# MS-E2122 - Nonlinear Optimization Lecture V

#### Fernando Dias

Department of Mathematics and Systems Analysis

Aalto University School of Science

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#### Outline of this lecture

#### Line search methods - univariate (single variable) optimisation

Line searches without derivatives

Line searches with derivatives

#### Methods for unconstrained optimisation

Coordinate descent

Gradient method

Newton's method

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## Last Week

- Optimality conditions;
- KKT conditions

Last week...

Fernando Dias

#### Outline of this lecture

### Line search methods - univariate (single variable) optimisation

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Line searches with derivatives

#### Methods for unconstrained optimisation

Coordinate descent

Gradient method

Newton's method

This pseudocode represents the general concept of optimization methods:

#### Algorithm Conceptual optimisation algorithm

- 1: **initialise.** iteration count k = 0, starting point  $x_0$
- 2: while stopping criteria are not met do
- 3: compute direction  $d_k$
- 4: compute step size  $\lambda_k$
- $5: x_{k+1} = x_k + \lambda_k d_k$
- 6: k = k + 1
- 7: end while
- 8: return  $x_k$

#### where

- k is an iteration counter;
- $\triangleright \lambda_k$  is a suitable step size;
- $ightharpoonup d_k$  is a direction vector;

Finding an optimal step size  $\lambda_k$  is in itself an optimisation problem called line search due to its **unidimensional** nature.

Line searches are the backbone of most optimisation methods.

Let  $\theta(\lambda)=f(x+\lambda d)$ . If f is **differentiable**, a straightforward approach is to find an optimal setup size  $\lambda$  is

$$\theta'(\lambda) = d^{\top} \nabla f(x + \lambda d) = 0$$

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However, one must bear in mind that:

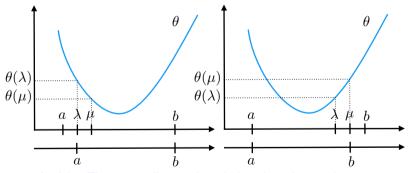
- ▶  $d^{\top} f(x + \lambda d)$  is often nonlinear in  $\lambda$ ;
- $\overline{\lambda} = \operatorname{argmin}_{\lambda} d^{\top} \nabla f(x + \lambda d) = 0$  is not necessarily optimal.

## Theorem 1 (Line search reduction)

Let  $\theta: \mathbb{R} \to \mathbb{R}$  be strictly quasiconvex over the interval [a,b], and let  $\lambda, \mu \in [a,b]$  such that  $\lambda < \mu$ . If  $\theta(\lambda) > \theta(\mu)$ , then  $\theta(z) \geq \theta(\mu)$  for all  $z \in [a,\lambda]$ . If  $\theta(\lambda) \leq \theta(\mu)$ , then  $\theta(z) \geq \theta(\lambda)$  for all  $z \in [\mu,b]$ .

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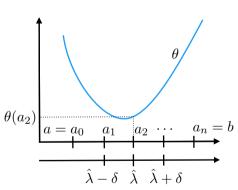
Applying Theorem 1 allows to iteratively reduce the search space.

### Line search methods - uniform search

Break [a,b] into n uniform intervals of size  $\delta$ , which leads to n+1 grid points  $a_k=a_0+k\delta$ , with  $a=a_0,b=a_n$ , and  $k=0,\ldots,n$ .

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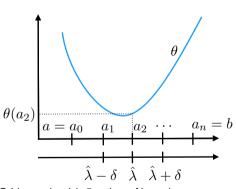


Let  $\hat{\lambda} = \operatorname{argmin}_{i=0,\dots,n} \theta(a_i)$ . We know that the optimal  $\overline{\lambda} \in [\hat{\lambda} - \delta, \hat{\lambda} + \delta]$ .

Grid search with 5 points; Note that  $\theta(a_2) = \min_{i=0,...,n} \theta(a_i)$ .

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#### Remarks:

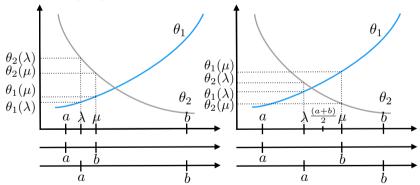
- The search can be repeated making  $a = \hat{\lambda} \delta$  and  $b = \hat{\lambda} + \delta$ .
- The number of grid points can increase dynamically, saving function evaluations.

More efficient methods can be devised by using information of the **previous** evaluation of  $\theta$ . These are known as sequential searches.

1. Dichotomous search: we place two points,  $\lambda$  and  $\mu$ , around the midpoint of [a,b] at a small distance  $\epsilon$ .

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Using the midpoint (a+b)/2 and Theorem 1 to reduce the search space.

#### Algorithm Dichotomous search

```
1: initialise. distance \epsilon > 0, tolerance l > 0, [a_0,b_0] = [a,b], k = 0
2: while b_k - a_k > l do
3: \lambda_k = \frac{a_k + b_k}{2} - \epsilon, \mu_k = \frac{a_k + b_k}{2} + \epsilon
4: if \theta(\lambda_k) < \theta(\mu_k) then
5: a_{k+1} = a_k, b_{k+1} = \mu_k
6: else
7: a_{k+1} = \lambda_k, b_{k+1} = b_k
8: end if
9: k = k + 1
10: end while
11: return \overline{\lambda} = \frac{a_k + b_k}{2}
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**Remark:** The number of steps (and evaluations of  $\theta$ ) can be predicted beforehand:

$$b_{k+1} - a_{k+1} = \frac{1}{2^k}(b_0 - a_0) + 2\epsilon \left(1 - \frac{1}{2^k}\right).$$

More than just function evaluations, it also uses derivative information. We assume  $\theta(\lambda)$  to be **differentiable** and **convex**.

- 2. **Bisection method:** The main idea is
  - 1. if  $\theta'(\lambda_k) = 0$ , then  $\lambda_k$  is a **minimiser**.

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  - 4. As in the dichotomous search, to maximise overall interval reduction, we set  $\lambda_k = \frac{1}{2}(b_k + a_k)$ .

#### Algorithm Bisection method

```
1: initialise. tolerance l > 0, [a_0, b_0] = [a, b], k = 0
 2: while b_k - a_k > l do
        \lambda_k = \frac{(b_k + a_k)}{2} and evaluate \theta'(\lambda_k)
         if \theta'(\lambda_k) = 0 then return \lambda_k
 5:
        else if \theta'(\lambda_k) > 0 then
               a_{k+1} = a_k, b_{k+1} = \lambda_k
        else
 8:
               a_{k+1} = \lambda_k, b_{k+1} = b_k
          end if
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# Inexact line searches - Armijo rule

Often the use of non-optimal (i.e., inexact) step sizes  $\lambda_k$  is enough to guarantee a good performance.

**Armijo's rule**: find acceptable step sizes by balancing the trade-off between convergence and numerical performance.

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$$f(\overline{x} + d\overline{\lambda}) - f(\overline{x}) \le \alpha \overline{\lambda} \nabla f(\overline{x})^{\top} d$$

which, at  $\lambda = 0$ , is the same as

$$\theta(\overline{\lambda}) - \theta(0) \le \alpha \overline{\lambda} \theta'(0)$$

$$\theta(\overline{\lambda}) \le \theta(0) + \alpha \overline{\lambda} \theta'(0)$$
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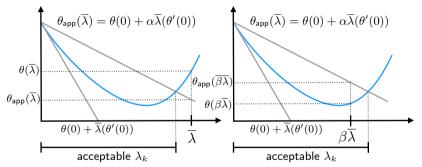
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 $\theta(\overline{\lambda}) \le \theta(0) + \alpha \overline{\lambda} \theta'(0)$ : Armijo's rule (AR)

If  $\overline{\lambda}$  does not satisfy AR,  $\overline{\lambda}$  is reduced by a factor  $\beta \in (0,1)$  and the test is repeated until (AR) is satisfied.

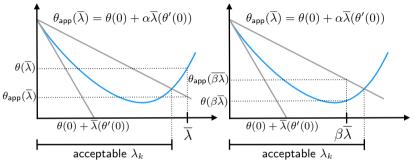
# Inexact line searches - Armijo's rule

Armijo's rule has a nice graphical interpretation:  $\overline{\lambda}$  is accepted if it is in an interval where the function  $\theta(\lambda)$  is below a deflected linear extrapolation (from 0).



At first  $\lambda_0 = \overline{\lambda}$  is not acceptable; after reducing the step size to  $\lambda_1 = \beta \overline{\lambda}$ , it enters the acceptable range where  $\theta(\lambda_k) \leq \theta_{\mathsf{app}}(\lambda_k) = \theta(0) + \alpha \lambda_k(\theta'(0))$ .

# Inexact line searches - Armijo's rule



#### Remarks:

- 1. Some variants also consider rules to guarantee that  $\overline{\lambda}$  is not too small, such as  $\theta(\delta\overline{\lambda}) \leq \theta(0) + \alpha\delta\overline{\lambda}\theta'(0)$ , with  $\delta > 1$ .
- 2. Also known in the literature as backtracking.
- 3. Typical values:  $\alpha \in [0.1, 0.5]$  and  $\beta \in [0.6, 0.99]$ . Very small  $\alpha$ , e.g.  $10^{-4}$  is often used as well.

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#### Line search methods - univariate (single variable) optimisation

Line searches without derivatives

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#### Methods for unconstrained optimisation

Coordinate descent

Gradient method

Newton's method

Next, we focus on optimizing functions  $f:\mathbb{R}^n\to\mathbb{R}$  with **more** than one dimension.

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1. Coordinate descent: the search direction is one coordinate axis per iteration. That is,  $d_i=1$  for coordinate i and  $d_{j\neq i}=0$ , for  $i,j\in\{1,\ldots,n\}$ .

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Several variants:

- 1. Cyclic: coordinates are considered in order  $1, \ldots, n$ ;
- Double-sweep: swap the coordinate order at each iteration;
- 3. **Gauss-Southwell:** choose components with largest  $\frac{\partial f(x)}{\partial x_i}$ ;
- 4. Stochastic: coordinates are selected at random

#### Algorithm Coordinate descent method (cyclic)

```
1: initialise. tolerance \epsilon > 0, initial point x^0, iteration count k = 0
2: while ||x^{k+1} - x^k|| > \epsilon do
3: for j = 1, \dots n do
4: d = \{d_i = 1, \text{ if } i = j; d_i = 0, \text{ if } i \neq j\}
5: \overline{\lambda}_j = \mathop{\mathrm{argmin}}_{\lambda \in \mathbb{R}} \{f(x_j^k + \lambda d_j)\}
6: x_j^{k+1} = x_j^k + \overline{\lambda}_j d_j
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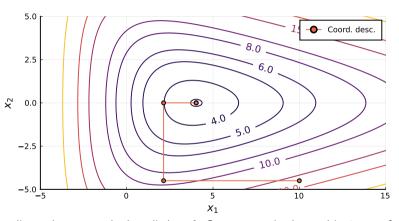
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#### Remarks:

- 1. The one-dimensional minimisation is called Gauss-Seidel step;
- 2. Block-coordinate methods use subgroups (blocks) of coordinates to define directions.

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



Coordinate descent method applied to f. Convergence is observed in 4 steps for a tolerance  $\epsilon=10^{-4}\,$ 

#### Gradient method

Recall that if d is a **descent direction**, there exists  $\delta>0$  such that  $f(x+\lambda d)< f(x)$  for all  $\lambda\in(0,\delta)$ . The following result provides directions of steepest descent.

## Lemma 2 (Steepest descent direction)

Suppose that  $f: \mathbb{R}^n \to \mathbb{R}$  is differentiable at  $x \in \mathbb{R}^n$  and  $\nabla f(x) \neq 0$ . Then  $\overline{d} = -\frac{\nabla f(x)}{||\nabla f(x)||}$  is the direction of steepest descent of f at x.

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#### Proof.

From differentiability of f, we have

$$f'(x;d) = \lim_{\lambda \to 0^+} \frac{f(x+\lambda d) - f(x)}{\lambda} = \nabla f(x)^{\top} d.$$

Thus, 
$$\overline{d} = \operatorname{argmin}_{||d|| \le 1} \left\{ \nabla f(x)^{\top} d \right\} = -\frac{\nabla f(x)}{||\nabla f(x)||}.$$

## Gradient method

#### Algorithm Gradient method

- 1: **initialise.** tolerance  $\epsilon > 0$ , initial point  $x_0$ , iteration count k = 0.
- 2: while  $||\nabla f(x_k)|| > \epsilon$  do
- 3:  $d = -\frac{\nabla f(x_k)}{||\nabla f(\overline{x})||}$
- 4:  $\overline{\lambda} = \operatorname{argmin}_{\lambda \in \mathbb{R}} \{ f(x_k + \lambda d) \}$
- 5:  $x_{k+1} = x_k + \overline{\lambda} d_j$
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### Remarks:

1. Steepest descent and gradient methods are different. When  $||d|| \le 1$  uses 2-norm (in Lemma 2), they are equivalent;

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- 1. Steepest descent and gradient methods are different. When  $||d|| \le 1$  uses 2-norm (in Lemma 2), they are equivalent;
- Poor convergence and zigzagging can be observed due to imprecise linear approximations (more on this later);

## Gradient method

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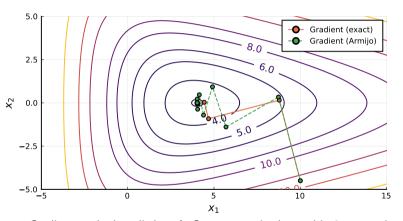


Figure: Gradient method applied to f. Convergence is observed in 9 steps using exact line search and 15 using Armijo's rule ( $\epsilon=10^{-4}$ )

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Deflection is achieved using the Hessian, which is equivalent to relying on quadratic approximations (rather than linear).

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Consider the 2<sup>nd</sup>-order approximation of f at  $x_k$ :

$$q(x) = f(x_k) + \nabla f(x_k)^{\top} (x - x_k) + \frac{1}{2} (x - x_k)^{\top} H(x_k) (x - x_k),$$

where  $H(x_k)$  is the Hessian at  $x_k$ .

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where  $H(x_k)$  is the Hessian at  $x_k$ . We require that  $\nabla q(x_{k+1})=0$ , which leads to

$$\nabla f(x_k) + H(x_k)(x - x_k) = 0.$$

Assuming that  $H^{-1}(x_k)$  exists, we obtain the update rule

$$x_{k+1} = x_k - H^{-1}(x_k)\nabla f(x_k).$$

#### Algorithm Newton's method

```
1: initialise. tolerance \epsilon > 0, initial point x_0, iteration count k = 0
```

2: while 
$$||\nabla f(x_k)|| > \epsilon$$
 do  
3:  $d = -H^{-1}(x_k)\nabla f(x_k)$ 

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- 1. Setting  $\overline{\lambda} = 1$  recovers the "pure" Newton's method;
- 2. As  $\nabla f(x_k)$  gets close to 0,  $H^{-1}(x_k)$  becomes singular;

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2: while ||\nabla f(x_k)|| > \epsilon do
3: d = -H^{-1}(x_k)\nabla f(x_k)
4: \overline{\lambda} = \mathop{\rm argmin}_{\lambda \in \mathbb{R}} \left\{ f(x_k + \lambda d) \right\}
5: x_{k+1} = x_k + \lambda d
6: k = k+1
7: end while
8: return x_k
```

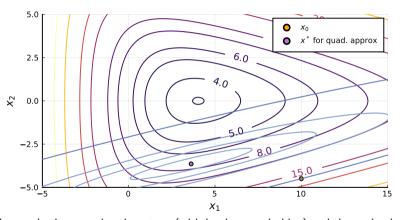
- 1. Setting  $\overline{\lambda} = 1$  recovers the "pure" Newton's method;
- 2. As  $\nabla f(x_k)$  gets close to 0,  $H^{-1}(x_k)$  becomes singular;
- 3. It might not converge if  $x_0$  is too far from optimal and fixed step size is used;

### Algorithm Newton's method

```
1: initialise. tolerance \epsilon>0, initial point x_0, iteration count k=0
2: while ||\nabla f(x_k)||>\epsilon do
3: d=-H^{-1}(x_k)\nabla f(x_k)
4: \overline{\lambda}=\mathrm{argmin}_{\lambda\in\mathbb{R}}\left\{f(x_k+\lambda d)\right\}
5: x_{k+1}=x_k+\overline{\lambda}d
6: k=k+1
7: end while
8: return x_k
```

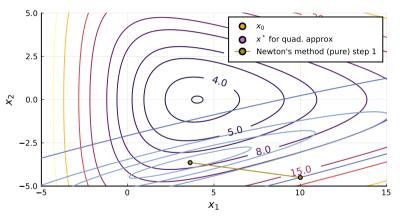
- 1. Setting  $\overline{\lambda} = 1$  recovers the "pure" Newton's method;
- 2. As  $\nabla f(x_k)$  gets close to 0,  $H^{-1}(x_k)$  becomes singular;
- 3. It might not converge if  $x_0$  is too far from optimal and fixed step size is used;
- Levenberg-Marquardt method and other trust-region method variants also address convergence issues of Newton's method.

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



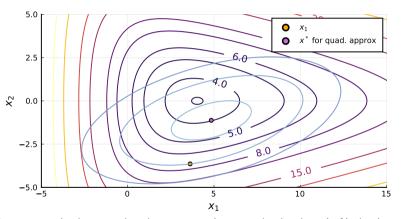
The quadratic approximation at  $x_0$  (with level curves in blue) and the optimal point  $x^{\ast}$ .

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



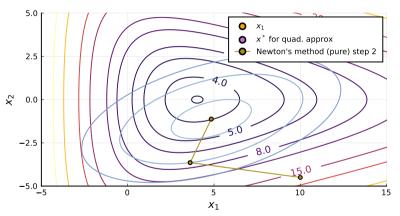
The new point  $x_1$  becomes  $x^*$ .

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



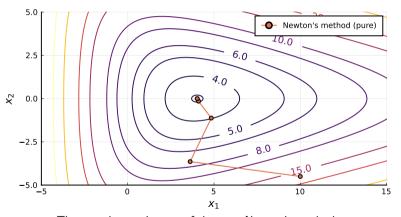
The new quadratic approximation at  $x_1$  and new optimal point  $x^*$ . Notice how the approximation improved.

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



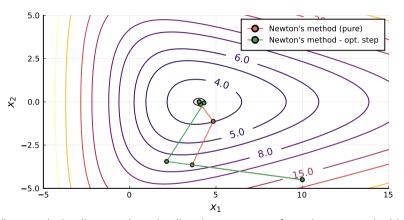
The new point  $x_2$  becomes  $x^*$ .

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



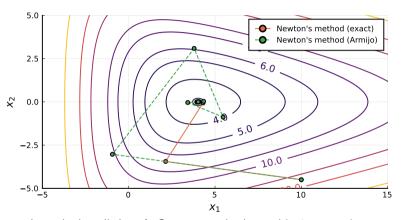
The complete trajectory of the pure Newton's method.

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



When employing line searches, the direction  $x_k-x_{k-1}$  from the pure method is used, but the actual step is optimised.

$$f(x) = e^{(-(x_1-3)/2)} + e^{((4x_2+x_1)/10)} + e^{((-4x_2+x_1)/10)}$$



Newton's method applied to f. Convergence is observed in 4 steps using exact line search and 27 using Armijo's rule ( $\epsilon=10^{-4}$ )

