

15-03-2022

yesterday

1) Revelation principle

↳ Allows us to look at direct, truthful mechanisms

2) used incentive constraints

(i) $q(\cdot)$ increasing

$$(ii) V(\theta) = V(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} q(s) ds$$

$$\& t(\theta) = V(\theta) - \theta q(\theta) = V(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} q(s) ds - \theta q(\theta)$$

(iii) $V(\underline{\theta}) \geq 0 \Rightarrow IR_{\underline{\theta}}$ holds for all θ

additional material

• Mailath p. 292 \rightarrow

• Bergess p. 07 \rightarrow

3) "order of integration trick"

to get

$$\int_{\underline{\theta}}^{\bar{\theta}} \left[\theta q(\theta) - \int_{\underline{\theta}}^{\theta} q(s) ds \right] f(\theta) d\theta$$

$= -t(\theta)$

$$= \int_{\underline{\theta}}^{\bar{\theta}} J(\theta) \cdot q(\theta) f(\theta) d\theta$$

with $J(\theta) = \theta - \frac{1-F(\theta)}{f(\theta)}$

today

1) some direct extensions

2) formal proofs of steps 2) & 3)

Screening – some immediate extensions of example

Divisible quantity instead of single indivisible good:

- nothing changes
- interpret $q \in [0, \bar{q}]$ as quantity instead of probability $q \in [0, 1]$

Production costs for the seller:

- suppose seller incurs cost $c(q)$ when producing quantity q
- seller's objective is now $-t(\theta) - c(q(\theta))$
- for buyer, nothing changes
- optimality condition for pointwise maximisation (if c convex increasing):

$q^*(\theta)$ determined by

$$\theta - \frac{1 - F(\theta)}{f(\theta)} - c'(q(\theta)) = 0.$$

Screening – incentive compatibility (formally)

We made our lives easy at several steps of the example:

- 1) did not proof formally that q must be increasing to fulfil IC
- 2) did not proof formally that $V'(\theta) = q(\theta)$ must hold to fulfil IC
- (3) Buyer's linear utility $\theta q + t$ seems like (very simple) special case

Let's provide complete proofs of (1) and (2) for more general case when buyer's utility is

$$u(\theta, q) + t. \quad - \text{ "quasi-linear utility"}$$

$$\hookrightarrow \text{yesterday } u(\theta, q) = \theta \cdot q$$

Screening – incentive compatibility – three results

We consider three fundamental results:

1. Envelope Theorem
2. Revenue Equivalence
3. 'Incentive Compatibility Characterisation'

Screening – Envelope Theorem

papers: Milgrom & Segal (2002) *Econometrica*. "Envelope Theorems"

Theorem (Envelope Theorem)

Assume that X is compact, and $\Theta = [\underline{\theta}, \bar{\theta}]$ and $g : \Theta \times X \rightarrow \mathbb{R}$ is differentiable in θ with uniformly bounded derivative*. Suppose the selection $x^*(\theta)$ solves

$$V(\theta) = \max_{x \in X} g(\theta, x).$$

Then we have

$$V'(\theta) = \frac{\partial}{\partial \theta} g(\theta, x^*(\theta)) \text{ a.e.} \quad \text{almost every}$$

and

$$V(\theta) = V(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} \frac{\partial}{\partial \theta} g(s, x^*(s)) ds.$$

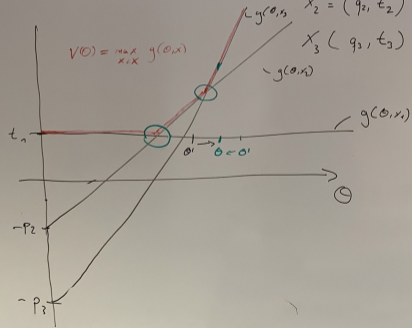
* X is our space of outcomes. In our example $X = [0, 1] \times \mathbb{R} \mid g(\theta; (q, t)) = \theta \cdot q + t$
* there exists $b \in \mathbb{R}$ s.t. $\max_{x \in X} \left| \frac{\partial}{\partial \theta} g(\theta, x) \right| \leq b \forall \theta$
 $x = (q, t)$

example of $V(\theta)$ with $g(\theta, x) = g(\theta, q, t) = \theta \cdot q + t$

suppose $X = \{x_1, x_2, x_3\}$: $x_1 = (q_1, t_1) = (0, t_1)$

$x_2 = (q_2, t_2) = (\frac{1}{2}, -p_2)$

$x_3 = (q_3, t_3) = (1, -p_3)$



Screening – Envelope Theorem – proof

First, we show that $V(\cdot)$ is absolutely continuous so that

$$V'(\cdot) \text{ exists a.e. and } V(\theta) = V(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} V'(s) ds.$$

• For any θ, θ'

$$\begin{aligned} |V(\theta') - V(\theta)| &\leq \max_{x \in X} |g(\theta', x) - g(\theta, x)| = \max_x \left| \int_{\theta}^{\theta'} \frac{\partial}{\partial \theta} g(s, x) ds \right| \\ &\leq \int_{\theta}^{\theta'} \underbrace{\max_x \left| \frac{\partial}{\partial \theta} g(s, x) \right|}_{\leq b} ds \leq \int_{\theta}^{\theta'} b ds \end{aligned}$$

$\Rightarrow V(\cdot)$ is absolutely continuous.

Second, we show that $V'(\theta) = \frac{\partial}{\partial \theta} g(\theta, x^*(\theta))$ a.e.

Take any $\theta \neq \theta'$

$V(\theta') \geq g(\theta', x^*(\theta))$ — true preference θ'
but selection for
other type θ

$$V(\theta') \geq \underbrace{V(\theta) - g(\theta, x^*(\theta))}_{=0 \text{ by definition}} + g(\theta', x^*(\theta))$$

$$\Leftrightarrow V(\theta') - V(\theta) \geq g(\theta', x^*(\theta)) - g(\theta, x^*(\theta))$$

case 1: $\theta' > \theta$: divide by $(\theta' - \theta) > 0$

$$\frac{V(\theta') - V(\theta)}{\theta' - \theta} \geq \frac{g(\theta', x^*(\theta)) - g(\theta, x^*(\theta))}{\theta' - \theta}$$

limit $\theta' \searrow \theta$: $V'(\theta_+) \geq \frac{\partial}{\partial \theta} g(\theta, x^*(\theta))$

case 2: $\theta' < \theta$: divide by $(\theta' - \theta) < 0$

$$\frac{V(\theta') - V(\theta)}{\theta' - \theta} \leq \frac{g(\theta', x^*(\theta)) - g(\theta, x^*(\theta))}{\theta' - \theta}$$

limit $\theta' \nearrow \theta$: $V'(\theta_-) \leq \frac{\partial}{\partial \theta} g(\theta, x^*(\theta))$

since $V'(\cdot)$ exists a.e. i.e. $V'(\theta) = V'(\theta_+) = V'(\theta_-) = \frac{\partial}{\partial \theta} g(\theta, x^*(\theta))$

and thus $V(\theta) = V(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} \frac{\partial}{\partial s} g(s, x^*(s)) ds$ □

Screening – Revenue Equivalence

Theorem (Revenue Equivalence)

Fix a function $q : \theta \rightarrow Q$. Suppose that Q is compact and $\Theta = [\underline{\theta}, \bar{\theta}]$.

Let the agent's utility be $u(\theta, q) + t$, where u is differentiable in θ with uniformly bounded derivative. Any incentive compatible mechanism that implements $q(\theta)$ gives agent payoff

$$V(\theta) = V(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} \frac{\partial}{\partial \theta} u(s, q(s)) ds,$$

transfers must satisfy $-t(\theta) = u(\theta, q(\theta)) - V(\underline{\theta}) - \int_{\underline{\theta}}^{\theta} \frac{\partial}{\partial \theta} u(s, q(s)) ds.$

- By IC, the allocation rule almost completely pins down agent's and principal's payoff
- Only 'degree of freedom' is the constant $V(\underline{\theta})$

proof idea: use envelope Theorem with $g(\theta; (q, t)) = u(\theta, q) + t$
 $(q(\theta), t(\theta))_{\theta}$ is incentive compatible if $x^*(\theta) = (q(\theta), t(\theta))$

Screening – characterisation of IC

The Revenue Equivalence Theorem provides a necessary condition:

'If mechanism is incentive compatible, then (q, t) satisfies...'

Two issues remain:

- (When) are these conditions sufficient for incentive compatibility?
- In our example we said q had to be increasing, where did that come from?

Screening – characterisation of IC

Theorem

Suppose that Q is compact and $\Theta = [\underline{\theta}, \bar{\theta}]$. Let the agent's utility be $u(\theta, q) + t$, where u is differentiable in θ with uniformly bounded derivative.

If $\frac{\partial^2 u(\theta, q)}{\partial q \partial \theta} > 0$, then $(q(\theta), t(\theta))$ is IC if and only if

with $u(\theta, q) = \theta \cdot q$,
 $\frac{\partial^2}{\partial \theta \partial q} = 1 > 0$.

$q(\theta)$ is non-decreasing

and

$$-t(\theta) = u(\theta, q(\theta)) - V(\underline{\theta}) - \int_{\underline{\theta}}^{\theta} \frac{\partial}{\partial \theta} u(s, q(s)) ds.$$

Screening – characterisation of IC – proof

Incentive compatibility means that for all θ and θ'

$$\underline{u(\theta, q(\theta)) + t(\theta)} - [u(\theta, q(\theta')) + \underline{t(\theta')}] \geq 0$$

$$= \underbrace{V(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} \frac{\partial}{\partial \theta} u(s, q(s)) ds}_{\text{orange}} - \left[u(\theta, q(\theta')) - u(\theta', q(\theta')) + \underbrace{V(\underline{\theta}) + \int_{\underline{\theta}}^{\theta'} \frac{\partial}{\partial \theta} u(s, q(s)) ds}_{\text{red}} \right]$$

$$= \underbrace{\int_{\underline{\theta}}^{\theta} \frac{\partial}{\partial \theta} u(s, q(s)) ds - \int_{\underline{\theta}}^{\theta'} \frac{\partial}{\partial \theta} u(s, q(s)) ds}_{\text{blue}} - \left[\underbrace{u(\theta, q(\theta')) - u(\theta', q(\theta'))}_{\text{blue}} \right]$$

$$= \int_{\theta'}^{\theta} \frac{\partial}{\partial \theta} u(s, q(s)) ds$$

$$= \int_{\theta'}^{\theta} \frac{\partial}{\partial \theta} u(s, q(\theta')) ds$$

$$= \int_{\theta'}^{\theta} \left(\frac{\partial}{\partial \theta} u(s, q(s)) - \frac{\partial}{\partial \theta} u(s, q(\theta')) \right) ds = \int_{\theta'}^{\theta} \int_{q(\theta')}^{q(s)} \underbrace{\frac{\partial^2}{\partial \theta \partial q} u(s, z)}_{> 0 \text{ by assumption}} dz ds$$

$$\geq 0 \text{ for all } \theta, \theta' \quad \text{iff } (\theta - \theta')(q(\theta) - q(\theta')) \geq 0 \text{ for all } \theta, \theta' \quad \square$$

Adverse Selection

Mechanism Design

Mechanism design

- how can we aggregate individual preferences into a collective decision?
- especially if individuals' preferences are private information

Compared to the screening problem, we now consider multiple agents

- interests may conflict with each other
- there is increased competition that a seller may exploit
- will inefficiencies increase/decrease?

Mechanism Design – General Setup

The Environment

- n agents
- each agent i has private information (his type) $\theta_i \in \Theta_i$
- set of possible alternatives/outcomes $x \in X$
- each agent is expected-utility maximiser with vNM utility function

$$u_i(\theta, x) \in \mathbb{R}, \quad \text{for } \theta \in \Theta = \Theta_1 \times \dots \times \Theta_n \text{ and } x \in X.$$

- the type profile $\theta = (\theta_1, \dots, \theta_n)$ is distributed according to F with density $f > 0$
- **notation:** we write

$$\theta_{-i} = (\theta_1, \dots, \theta_{i-1}, \theta_{i+1}, \dots, \theta_n) \quad \text{and} \quad (\theta_i, \theta_{-i}) = \theta$$

\downarrow
 θ_i is missing

Mechanism Design – Setup – some terminology

Private Values

IPV = independent
Private values

- i 's preferences depend only on θ_i :

$$u_i(\theta, x) = u_i(\theta_i, x)$$

- 'interdependent values' otherwise

Quasi-linear Utilities

- outcomes $X = K \times \mathbb{R}^n$, where

$k \in K$ some physical allocation,

$t = (t_1, \dots, t_n) \in \mathbb{R}^n$ transfers

- i 's utility is linear in money (his transfer):

$$u_i(\theta, x) = v_i(\theta, k) + t_i$$

Independent Types

- θ_i 's distribution indep. of other types θ_{-i} :

$$f(\theta) = \prod_{i=1}^n f_i(\theta_i)$$

- 'correlated' types otherwise

Social Choice/Unrestricted Domain

- $X = \{a, b, \dots\}$ finite set of alternatives

- θ_i gives ranking over alternatives:

$$a \theta_i b \Leftrightarrow a \succ_i b$$

- Unrestricted domain if

Θ_i contains all possible rankings over X

Mechanism Design – environment examples

Ex 1. Public good

- outcomes $(k, t) \in X = \{0, 1\} \times \mathbb{R}^n$
 $k \in \{0, 1\}$ with $k = 1$ if bridge is built
 $t_i \in \mathbb{R}$ transfer to agent i
- θ_i is i 's willingness to pay for bridge
 $u_i(\theta, x) = \theta_i k + t_i$

private values

quasi-linear utilities

Ex. 2 Allocation with externalities

- outcomes $(k, t) \in X = \{0, 1, \dots, n\} \times \mathbb{R}^n$
 $k = \begin{cases} 0 & \text{if nobody gets object} \\ i & \text{if agent } i \text{ gets object} \end{cases}$
 $t_i \in \mathbb{R}$ transfer to agent i
- $\theta_i = (\theta_i^i, \theta_i^x)$ with utility

$$u_i(\theta, x) = \begin{cases} t_i & \text{if } k = 0 \\ \theta_i^i + t_i & \text{if } k = i \\ -\theta_i^x + t_i & \text{if } k \notin \{0, i\} \end{cases}$$

quasi-linear utility

private values

Mechanism Design – Social Choice Functions

Our goal is generally to choose a **good** outcome $x \in X$ given the realised preferences $\theta \in \Theta$

Definition (Social Choice Function)

A **social choice function** (scf) $\xi: \Theta \rightarrow X$ assigns to each type profile $\theta \in \Theta$ an alternative $\xi(\theta) \in X$.

The problem of the mechanism designer is not 'lack of power'

- if the designer knew θ , she could always choose the 'optimal' outcome

The problem is 'just' the asymmetric information

Mechanism Design – Mechanisms (general/indirect)

Typically, social (collective) outcomes are determined through interaction in some institution

Definition (Mechanism)

A **mechanism** $\Gamma = (S_1, \dots, S_n, g)$ consists of

- a strategy space S_i for each agent i
- an outcome function $g : S_1 \times \dots \times S_n \rightarrow X$.

A mechanism $\Gamma = (S, g)$ together with the environment induces a Bayesian game:

$G_\Gamma = (n, \{S_i\}_{i \leq n}, \{\tilde{u}_i\}_{i \leq n}, \Theta, F)$, with payoffs $\tilde{u}_i(\theta, s_1, \dots, s_n) = u_i(\theta, g(s_1, \dots, s_n))$.

Mechanism Design – Incentive Compatibility

We have several solution concepts: Let $(s_i^*)_{i=1}^n$ be a strategy profile, where $\forall i: s_i : \Theta_i \rightarrow S_i$

- Dominant strategy equilibrium: for all i, θ_i, s_i :

$$u_i((\theta_i, \theta_{-i}), g(s_i^*(\theta_i), s_{-i})) \geq u_i((\theta_i, \theta_{-i}), g(s_i, s_{-i})) \quad \forall \theta_{-i}, s_{-i}$$

- Ex-post equilibrium: for all i, θ_i, s_i :

$$u_i((\theta_i, \theta_{-i}), g(s_i^*(\theta_i), s_{-i}^*(\theta_{-i}))) \geq u_i((\theta_i, \theta_{-i}), g(s_i, s_{-i}^*(\theta_{-i}))) \quad \forall \theta_{-i}$$

- Bayes-Nash equilibrium: for all i, θ_i, s_i :

$$\begin{aligned} & \int_{\Theta_{-i}} u_i((\theta_i, \theta_{-i}), g(s_i^*(\theta_i), s_{-i}^*(\theta_{-i}))) \, dF_{-i}(\theta_{-i}|\theta_i) \\ & \geq \int_{\Theta_{-i}} u_i((\theta_i, \theta_{-i}), g(s_i, s_{-i}^*(\theta_{-i}))) \, dF_{-i}(\theta_{-i}|\theta_i) \end{aligned}$$

Mechanism Design – Participation Constraints

Definition

We say that mechanism $\Gamma = (S, g)$ [...] **implements** scf ξ if there exists a [...] equilibrium strategy profile $(s_i^*)_{i=1}^n$ such that

$$g(s_1^*(\theta_1), \dots, s_n^*(\theta_n)) = \xi(\theta) \quad \text{for all } \theta \in \Theta.$$

where [...] \in {dominant strategy, ex-post, Bayes}

- Full implementation: every equilibrium results in $\xi(\theta)$
- Partial implementation: there is an equilibrium that results in $\xi(\theta)$

We focus on partial implementation

Mechanism Design – example second-price auction

Theorem (Revelation Principle)

For any mechanism $\Gamma = (S, g)$ and $[\dots]$ -equilibrium strategy profile $(s_i^*)_{i=1}^n$ that implements scf ξ , there exists a **direct** mechanism $\hat{\Gamma} = (\Theta, \xi)$ such that **truth-telling** is a $[\dots]$ equilibrium.

- Only ensures that there is AN equilibrium
- In different (indirect) mechanisms sharing the same direct mechanism other equilibria may arise