



Aalto University
School of Science

REA method: Preference Programming in efficiency analysis

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Notation

K DMUs

M Inputs

N Outputs

$$x_k = (x_1, x_2, \dots, x_m)$$

Vector of inputs for DMU_k

$$y_k = (y_1, y_2, \dots, y_n)$$

Vector of outputs for DMU_k

$$v = (v_1, v_2, \dots, v_m)$$

Weight vector for the inputs

$$u = (u_1, u_2, \dots, u_n)$$

Weight vector for the outputs

CCR-DEA

Maximize the efficiency for each DMU by choosing the most favourable input and output weights.

The efficiency of DMU_k is calculated as so:

$$E_k(u, v) = \frac{\text{output}}{\text{input}} = \frac{u^T y_k}{v^T x_k} = \frac{\sum_{i=1}^n v_i y_{ik}}{\sum_{i=1}^m u_i x_{ik}}$$

Shortcomings of CCR-DEA

- **Other combinations of input/output weights are neglected**
 - Relevant preference information is lost
- **Method is sensitive to outliers**
 - An outlier DMU can change the efficient frontier drastically
- **Requires a large number of DMUs to form the efficient frontier accurately**

Salo and Punkka (2011)

Ratio-based efficiency analysis (REA)

- **Pairwise comparisons between DMUs using ALL feasible sets of weights**
- **Important concepts:**
 - Ranking intervals
 - Efficiency dominance
 - Efficiency bounds

Ranking intervals

$$E_k(u, v) = \frac{u^T y_k}{v^T x_k}$$

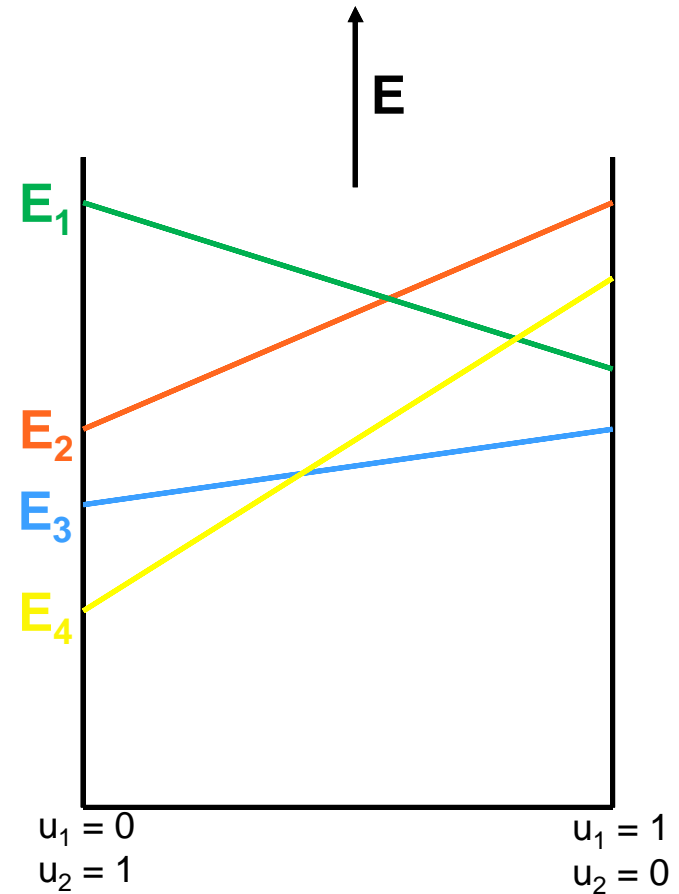
- **DMUs are ranked by their efficiency scores for all feasible weights**
- **What are the best and worst ranks a given DMU can get when the weights are varied?**
 - Can be solved with linear optimization

A1

Ranking intervals

- 4 DMUs, one input, two outputs
- Output weights u_1 u_2 are varied

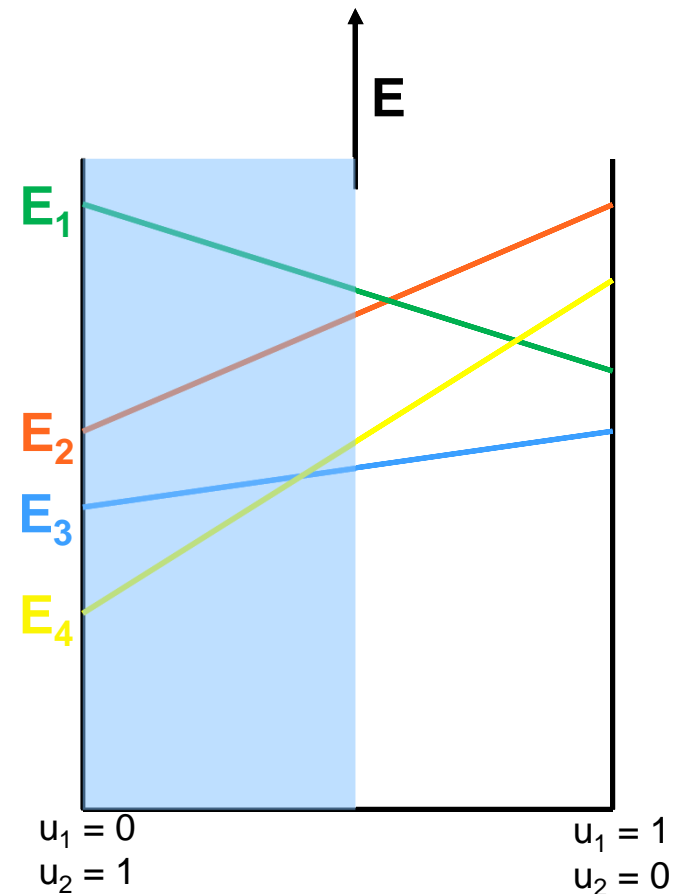
DMU	1	2	3	4
Max rank	1	1	3	2
Min rank	3	2	4	4
Interval	[1,3]	[1,2]	[3,4]	[2,4]



Ranking intervals

- Additional preference information narrows the intervals
- If $u_1 < u_2$ we are in the blue region

DMU	1	2	3	4
Max rank	1	2	3	3
Min rank	1	2	4	4
Interval	[1,1]	[2,2]	[3,4]	[3,4]



Ranking intervals

$r_k^{min}:$

$$\min_{x,y,z} 1 + \sum_{l \neq k} z_l$$

$$\sum_n u_n y_{nl} \leq \sum_m v_m x_{ml} + C z_l$$

$$l \in \{1, \dots, k\}, l \neq k$$

$$\sum_n u_n y_{nk} = \sum_m v_m x_{mk}$$

$$z_l \in \{0, 1\}, l \neq k$$

$$(u, v) \in (S_u, S_v)$$

$r_k^{max}:$

$$\max_{x,y,z} 1 + \sum_{l \neq k} z_l$$

$$\sum_n v_m x_{ml} \leq \sum_m u_n y_{nl} + C(1 - z_l)$$

$$l \in \{1, \dots, k\}, l \neq k$$

$$\sum_n u_n y_{nk} = \sum_m v_m x_{mk}$$

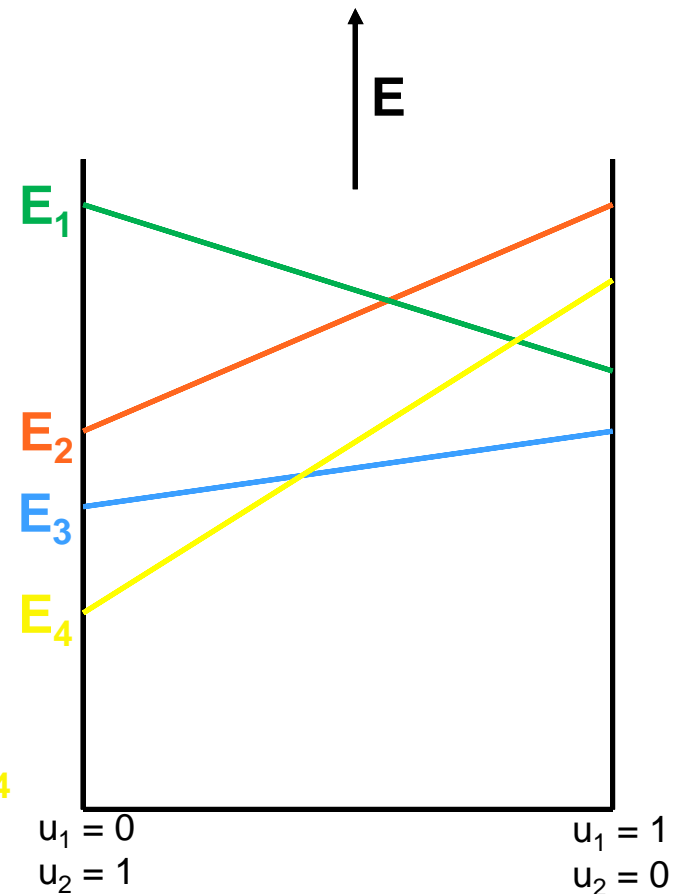
$$z_l \in \{0, 1\}, l \neq k$$

$$(u, v) \in (S_u, S_v)$$

Efficiency dominance

- DMU_k dominates DMU_l if
 1. $E_k(u, v) \geq E_l(u, v)$ for all feasible u and v
 2. $E_k(u, v) > E_l(u, v)$ for some feasible u and v

DMU₂ dominates **DMU₃** and **DMU₄**



Efficiency dominance

- The pairwise dominance relation is defined as so

$$D_{k,l}(u, v) = \frac{E_k(u, v)}{E_l(u, v)}$$

- If $\min D_{k,l}(u, v) \geq 1$ and $\max D_{k,l}(u, v) > 1$, DMU_k dominates DMU_l
- Can be solved with linear programming

Efficiency dominance

$$\begin{aligned}
 D_{k,l}(u, v) &= \frac{E_k(u, v)}{E_l(u, v)} \\
 &= \frac{\sum_n u_n y_{nk}}{\sum_m v_m x_{mk}} \bigg/ \frac{\sum_n u_n y_{nl}}{\sum_m v_m x_{ml}}
 \end{aligned}$$

Diagram illustrating the components of the efficiency dominance ratio $D_{k,l}(u, v)$:

- The numerator $\sum_n u_n y_{nk}$ is associated with the objective function: $\min_{u,v} \sum_n u_n y_{nk}$
- The denominator $\sum_m v_m x_{ml}$ is associated with the constraint: $\sum_n u_n y_{nl} = \sum_m v_m x_{ml}$
- The denominator $\sum_m v_m x_{mk}$ is associated with the constraint: $\sum_m v_m x_{mk} = 1$

If the solution is exactly 1, the same problem is maximized. If that solution is >1 , we have dominance.

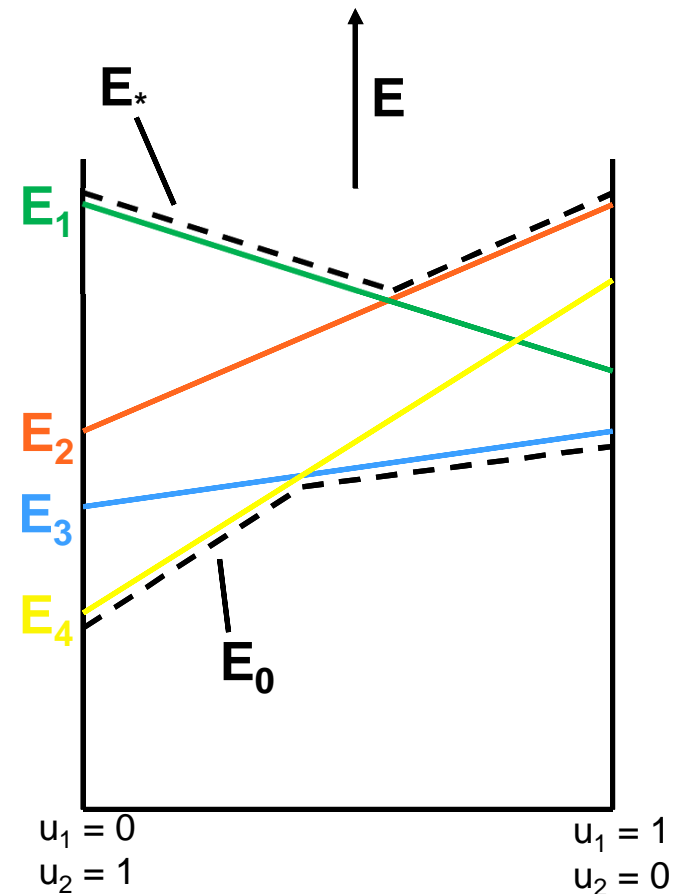
$$(u, v) \in (S_u, S_v)$$

Efficiency bounds

- How efficient can a given DMU be with respect to the most (E_*) and least (E_0) efficient DMUs?
- Another optimization problem over u, v .

$$E_2 \in [0.6, 1.0] E_*$$

$$E_2 \in [1.3, 1.8] E_0$$



Efficiency bounds

$$D_{k,\bar{L}}(u, v) = \frac{E_k(u, v)}{\max_{l \in L} E_l(u, v)}$$

$$D_{k,\underline{L}}(u, v) = \frac{E_k(u, v)}{\min_{l \in L} E_l(u, v)}$$

$$\underline{D}_{k,\bar{L}}:$$

$$\max_{u,v} \sum_n u_n y_{nk}$$

$$\sum_m v_m x_{mk} = 1, \quad (u, v) \in (S_u, S_v)$$

$$\sum_n u_n y_{nl} \leq \sum_m v_m x_{ml}, \quad l \in L$$

$$\bar{D}_{k,\underline{L}}:$$

$$\min_{u,v} \sum_n u_n y_{nk}$$

$$\sum_m v_m x_{mk} = 1, \quad (u, v) \in (S_u, S_v)$$

$$\sum_n u_n y_{nl} \leq \sum_m v_m x_{ml}, \quad l \in L$$

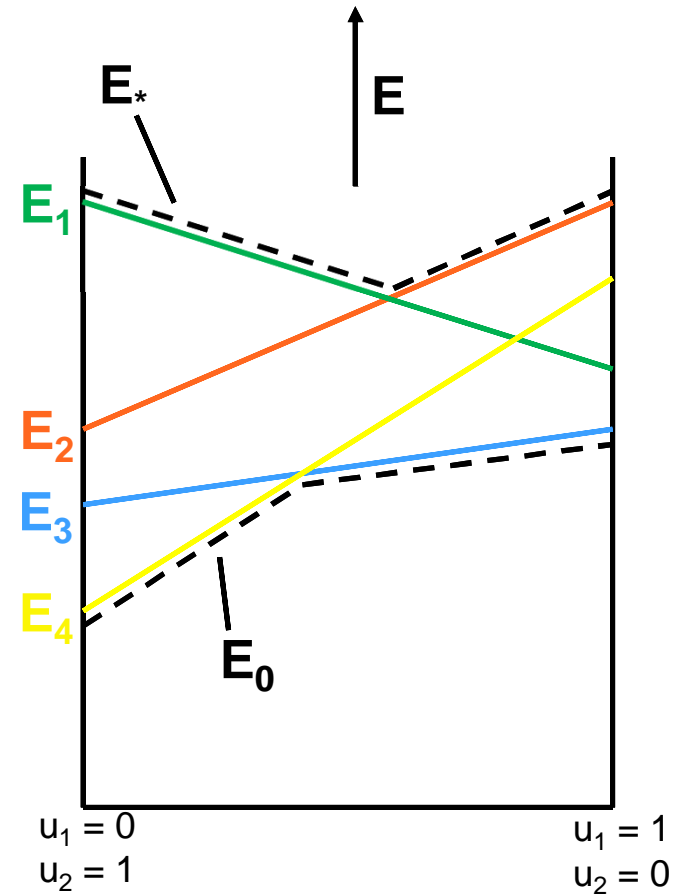
Efficiency bounds

Recall: $E_2 \in [0.6, 1.0]E_*$
 $E_2 \in [1.3, 1.8]E_0$

In this case

$$\bar{D}_{2,\bar{L}} = 1.0 \quad \underline{D}_{2,\bar{L}} = 0.6$$

$$\bar{D}_{2,\underline{L}} = 1.8 \quad \underline{D}_{2,\underline{L}} = 1.3$$



Applying REA to Scottish Health Boards

Schang et. al. (2016)

Application

- Health care quality of the 14 boards are compared
- Inputs and outputs are normalized to $[0,1]$ intervals.
- Difficulties in defining weights → Need for robust ranking methods, aka. REA

Inputs

- Resident population
- Number of occupied bed days (OBD)

Outputs

- Over 18 week wait from referral to treatment
- Longer than 4 hour wait in A&E
- Emergency admission
- Delayed discharge
- C.difficile infection
- MRSA/MSSA infection

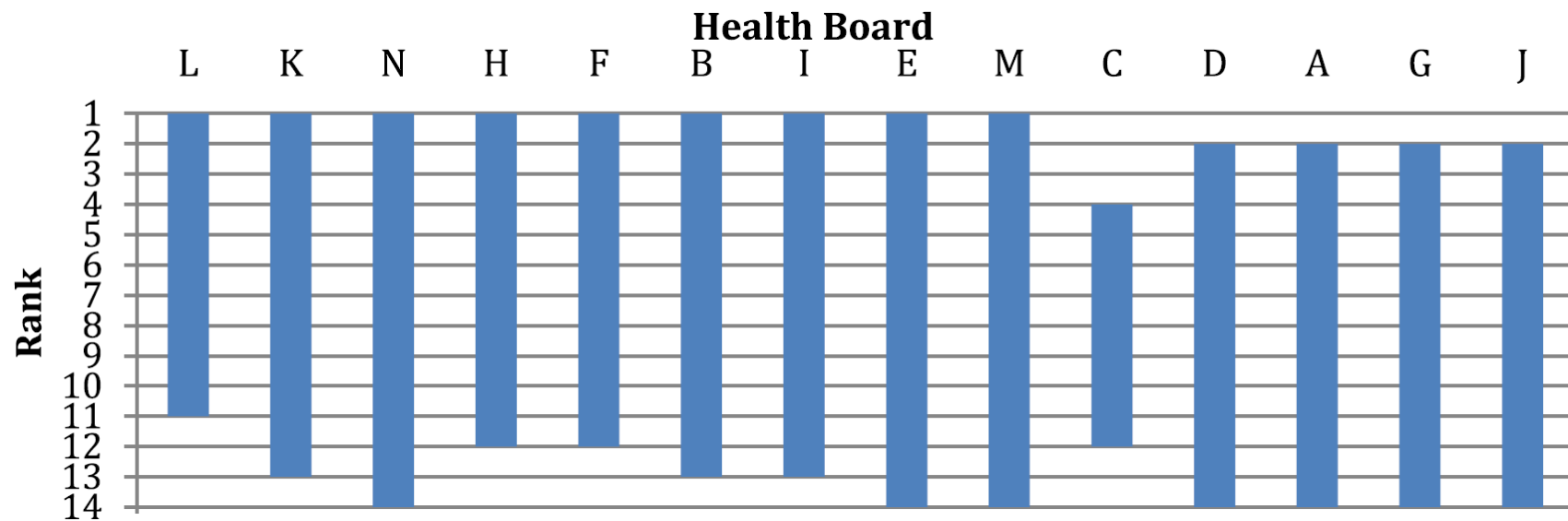
Application

		4-hour A&E waiting	Emergency admissions	MRSA/MSSA	
A	Ayrshire & Arran	8312	3646	23	
B	Borders	3267	3612	21	
C	Dumfries & Galloway	5987	3130	27	
D	Fife	4559	2725	35	
E	Forth Valley	8238	2513	26	
F	Grampian	3812	2239	25	
G	Greater Glasgow & Clyde	6956	3061	34	...
H	Highland	2199	2825	17	
I	Lanarkshire	8667	2671	24	
J	Lothian	9172	2495	30	
K	Orkney	1663	2661	9	
L	Shetland	730	2555	13	
M	Tayside	1119	2964	36	
N	Western Isles	1666	3320	4	

Application

- **Ranking intervals and dominances were solved with three different sets of weights. Input variables were unrestricted in each scenario.**
 - **Set 1: No weight restrictions**
 - **Set 2: Ordinal weights restrictions**
 - **Set 3: Ordinal and proportional weight restrictions**

No weighting – Ranking intervals



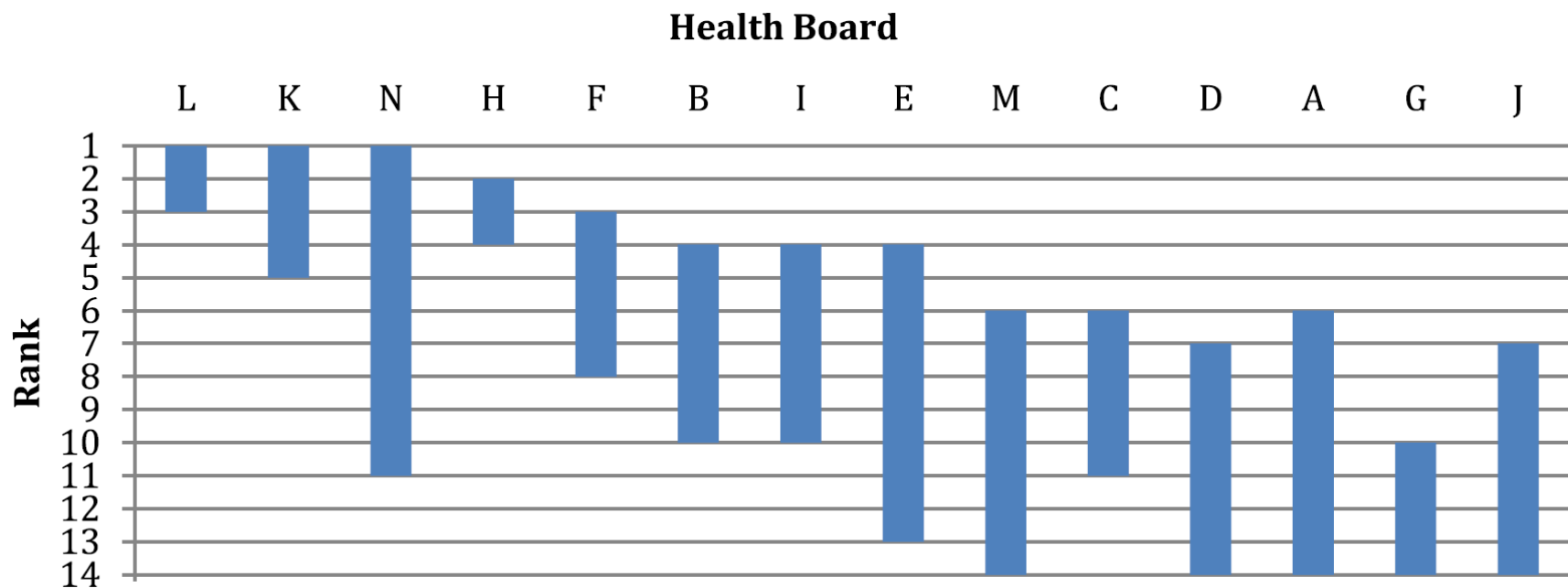
Schang et al. (2016)

Ordinal weighting

DMs were asked to rank the outputs from worst to least bad

- 1. an MRSA/MSSA infection**
- 2. an emergency admission**
- 3. a C.difficile infection**
- 4. longer than 18 week wait from referral to treatment**
- 5. longer than 4 hour wait in A&E**
- 6. a delayed discharge**

Ordinal weighting – Rank intervals



Schang et al. (2016)

Proportional weighting

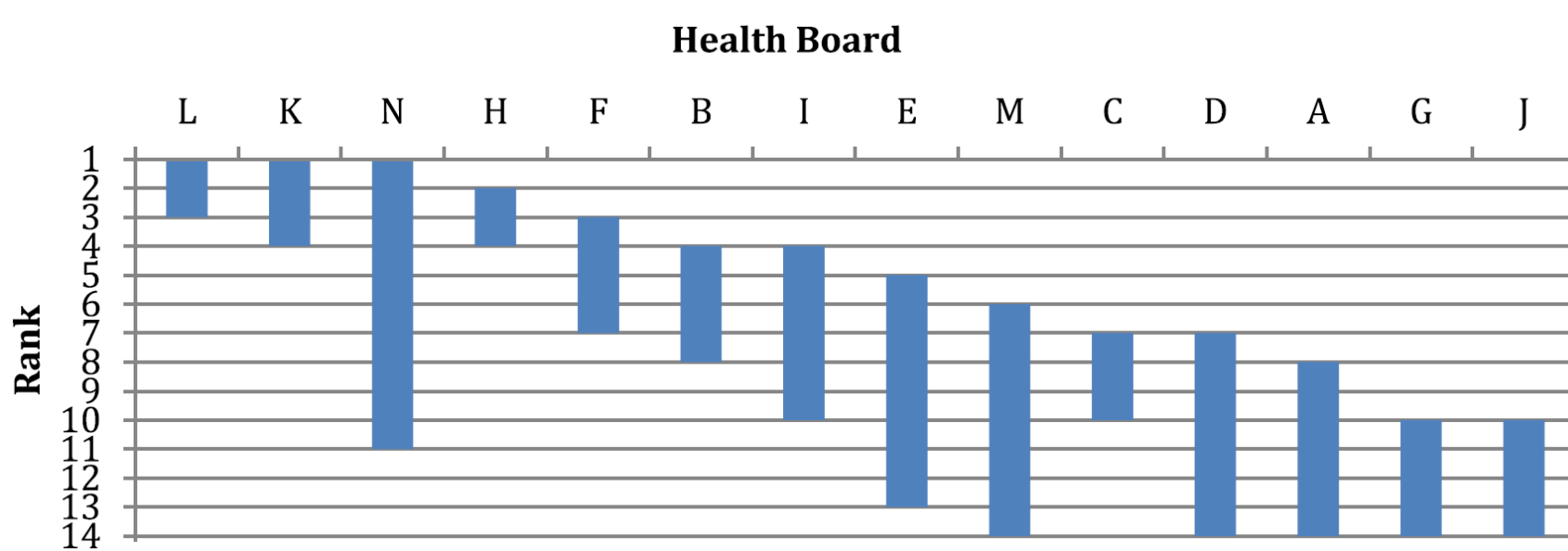
” Avoiding an event of the worst healthcare quality measure cannot be more than ten times as valuable as avoiding an event of the least bad quality measure”

$$u_n \geq 0.1 \forall n$$

“Avoiding C.difficile infection must be at least $\frac{1}{4}$ as valuable as avoiding MRSA/MSSA infection”

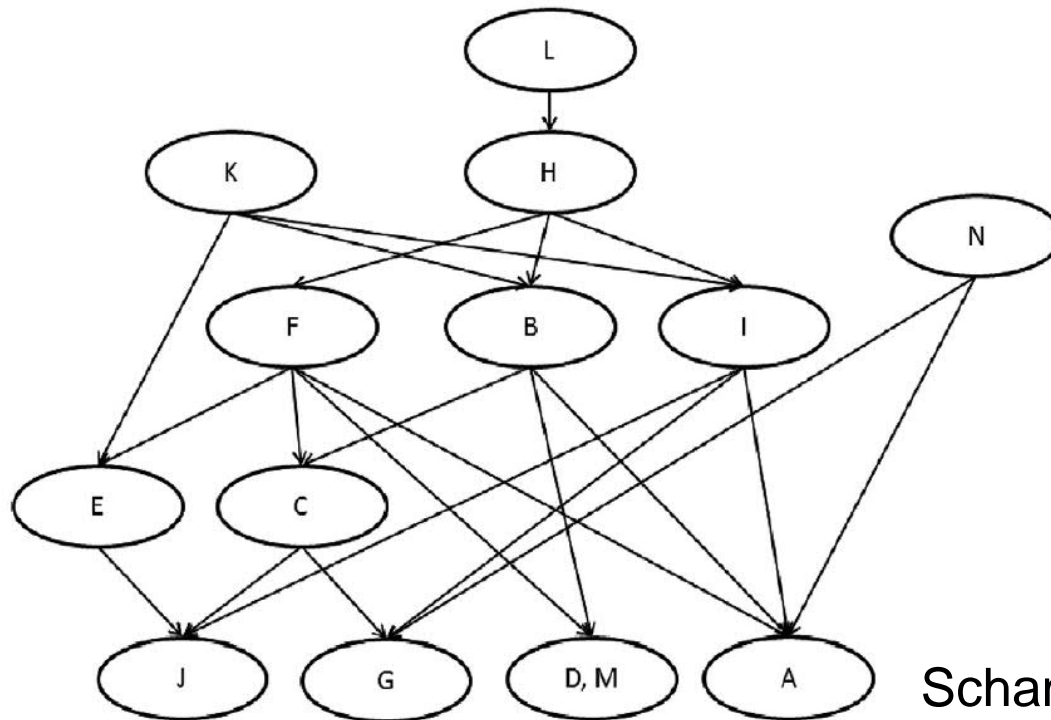
$$u_3 \geq \frac{1}{4} u_1$$

Ordinal & Proportional weighting – Rank intervals



Schang et al. (2016)

Ordinal & Proportional weighing - Dominances



Schang et al. (2016)

Homework

Are the following statements true or false? Justify your answers.

- a) DMU_k can not dominate DMU_l if their ranking intervals overlap.
- b) If DMU_k is dominated by n DMUs. The highest rank DMU_k can possibly have for some u, v is $n+1$.
- c) It always holds that $\bar{D}_{k,\bar{L}} \leq \underline{D}_{k,\underline{L}}$
- d) There can not be any dominances if weights are non-restricted.

Send the answers to lauri.vaara@aalto.fi

DL: 26.11 at 9:00

References

- [1] Salo, A., Punkka, A., 2011: Ranking Intervals and Dominance Relations for Ratio-Based Efficiency Analysis, *Management Science* 57/1, s. 200-214
- [2] Schang, L., Hynninen, Y., Morton, A., Salo, A., 2016: Developing robust composite measures of healthcare quality - Ranking intervals and dominance relations for Scottish Health Boards, *Social Science & Medicine* 162, s. 59-67