



Aalto University  
School of Business

# Session 6: Vehicle routing problems

## 35E00750 Logistics Systems and Analytics

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# Learning objectives

- Understand the Traveling Salesman Problem (TSP) formulations
- Understand Vehicle Routing Problem (VRP) formulations



After this lecture, VRP will not be a “black box”

*Slide courtesy of Dr. Ilke Bakir, Associate Professor, Department of Operations, Faculty of Economics and Business, University of Groningen*

# Course contents

## Part 1. Background

1. Understanding supply chains
2. Achieving supply chain fit
3. Mathematical programming for Logistics & SCM
4. Guest lecture: Janne Kilpua

## Part 2. Transportation

5. Urban logistics
- 6. Vehicle routing problems**

## Part 3. Facilities

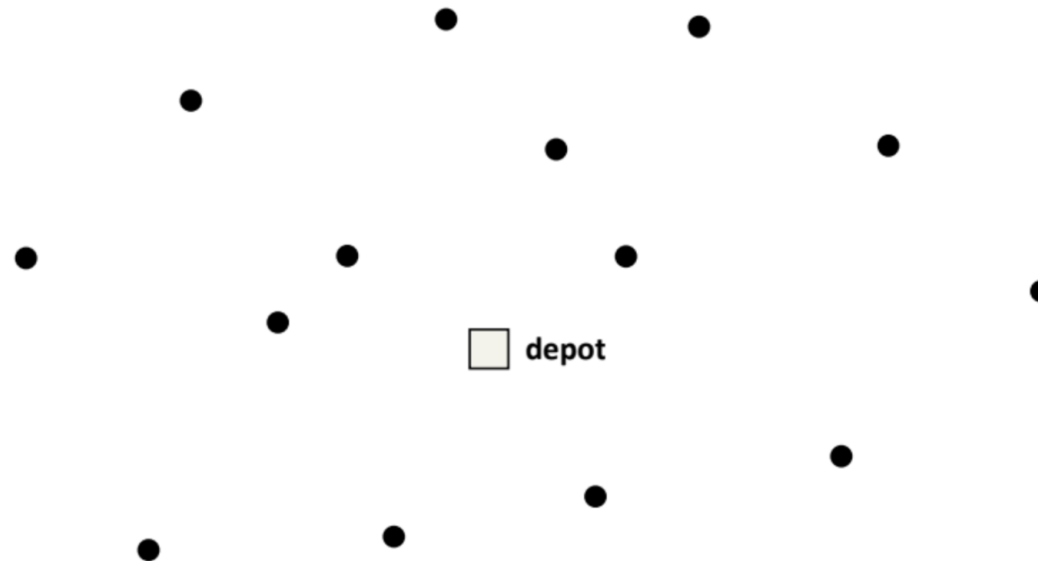
7. Warehousing technologies
8. Guest lecture: Vesa Hämetvaara (Konecranes)
9. Facility location problems

## Part 4. Data

10. Digital logistics
11. Logistical drivers and metrics

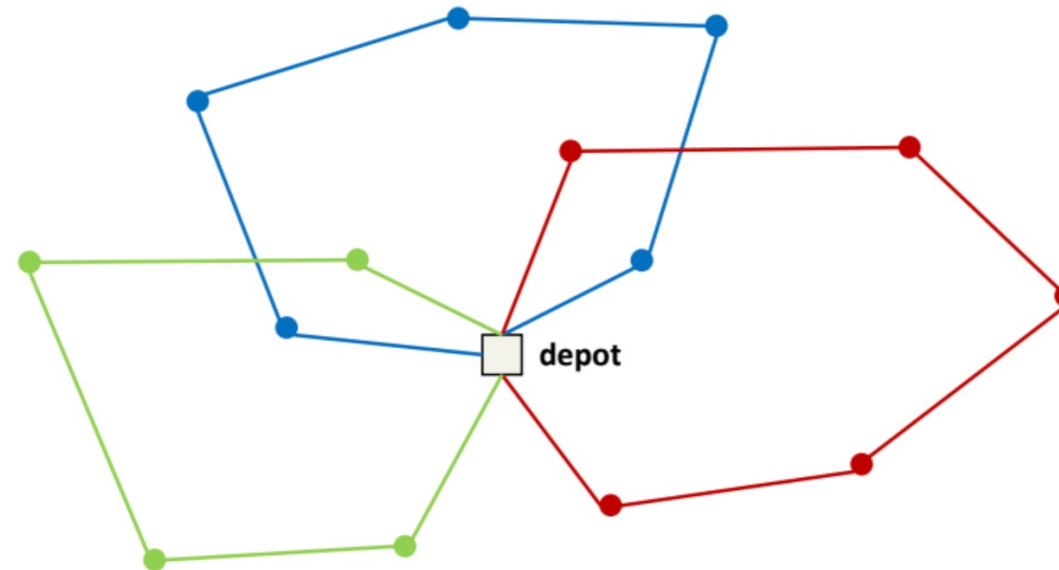
# Vehicle Routing Problems (VRP)

The problem of **constructing routes** from/to one or several depot(s) to visit a number of geographically dispersed locations to pick up or deliver goods



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The problem of **constructing routes** from/to one or several depot(s) to visit a number of geographically dispersed locations to pick up or deliver goods

**So that:**

- Each city is visited exactly once by exactly one vehicle
- All routes start and end at a depot

# Single Vehicle Route: Traveling Salesman Problem (TSP)

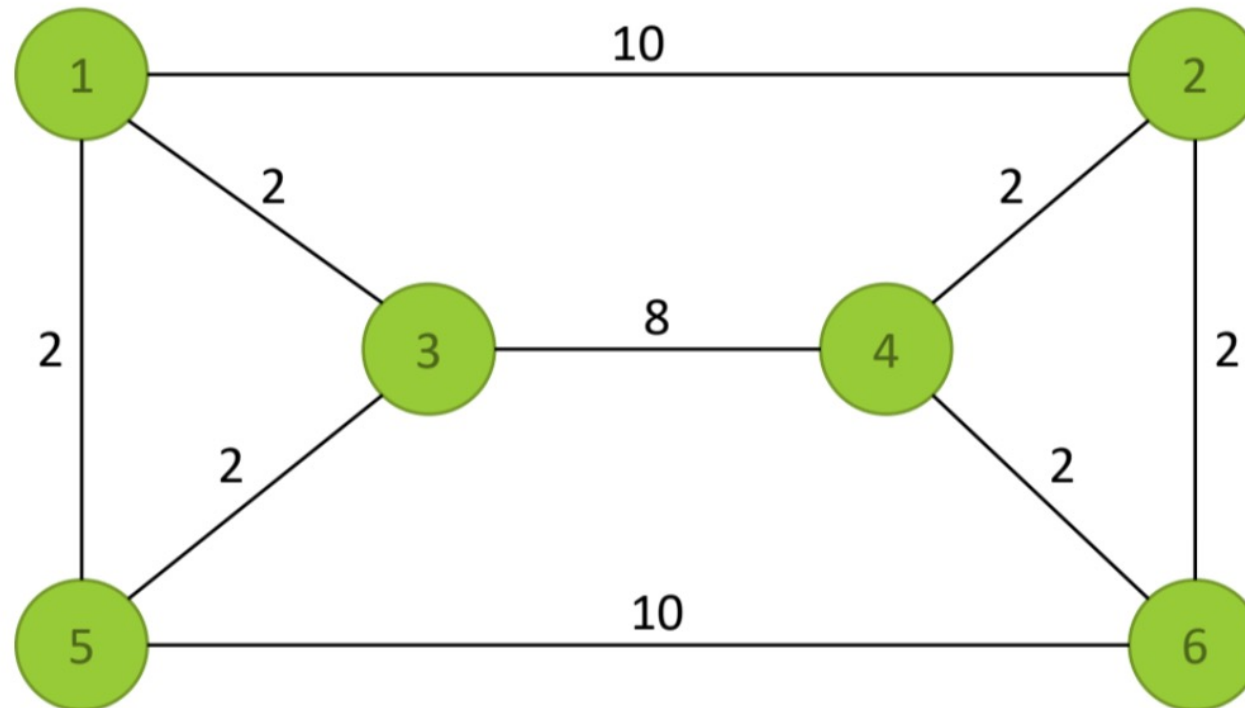
- The Traveling Salesman Problem (TSP) seeks a route with a **minimum total length** visiting **every point** in a given set **exactly once**.
- **More formal description:**
  - *“Suppose a salesman would like to visit several cities. Find the shortest route that visits each city exactly once and returns the salesman back to where they started”*

# TSP Networks



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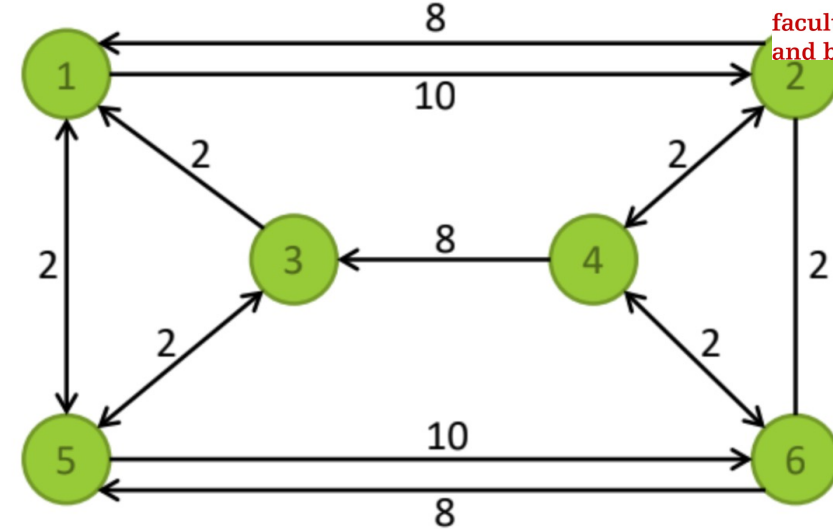
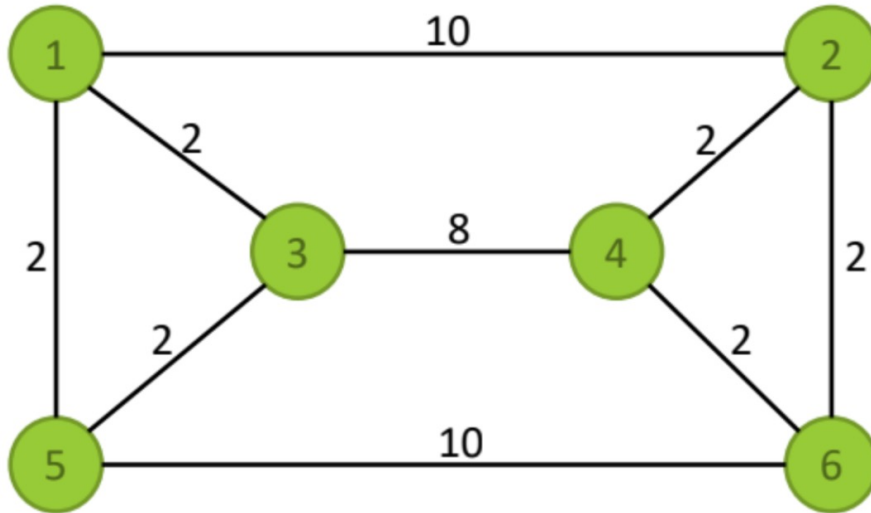


# TSP Networks



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## Symmetric TSP

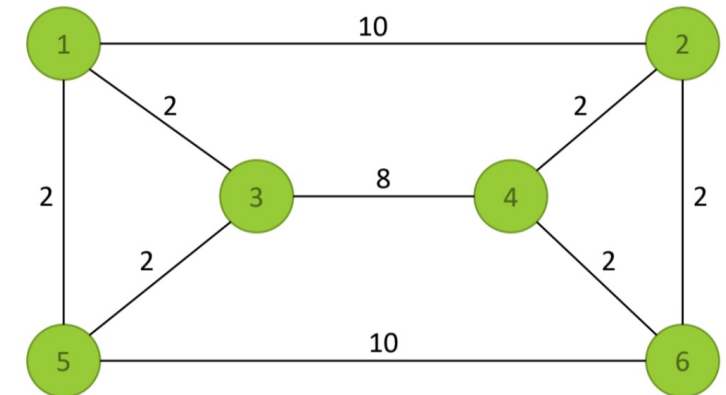
- The distance between all pairs of two cities  $(i,j)$  is the same in both direction

## Asymmetric TSP

- Distances may differ between the  $i-j$  and  $j-i$  direction
- Paths may not exist in both directions  $i-j$  and  $j-i$

# TSP: LP Formulation (Variables and Parameters)

- **What are the variables?**
  - If tour uses the arc from city  $i$  to city  $j$
  - $x_{ij} = 1$  if you passes from city  $i$  to city  $j$
  - $x_{ij} = 0$  otherwise
- **Indicate the valid range of the variables**
  - $x_{ij}$  are binary (0 or 1) variables
  - Denoted as  $x_{ij} \in \{0, 1\} \forall i, j$
- **Parameters**
  - $N$ : set of nodes
  - $d_{ij}$ : Travel distance of arc  $(i, j)$



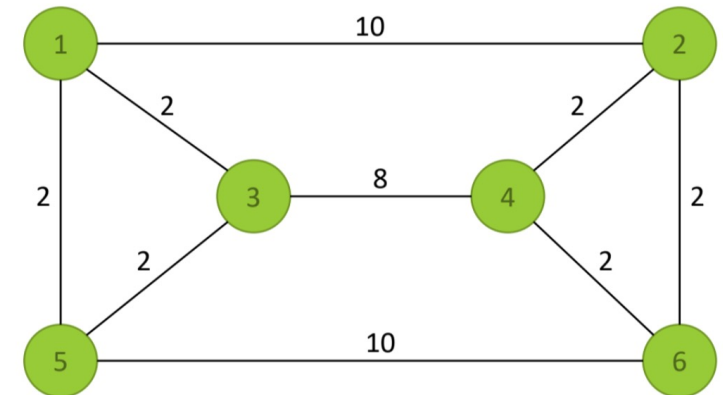
# TSP: LP Formulation (Objective FUNCTION)

- **Objective function**

- Minimize route length
- Minimize  $10x_{12} + 2x_{13} + 2x_{15} + 10x_{21} + 2x_{24} + 2x_{26} + 2x_{31} + 2x_{35} + 8x_{34} + 2x_{42} + 8x_{43} + 2x_{46} + 2x_{51} + 2x_{53} + 10x_{56} + 2x_{62} + 2x_{64} + 10x_{65}$
- Or in summation form:

$$\text{Min} \sum_{i \in N} \sum_{j \in N} d_{ij} x_{ij}$$

where  $d_{ij}$  denotes the length of the arc between city  $i$  and city  $j$



# TSP: LP Formulation (Constraints)

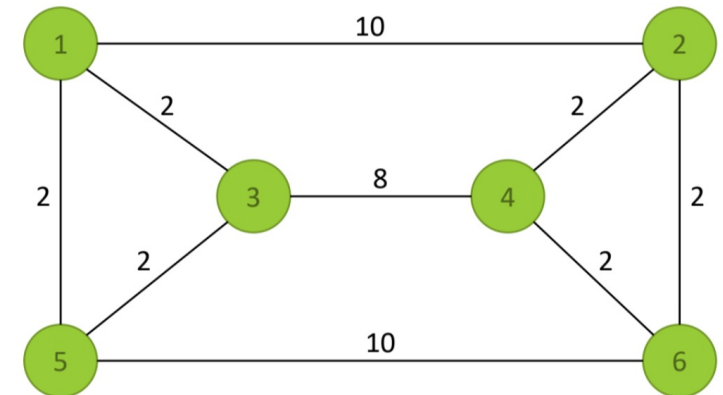
- **Formulate the constraints**

- For each city  $i$ , there is exactly **one** city that is its predecessor:

$$\sum_j x_{ji} = 1, \forall i \in N$$

- For each city  $i$ , there is exactly **one** city that is its successor:

$$\sum_j x_{ij} = 1, \forall i \in N$$



# TSP: LP Formulation (Constraints)

- Example (for the given network):

$$\sum_{j \in N} x_{ji} = 1, \forall i \in N$$

$$x_{21} + x_{31} + x_{51} = 1$$

$$x_{12} + x_{42} + x_{62} = 1$$

$$x_{13} + x_{43} + x_{53} = 1$$

$$x_{24} + x_{34} + x_{64} = 1$$

$$x_{15} + x_{35} + x_{65} = 1$$

$$x_{26} + x_{46} + x_{56} = 1$$

$$\sum_{j \in N} x_{ji} = 1, \forall i \in N$$

$$x_{12} + x_{13} + x_{15} = 1$$

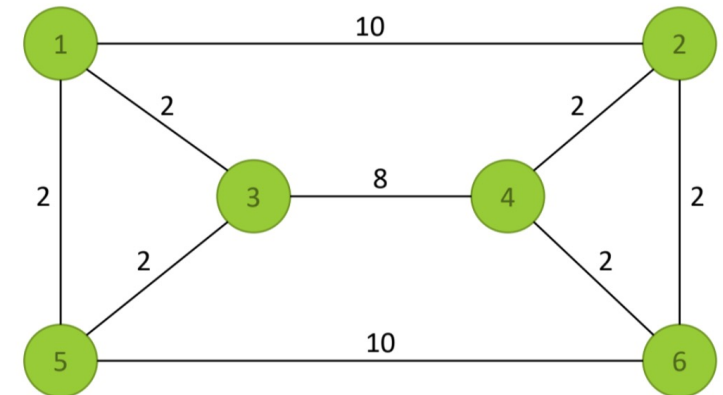
$$x_{21} + x_{24} + x_{26} = 1$$

$$x_{31} + x_{34} + x_{35} = 1$$

$$x_{42} + x_{43} + x_{46} = 1$$

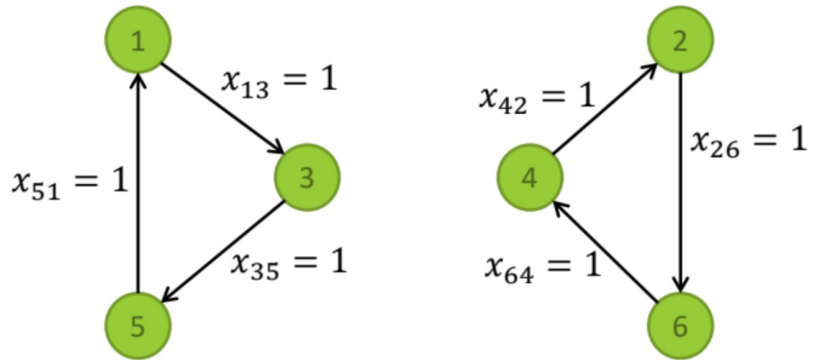
$$x_{51} + x_{53} + x_{56} = 1$$

$$x_{62} + x_{64} + x_{65} = 1$$



# TSP: LP Formulation (Constraints)

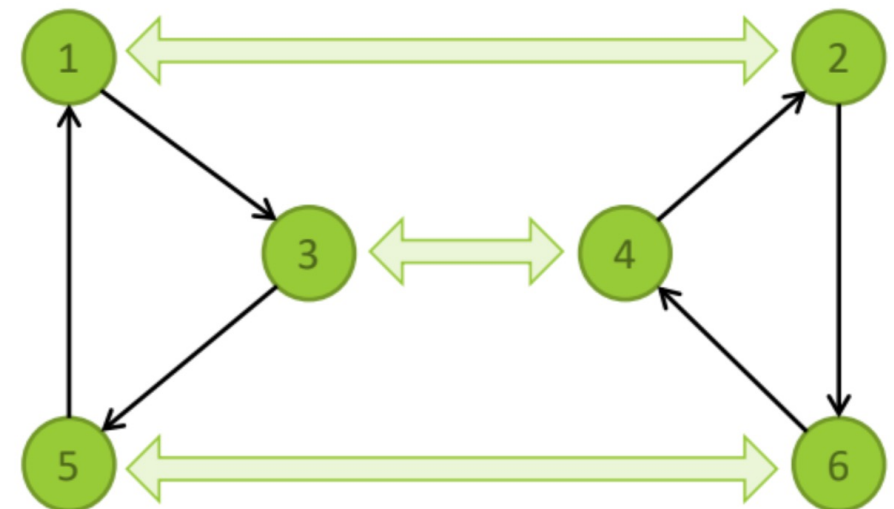
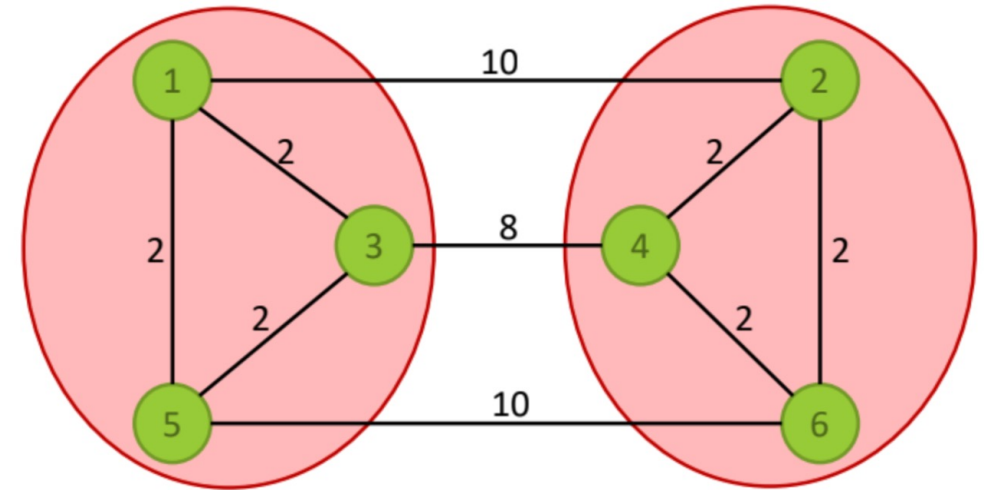
- So far, we have introduced two sets of constraints
  - Every node  $i$  must have a predecessor node
  - Every node  $i$  must have a successor node
- Now, consider this solution



- It satisfies the two sets of constraints
  - Every node has exactly one predecessor and exactly one successor
- But it is NOT a valid TSP tour!
  - **Remember:** TSP tour must be a route, i.e., the salesman must be able to start at one of the cities, visit all other cities, and come back to the city where they started.
- So, we need more constraints.

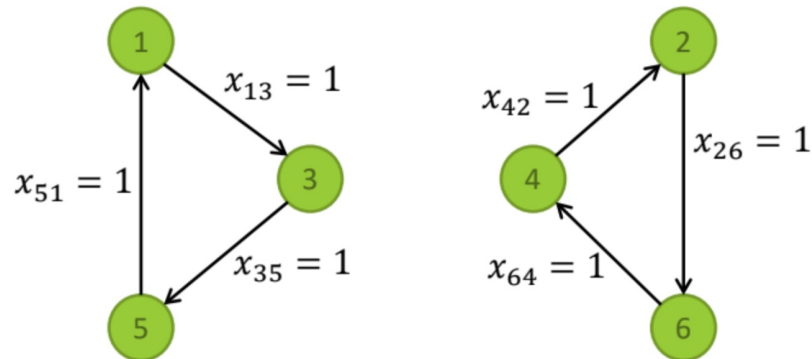
# TSP: LP Formulation (SECs)

- We need **subtour elimination constraints** (SECs) to prevent solutions that contain subtours
- At least one arc that leaves one of the nodes 1, 3, 5 and enters one of the nodes 2, 4, or 6 must be in the solution.
- In this network, only arcs 1-to-2, 3-to-4, and 5-to-6 apply
- Mathematically
  - $x_{12} + x_{34} + x_{56} \geq 1$



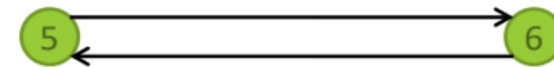
# TSP: LP Formulation (SECs)

- Once we add the constraint  $x_{12} + x_{34} + x_{56} \geq 1$  to the formulation, this solution

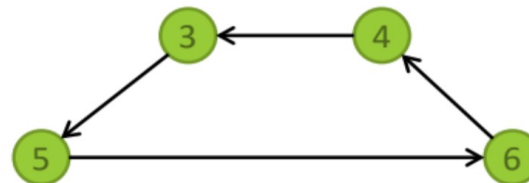
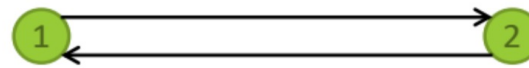


is not feasible anymore.

- But what about this one



- Or this one



- These solutions satisfy all the constraints defined so far, including  $x_{12} + x_{34} + x_{56} \geq 1$ . But they are still NOT valid TSP tours.



# TFC: LP Formulation (SECs)

- The SECs need to be defined so that they prevent **ALL POSSIBLE SUBTOURS!**
  - To guarantee, this, we need to define SECs for every node subset.
  - Remember:
    - $N$  is the set of all nodes (cities) in the network.
    - When we write  $S \subset N$ , it means that  $S$  is a **subset** of set  $N$ .
      - For example, if  $N = \{1, 2, 3, 4\}$  and  $S = \{1, 2\}$ , then  $S \subset N$ .
    - SECs defined mathematically as follows.

$$\sum_{i \in S} \sum_{j \in N \setminus S} x_{ij} \geq 1, \quad \forall S \subset N: 2 \leq |S| \leq |N| - 1$$

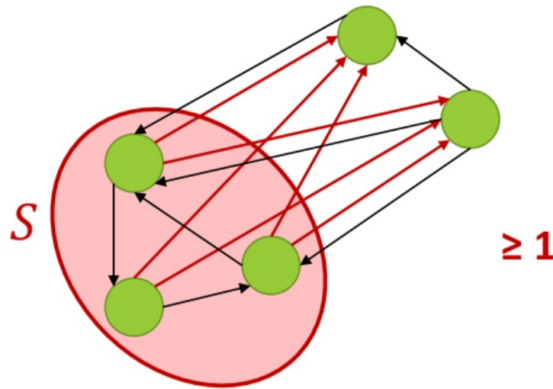
There must be **at least one** arc **leaving a node in  $S$**  and **entering a node outside of  $S$** . These constraints must be defined **for all subsets  $S$**  that contain **at least two nodes** (because a node subset with a single node cannot be a subtour) and **at most the number of nodes minus 1 node**.

# TSP: LP Formulation (SECs)

- For every node subset  $S$  such that  $2 \leq |S| \leq |N| - 1$ ,

$$\sum_{i \in S} \sum_{j \in N \setminus S} x_{ij} \geq 1$$

must hold.



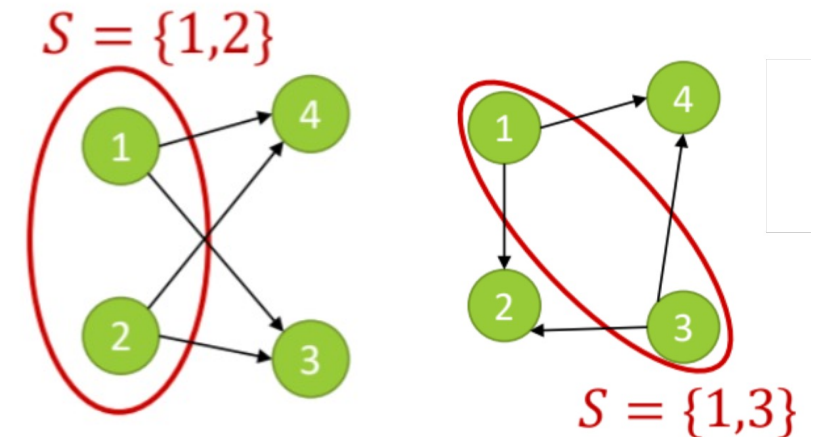
# Recall: Subsets

- **As an example, let  $N = \{1, 2, 3, 4\}$ .**
  - This set  $N$  has 4 elements, and therefore it has subsets containing 0, 1, 2, 3, and 4 elements
    - *0-element subset:  $\emptyset$  {empty set = a set with no elements in it}*
    - *1-element subsets:  $\{1\}$ ,  $\{2\}$ ,  $\{3\}$ , and  $\{4\}$*
    - *2-element subsets:  $\{1,2\}$ ,  $\{1,3\}$ ,  $\{1,4\}$ ,  $\{2,3\}$ ,  $\{2,4\}$ , and  $\{3,4\}$*
    - *3-element subsets:  $\{1,2,3\}$ ,  $\{1,2,4\}$ ,  $\{1,3,4\}$ , and  $\{2,3,4\}$*
    - *4-element subset:  $\{1,2,3,4\}$  (This is the set itself. Each set is a subset of itself.)*
  - A set with 4 elements has, in total,  $2^4 = 16$  subsets.
  - We want the subsets of  $N$  with at least 2 and at most  $|N| - 1 = 3$  elements for the SECs (specified by the expression  $\forall S \subset N: 2 \leq |S| \leq |N| - 1$ ).
    - *Thus, we want 2-element and 3-element subsets*
    - *If there were, for example, 5 elements in  $N$  (instead of 4), then we would need to define SECs for all 2-element, 3-element, and 4-element subsets of  $N$ .*

# TSP: Formulation (SECs)

$$\sum_{i \in S} \sum_{j \in N \setminus S} x_{ij} \geq 1,$$
$$\forall S \subset N: 2 \leq |S| \leq |N| - 1$$

- **Example:**  $N = \{1, 2, 3, 4\}$ 
  - Which subsets  $S$  to define these constraints?
    - $S \subset N: 2 \leq |S| \leq |N| - 1$
    - $|N| = 4$ . So, all 2-node and 3-node subsets must be considered.
  - The subsets and the corresponding SECs
    - $S = \{1,2\} \Rightarrow x_{13} + x_{14} + x_{23} + x_{24} \geq 1$
    - $S = \{1,3\} \Rightarrow x_{12} + x_{14} + x_{32} + x_{34} \geq 1$
    - Do the same for  $S = \{1,4\}$ ,  $S = \{2,3\}$ ,  $S = \{2,4\}$ , and  $S = \{3,4\}$ .
    - $S = \{1,2,3\} \Rightarrow x_{14} + x_{24} + x_{34} \geq 1$
    - $S = \{1,2,4\} \Rightarrow x_{13} + x_{23} + x_{43} \geq 1$
    - Do the same for  $S = \{1,3,4\}$  and  $S = \{2,3,4\}$ .



# Complete LP Formulation

$$\text{Min} \sum_{i \in N} \sum_{j \in N} d_{ij} x_{ij}$$

subject to

$$\sum_{j \in N} x_{ij} = 1,$$

$$\forall i \in N$$

successor constraint

$$\sum_{j \in N} x_{ji} = 1,$$

$$\forall i \in N$$

predecessor constraint

$$\sum_{i \in S} \sum_{j \in N \setminus S} x_{ij} \geq 1,$$

$$\forall S \subset N: 2 \leq |S| \leq |N| - 1$$

SECs

$$x_{ij} \in \{0, 1\},$$

$$\forall i, j \in N$$

# Classical Vehicle Routing Problem (VRP)

- **Construct a set of delivery routes** for vehicles stationed at a central depot which service all nodes and minimizes routing costs
- **From the TSP to a more general VRP**
  - Construct multiple routes/routes for multiple vehicles
  - Divide stops over routes
  - Find sequence for each route
- **Input**
  - Customers/demand is known and deterministic
  - Vehicle capacity known
  - $M$  vehicles
- **Objective function**
  - Minimize the total travel distance (time/length), and/or routing costs, and/or the number of vehicles needed

# VRP (Possible Additional Constraints)

- **Fleet**
  - Size of fleet, homogenous or heterogeneous fleet, vehicle capacity, maximum driving time
- **Nature of demand**
  - Pickups or deliveries, size of loads, due dates, spreading of locations, precedence relations
- **Underlying network**
  - Single or multiple depot(s), directed or undirected arcs, number of arcs on which vehicles can travel
- ...



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# Thank you!

# Questions?

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