<u> 28-3-22</u>

lecture 5

last week

- » Efficient (= welfase max.) mechanisms
  - VCG mechanisms for DIC (pivot mechanism in porticulas)
  - · Expected externality mechanism for BIC
- A Budget Balance and IR P Ryeison-Sattesthwaite Jhm

today

Revenue maximisation "optimal auctions"

## The auction problem:

- single indivisible object
- seller cost<u></u>
- n potential buyers with type  $heta_i$
- utility  $\theta_i q_i + t_i$
- types are independently distributed on  $[\underline{\theta}_i, \overline{\theta}_i]$  according to  $F_i$  with density  $f_i > 0$

 $\mathbb{E}\left[-\overset{\circ}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\overset{i=1}{\underset{i=1}{\atopi=1}{\underset{i=1}{\atopi=1}{\underset{i=1}{\underset{i=1}{\atopi=1}{\underset{i=1}{\underset{i=1}{\underset{i=1}{\atopi=1}{\underset{i=1}{\atopi=1}{\atopi=1}{\underset{i=1}{\atopi=1}{\underset{i=1}{\atopi=1}{\underset{i=1}{\atopi=1}{\underset{i=1}{\atopi=1}{\underset{i=1}{\atopi=1}{\underset{i=1}{\atopi=1}{\atopi=1}{\underset{i=1}{\atopi=1}{\underset{i=1}{\atopi=1}$ 

• feasible allocation probabilities:  $q_i(\theta) \in [0,1]$  with  $\sum_{i=1}^n q_i(\theta) \leq 1$ 

Seller commits to mechanism  $(q,t) \colon \Theta \to [0,1]^n \times \mathbb{R}^n$  to maximise revenue

# Mechanism Design – Revenue Maximisation – Optimal Auctions

Let's compare different auction formats for the example n = 2,  $\theta_i \sim U([0, 1])$ , c = 0

#### 1. First-price auction

- Each bidder makes a bid  $b = s^{\text{FPA}}(\theta_i)$ . The highest bid wins. The winner pays his bid.
- Find the symmetric equilibrium bid function s<sup>FPA</sup> (hint: linear function)

## 2. All-pay auction (=contest)

- Each bidder makes a bid  $b = s^{APA}(\theta_i)$ . The highest bid wins. Each bidder pays his bid.
- Find the symmetric equilibrium bid function  $s^{APA}$ .

## 3. English auction (=ascending-clock auction)

• A price is publicly displayed and increases continuously from  $p_0 = 0$ . Bidder i drops out when the price reaches  $s^{\text{SPA}}(\theta_i)$ .

When i drops out (first),  $j \neq i$  wins and pays p at which i dropped out.

 $\rightarrow$  • What is the weakly dominant stopping strategy  $s^{
m SPA}$ ?

#### **Q:** What is the expected revenue of the seller?

First-price - auction expected revenue • What is biddes I's expected payoff if his valuation is  $\Theta_r$ , bidde 2 uses bidding strategy  $S(\cdot)$ , and bidder 1 bids  $D_r$ ?  $(\Theta_r - b_r) \cdot \Pr[S(\Theta_2) \leq b_r]$   $= (\Theta_r - b_r) \Pr[S(\Theta_2) \leq S(\Theta_r)]$  $\max_{b_{1}} \left\{ (0, -b_{1}) : s^{-1}(b_{2}) \right\} = s^{-1}(b_{2}) \left\{ s(x) \in a + m \cdot 0 \\ s(0) = a + m \cdot 0 \\ s(0)$  $F\alpha:\frac{\partial}{\partial b_{n}} \frac{1}{2} \dots \frac{1}{2} = 0 \quad <=> -\frac{5(b_{n})}{5} + \frac{b_{n}}{5(5'(b_{n}))} = 0 \quad S(0) \quad must \quad be \quad 0$  $(z-y) - \Theta_{1} + \frac{\Theta_{1} - y(\Theta_{1})}{y'(\Theta_{1})} = (y-y)(\Theta_{1}) + y(\Theta_{1})$  $\begin{array}{c} \zeta = \gamma \quad S'(\Theta_{n}) \cdot \Theta_{n} = \Theta_{n} - S(\Theta_{n}) \\ \end{array} \begin{array}{c} \zeta = \gamma \quad S^{FM}(\theta) - \frac{1}{2} \Theta_{n} \\ \end{array}$ 

All-pay - auction expected revenue . What is Siddes I's expected payoff What is biddes I's expected payoff if his valuation is  $\Theta_r$ , sidle 2 uses bidding stratagy S(i), and bidde 1 bids  $b_r$ ?  $\Theta_r \operatorname{Fi}[S(\Theta_2) \leq b_r] = b_r$   $E_{an}[\max \{i \Theta_r, i\Theta_i\}]$   $= \frac{i}{2} \int (1 - F_r^2 \Theta_r) d\Theta = \frac{i}{3}$   $= \frac{i}{2} \int (1 - F_r^2 \Theta_r) d\Theta = \frac{i}{3}$ - {O, R [O2 = 51(6)] - 61 }  $\begin{array}{c}
\underline{APA} \\
\underline{E}_{\boldsymbol{\sigma},\boldsymbol{\sigma}_{1}}\left[\frac{1}{2}\boldsymbol{\sigma}_{1}^{2}+\frac{1}{2}\boldsymbol{\sigma}_{2}^{2}\right] = E\left[\boldsymbol{\sigma}_{1}^{2}\right], \frac{1}{3}
\end{array}$  $\frac{D}{\partial L_n} \begin{bmatrix} 1 \\ -1 \end{bmatrix} = \underbrace{\Theta_n}_{S'(G_n)} - 1 = 0 \iff \underbrace{S'(G_n)}_{S'(G_n)} = \Theta_n$ <

English - auction expected revenue Strategically equivalent to 2<sup>nd</sup> price anchien. S<sup>PM</sup> (0) = 0, is dominant strategy. => <u>expected Revenue</u> with ...2 (in <u>TPV</u>  $E_{a,a}[mix]_{a,b}^{a}(b_{1})]$ =>  $\frac{expected Revenue}{E_{a,a}[mix]_{a}^{b}(b_{1})]} = \frac{1}{3}$  $\begin{array}{c} \underline{APA} \\ \underline{E}_{a,a} \begin{bmatrix} \frac{1}{2}a^{2} + \frac{1}{2}a_{1}^{2} \end{bmatrix} = E\begin{bmatrix} a^{2} \end{bmatrix} \cdot \frac{1}{3} \end{array}$ expected payof to seller.