}, \bar{S}_{i}\right)}{\left|S_{i}\right|}$,
$-\sum_{i} x_{i}^{\top} L x_{i}=\sum_{i} \frac{\operatorname{cut}\left(S_{i}, \bar{S}_{i}\right)}{\left|S_{i}\right|}$,
- $x_{i}^{\top} x_{j}=0, i \neq j$,
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$$

Objective:

$$
\begin{aligned}
\min _{X} & \operatorname{Tr}\left(X^{T} L X\right) \\
\text { subject to } & X^{T} X=I \\
& X_{i j}=x_{i j} \text { as defined above. }
\end{aligned}
$$

## Clustering with cuts

## Optimizing NCut

$\operatorname{NCut}(S, \bar{S})=\left(\frac{1}{\operatorname{vol}(S)}+\frac{1}{\operatorname{vol}(\bar{S})}\right) \operatorname{cut}(S, \bar{S})$.
$\operatorname{vol}(S)=\sum_{v \in S} d(v)$.

## Clustering with cuts

## Optimizing NCut

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-\sqrt{\operatorname{vol}(S) / \operatorname{vol}(\bar{S})} & \text { if } v_{i} \in \bar{S} .\end{cases} \\
& \operatorname{vol}(S)=\sum_{v \in S} d(v) .
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We have

- $x^{\top} L x=\operatorname{vol}(V) \cdot \operatorname{NCut}(S)$,
- $(D x)^{T} 1=0$,
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- $x^{\top} L x=\operatorname{vol}(V) \cdot \operatorname{NCut}(S)$,
- $(D x)^{T} 1=0$,
- $x^{T} D x=\operatorname{vol}(V)$.

Note: $D^{-1 / 2} L D^{-1 / 2}$ is known as the normalized Laplacian.

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$$

We have

- $x_{i}^{T} L x_{i}=\frac{\operatorname{cut}\left(S_{i}, \bar{S}_{i}\right)}{\operatorname{vol}\left(S_{i}\right)}$,
$-\sum_{i} x_{i}^{T} L x_{i}=\sum_{i} \frac{\operatorname{cut}\left(S_{i}, \bar{S}_{i}\right)}{\operatorname{vol}\left(S_{i}\right)}$,
- $x_{i}^{\top} x_{j}=0, i \neq j$,
- $x_{i}^{T} D x_{i}=1$.

$$
x_{i j}= \begin{cases}\frac{1}{\sqrt{\mathrm{vol}\left(S_{i}\right)}} & \text { if } v_{j} \in S_{i} \\ 0 & \text { if } v_{j} \in \bar{S}_{i}\end{cases}
$$

Objective:

$$
\min _{X} \operatorname{Tr}\left(X^{T} D^{-1 / 2} L D^{-1 / 2} X\right)
$$

subject to $\quad X^{\top} X=I$,
$X_{i j}=x_{i j}$ as defined above.

Unnormalized spectral clustering .
Input: Graph $G=(V, E)$ with adjacency matrix $W$, number of clusters $k$.

1. Compute the Laplacian $L=D-W$.
2. Compute the eigenvectors $v_{1}, \ldots, v_{k}$ of $L$ corresponding to the $k$ smallest eigenvalues.
3. Consider a matrix $X$ whose columns are $v_{1}, \ldots, v_{k}$.
4. Run $k$-means on $X$.

Normalized variants:

## Spectral clustering using normalized Laplacian.

Input: Graph $G=(V, E)$ with adjacency matrix $W$, number of clusters $k$.

1. Compute the normalized Laplacian $L_{s y m}=D^{-1 / 2}(D-W) D^{-1 / 2}$.
2. Compute the eigenvectors $v_{1}, \ldots, v_{k}$ of $L_{\text {sym }}$ corresponding to the $k$ smallest eigenvalues.
3. Consider a matrix $X$ whose columns are $v_{1}, \ldots, v_{k}$. Normalize the rows of $X$ to unit norm.
4. Run $k$-means on $X$.

Note: $L_{\text {sym }}=D^{-1 / 2} L D^{-1 / 2}=I-D^{-1 / 2} W D^{-1 / 2}$.
Further reading:

- Von Luxburg, Ulrike. "A tutorial on spectral clustering." Statistics and computing 17.4 (2007)
- Ng, Jordan, and Weiss. "On spectral clustering: Analysis and an algorithm." NIPS 2002.

Normalized variants:

## Spectral clustering using random walk Laplacian.

Input: Graph $G=(V, E)$ with adjacency matrix $W$, number of clusters $k$.

1. Compute the random walk Laplacian $L_{r w}=D^{-1}(D-W)$.
2. Compute the right eigenvectors $v_{1}, \ldots, v_{k}$ of $L_{r w}$ corresponding to the $k$ smallest eigenvalues.
3. Consider a matrix $X$ whose columns are $v_{1}, \ldots, v_{k}$. Normalize the rows of $X$ to unit norm.
4. Run $k$-means on $X$.

Note: $L_{r w}=D^{-1} L=I-D^{-1} W$. This is related to the random walk matrix.

> Further reading:

- Von Luxburg, Ulrike. "A tutorial on spectral clustering." Statistics and computing 17.4 (2007)
- Dhillon, Inderjit S., Yuqiang Guan, and Brian Kulis. "Kernel k-means: spectral clustering and normalized cuts." KDD 2004.

Recommended read: "A tutorial on spectral clustering." by Ulrike Von Luxburg.
As a rule of thumb, use the random walk variant.

## Spectral clustering in practice

We have seen how to use spectral clustering on graphs.
But can we use it on any type of data? E.g. points in $\mathbb{R}^{d}$.


Example by Von Luxburg ${ }^{1}$.

We can create a graph based on these data.


[^0]
## Example by Von Luxburg ${ }^{1}$.

epsilon-graph: there is an edge between $x$ and $y$ if and only if $\|x-y\|_{2} \leq \epsilon$.
epsilon-graph, epsilon=0.3


[^1] 395-416.

Example by Von Luxburg ${ }^{1}$.
$k$-nearest-neighbours: there is an edge between $x$ and $y$ if and only if $x$ is one of the $k$ nearest neighbours of $y$ or $y$ is one of the $k$ nearest neighbours of $x$.


[^2]Example by Von Luxburg ${ }^{1}$.
mutual $k$-nearest-neighbours: there is an edge between $x$ and $y$ if and only if $x$ is one of the $k$ nearest neighbours of $y$ and $y$ is one of the $k$ nearest neighbours of $x$.


[^3]Alternatively, we can consider a fully-connected weighted graph, by using a similarity function.

We will use the Gaussian (or RBF) kernel, defined as follows:

$$
\begin{aligned}
\kappa: \mathbb{R}^{d} \times \mathbb{R}^{d} & \rightarrow \mathbb{R} \\
x, y & \mapsto \kappa(x, y)=\exp \left(\frac{-\|x-y\|_{2}^{2}}{2 \sigma^{2}}\right) .
\end{aligned}
$$

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Recall density for

$$
\mathcal{N}(\mu, \sigma): f(x)=\frac{1}{\sigma \sqrt{2 \pi}} \exp \left(\frac{-(x-\mu)^{2}}{2 \sigma^{2}}\right)
$$


$\kappa(x, y)=\exp \left(\frac{-\|x-y\|_{2}^{2}}{2 \sigma^{2}}\right)$
We set $x$ to be the red point and compute the value of $\kappa(x, y)$ for all $y$, with different values of $\sigma$. We plot each point with opacity equal to $\kappa(x, y)$.

$\kappa(x, y)=\exp \left(\frac{-\|x-y\|_{2}^{2}}{2 \sigma^{2}}\right)$
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- $\sigma=0.1$

$\kappa(x, y)=\exp \left(\frac{-\|x-y\|_{2}^{2}}{2 \sigma^{2}}\right)$
We set $x$ to be the red point and compute the value of $\kappa(x, y)$ for all $y$, with different values of $\sigma$. We plot each point with opacity equal to $\kappa(x, y)$.
- $\sigma=0.2$

$\kappa(x, y)=\exp \left(\frac{-\|x-y\|_{2}^{2}}{2 \sigma^{2}}\right)$
We set $x$ to be the red point and compute the value of $\kappa(x, y)$ for all $y$, with different values of $\sigma$. We plot each point with opacity equal to $\kappa(x, y)$.
- $\sigma=0.4$

$\kappa(x, y)=\exp \left(\frac{-\|x-y\|_{2}^{2}}{2 \sigma^{2}}\right)$
We set $x$ to be the red point and compute the value of $\kappa(x, y)$ for all $y$, with different values of $\sigma$. We plot each point with opacity equal to $\kappa(x, y)$.
- $\sigma=0.8$

$\kappa(x, y)=\exp \left(\frac{-\|x-y\|_{2}^{2}}{2 \sigma^{2}}\right)$
We set $x$ to be the red point and compute the value of $\kappa(x, y)$ for all $y$, with different values of $\sigma$. We plot each point with opacity equal to $\kappa(x, y)$.

- $\sigma=1.6$
$\kappa(x, y)=\exp \left(\frac{-\|x-y\|_{2}^{2}}{2 \sigma^{2}}\right)$
We set $x$ to be the red point and compute the value of $\kappa(x, y)$ for all $y$, with different values of $\sigma$. We plot each point with opacity equal to $\kappa(x, y)$.

- $\sigma=1.6$

Note: the matrix of a fully connected graph might be too large to store. Consider thresholding or using nearest neighbours.

## Practical considerations

Some recommendations:

- Scalability: the similarity (adjacency) matrix is of size $n \times n$. Thus, spectral clustering requires at least $O\left(n^{2}\right)$ computations just for the preliminary phase. The spectral decomposition requires $O\left(n^{3}\right)$ work in general. Always use sparse matrices if possible.

[^4]
## Practical considerations

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- Use a nearest neighbour graph to avoid storing the $n \times n$ matrix.

[^5]
## Practical considerations

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- Use a nearest neighbour graph to avoid storing the $n \times n$ matrix.
- Compute eigenvectors efficiently ${ }^{2}$

[^6]
## Practical considerations

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- Use a nearest neighbour graph to avoid storing the $n \times n$ matrix.
- Compute eigenvectors efficiently ${ }^{2}$
- Sample random subgraphs.

[^7]
## Practical considerations

Some recommendations:

- Scalability: the similarity (adjacency) matrix is of size $n \times n$. Thus, spectral clustering requires at least $O\left(n^{2}\right)$ computations just for the preliminary phase. The spectral decomposition requires $O\left(n^{3}\right)$ work in general. Always use sparse matrices if possible.
- Use a nearest neighbour graph to avoid storing the $n \times n$ matrix.
- Compute eigenvectors efficiently ${ }^{2}$
- Sample random subgraphs.
- Read more ${ }^{3}$

[^8]Take-aways from this lecture:

- Derivation of spectral clustering from first principles.
- Derivation of spectral clustering from cut objectives:
- RatioCut
- NCut
- Spectral clustering algorithms.
- Building a graph from vector data.
- Practical considerations.


[^0]:    ${ }^{1}$ Von Luxburg, Ulrike. "A tutorial on spectral clustering." Statistics and computing 17.4 (2007): 395-416.

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[^2]:    ${ }^{1}$ Von Luxburg, Ulrike. "A tutorial on spectral clustering." Statistics and computing 17.4 (2007): 395-416.

[^3]:    ${ }^{1}$ Von Luxburg, Ulrike. "A tutorial on spectral clustering." Statistics and computing 17.4 (2007): 395-416.

[^4]:    ${ }^{2}$ https://en.wikipedia.org/wiki/LOBPCG
    ${ }^{3}$ Yan, Donghui, Ling Huang, and Michael I. Jordan. "Fast approximate spectral clustering." KDD 2009.

[^5]:    ${ }^{2} h t t p s: / / e n . w i k i p e d i a . o r g /$ wiki/LOBPCG
    ${ }^{3}$ Yan, Donghui, Ling Huang, and Michael I. Jordan. "Fast approximate spectral clustering." KDD 2009.

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