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# Session 9: Facility location problems

## 35E00750 Logistics Systems and Analytics

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# Discrete and continuous methods

- **There are two types of facility location methods**
  - Methods for continuous facility location problems seek to find a location (or multiple locations) anywhere in a two-dimensional plane
  - Methods for discrete facility location problems evaluate known (candidate) locations
- **In this course, we consider the following methods**
  - Continuous: Center-of-gravity method
  - Discrete: Factor rating method
  - Discrete: Cost-volume analysis
  - Discrete: Linear programming formulation of facility location problem



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# Center-of-gravity method

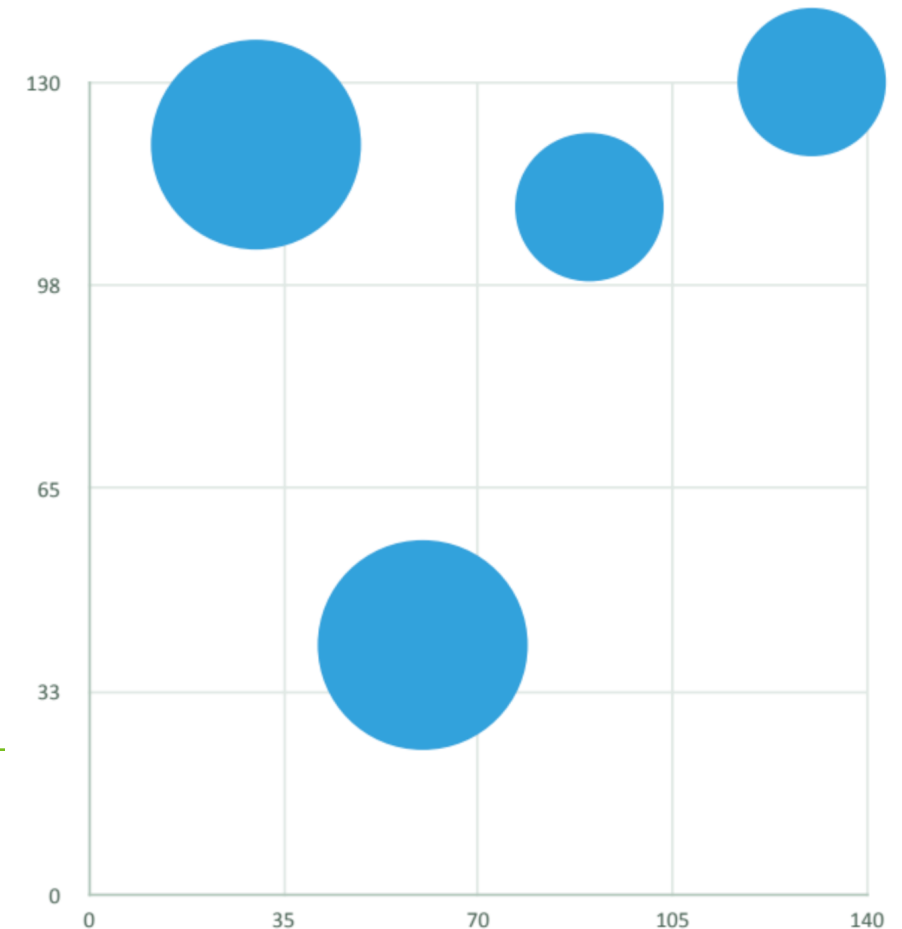
# Center-of-gravity method

- **The center-of-gravity method is a mathematical technique for finding an optimal location for a single facility**
  - Ideal location minimizes weighted (with volume of goods) distance between, for example, warehouse and retailers
- **Method is using**
  - Location of the network
  - Volume of goods to be shipped to those locations
  - (Transportation costs)
- **Objective**
  - Minimize distance/costs (directly proportional to distance and volume)

# Center-of-gravity method

- Determine an optimal location  $(x, y)$  for a warehouse location given the following demand volumes

City	Demand	Coordinates $(x,y)$
Amsterdam	2000	$(30,120)$
Berlin	1000	$(90,110)$
Osnabruck	1000	$(130,130)$
Brussels	2000	$(60,40)$



# Center-of-gravity method



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$C_x$  = X-coordinate  
 facility location

$$C_x = \frac{\sum_i d_{ix} W_i}{\sum_i W_i}$$

$d_{ix}$  = X-coordinate  
 of location  $i$

$W_i$  = Volume of goods moved  
 to or from location  $i$

$C_y$  = Y-coordinate  
 facility location

$$C_y = \frac{\sum_i d_{iy} W_i}{\sum_i W_i}$$

$d_{iy}$  = Y-coordinate  
 of location  $i$

# Center-of-gravity method (example)

- **Four cities with given coordinates and demand**

City	Demand	Coordinates (x,y)
Amsterdam	2000	(30,120)
Berlin	1000	(90,110)
Osnabruck	1000	(130,130)
Brussels	2000	(60,40)

- **Determine an optimal location (x,y) for a warehouse location**

# Center-of-gravity method (example)

- Determine an optimal location (x,y) for a warehouse location

City	Demand	Coordinates (x,y)
Amsterdam	2000	(30,120)
Berlin	1000	(90,110)
Osnabruck	1000	(130,130)
Brussels	2000	(60,40)

- Step 1: Total sum of demand ( $\sum_i w_i$ )=6000
- Step 2a: Total sum of demand moved to x-coordinate:
  - $\sum_i d_{ix}w_i = 30 * 2000 + 90 * 1000 + 130 * 1000 + 60 * 2000 = 400,000$
- Step 2b: Total sum of demand moved to y-coordinate
  - $\sum_i d_{iy}w_i = 120 * 2000 + 110 * 1000 + 130 * 1000 + 40 * 2000 = 560,000$



# Center-of-gravity method (example)

- **Determine an optimal location (x,y) for a warehouse location**
  - Step 2 (summary): 400,000 of goods to x-coordinate; 560,000 to y-coordinate
  - Step 3a: x-coordinate warehouse:

$$\frac{\sum_i d_{ix} w_i}{\sum_i w_i} = \frac{400,000}{6,000} = 66.7$$

- Step 3b: y-coordinate warehouse:

$$\frac{\sum_i d_{iy} w_i}{\sum_i w_i} = \frac{560,000}{6,000} = 93.3$$

City	Demand	Coordinates (x,y)
Amsterdam	2000	(30,120)
Berlin	1000	(90,110)
Osnabruck	1000	(130,130)
Brussels	2000	(60,40)

# Center-of-gravity method

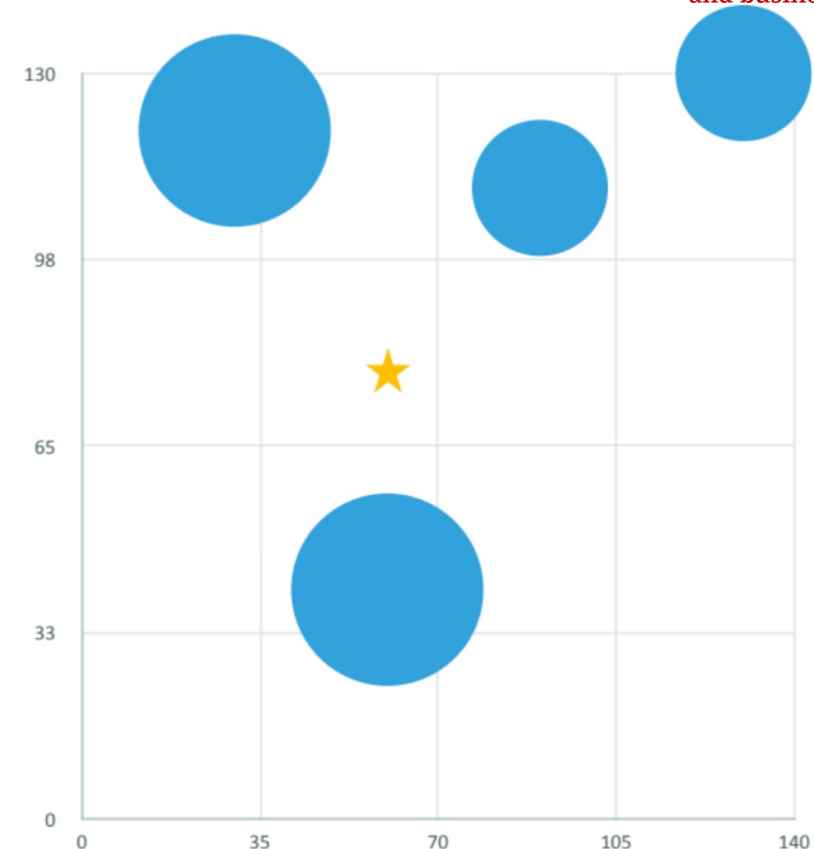


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- Optimal location is (66.7, 93.3)

City	Demand	Coordinates (x,y)
Amsterdam	2000	(30,120)
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# Factor rating method

# Factor rating method



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# Factor rating method

- **Deciding between a set of known candidate locations based on qualitative and quantitative factors**
  - List of relevant factors (qualitative)
  - Assign importance weight to each factor (0-1)
  - Develop a scale for each factor (1–100)
  - Score each location using factor scale
  - Multiply scores by weights for each factor and total
  - Select location with maximum total score
- **No exact results due to subjectivity of factors, weights, scales and scores**

# Factor rating method



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- Example inspired from Alibaba warehouse location decision

Factor	Weight	Score Netherlands	Score Belgium	Weighted score Netherlands	Weighted score Belgium
<b>Total</b>	<b>1.00</b>			<b>50</b>	<b>57</b>

# Factor rating method

- Example inspired from Alibaba warehouse location decision

Factor	Weight	Score Netherlands	Score Belgium	Weighted score Netherlands	Weighted score Belgium
Geographical positioning	0.30	60	70	$0.30 * 60 = 18$	21
Possibilities offered by airport authorities	0.25	10	70	$0.25 * 10 = 2.5$	17.5
Tax environment	0.15	80	20	$0.15 * 80 = 12$	3
State of the infrastructure	0.10	85	25	$0.10 * 85 = 8.5$	2.5
Availability of labour	0.20	45	65	$0.20 * 45 = 9$	13
<b>Total</b>	<b>1.00</b>			<b>50</b>	<b>57</b>



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# Cost-volume analysis



# Cost-volume analysis

- Analysis to make an **economic comparison** of a set of known/candidate locations
- **Determine fixed and variable costs for each known location**
  - Fixed costs (for example, costs of opening a warehouse, or costs of acquiring a truck)
  - Variable costs (for example, cost per product, or cost per kilometer)

# Cost-volume analysis (example)

- **Three candidate locations for a warehouse**
  - Assen (fixed cost of €30,000.00 and variable cost of €65.00 per pallet)
  - Heerenveen (fixed cost of €55,000.00 and variable cost of €30.00 per pallet)
  - Groningen (fixed cost of €110,000.00 and variable cost of €10.00 per pallet)
- **Which location would be best at which number of pallets?**

# Cost-volume analysis (example)

- Let  $x$  be the number of pallets
- Assen vs. Heerenveen

$$\begin{aligned} 30000 + 65x &= 55000 + 30x \\ \Leftrightarrow 35x &= 25000 \\ \Leftrightarrow x &= 714.29 \end{aligned}$$

- Hence, from 715 pallets, Heerenveen becomes the preferred location

- Heerenveen vs. Groningen

$$\begin{aligned} 55000 + 30x &= 110000 + 10x \\ \Leftrightarrow 20x &= 55000 \\ \Leftrightarrow x &= 2750 \end{aligned}$$

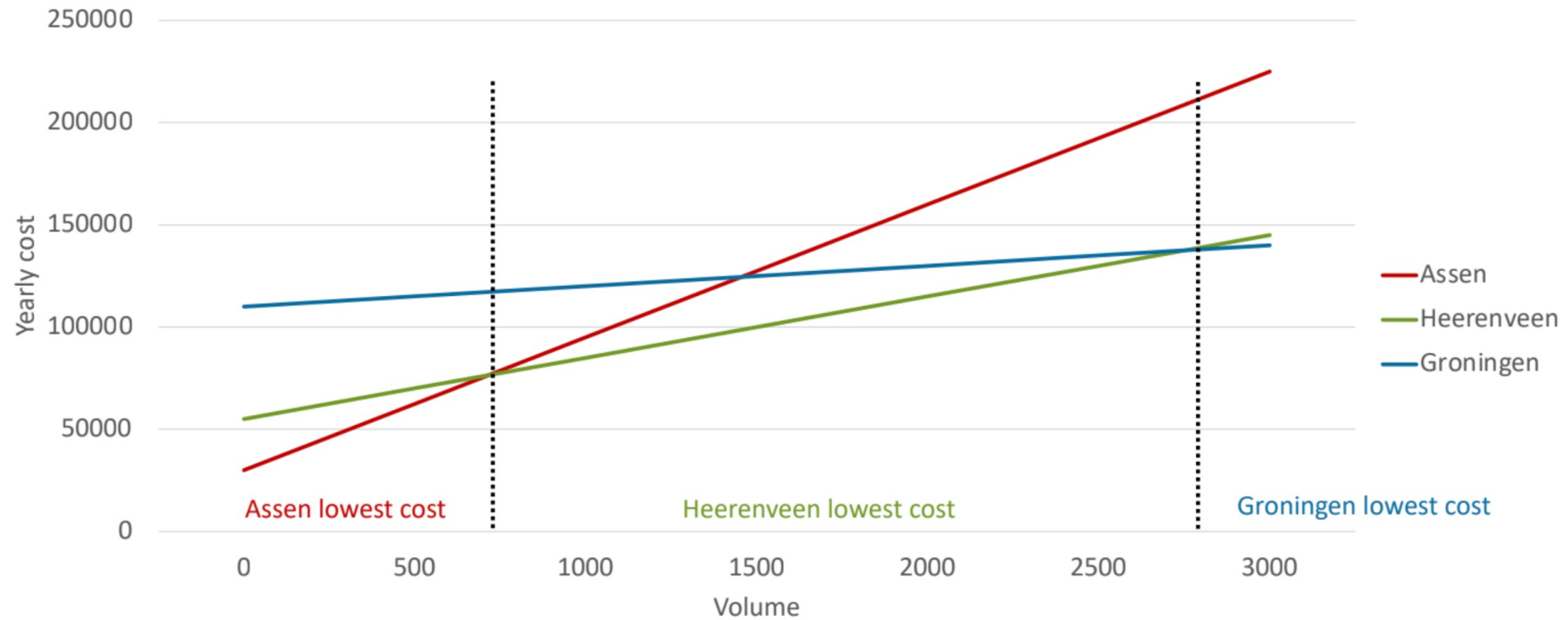
- Hence, from 2751 pallets, Groningen becomes the preferred location

# Cost-volume analysis (example)



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# Linear Programming (LP) Formulation of the Facility Location Problem (FLP)

# LP formulation of the FLP

- **General FLP**
  - Set of spatially distributed customers
  - Set of candidate facilities to serve customer demands
  - Distances, time, and costs are measured by a given metric
- **Main questions**
  - Number of facilities
  - Location(s) of facilities
- **Typical objective**
  - Minimize costs (facility costs, transportation costs, inventory costs)

# Facility Location Problem

- **Sets:**

- $I$ : set of customers
- $J$ : set of candidate locations

- **Parameters:**

- $D_i$ : demand amount of customer  $i$
- $K_j$ : capacity of facility  $j$
- $F_j$ : fixed cost for opening facility  $j$
- $c_{ij}$ : cost of sending **one unit of product** from location  $j$  to customer  $i$

- **Variables:**

- $y_j$ : whether or not to open a facility at location  $j$
- $x_{ij}$ : amount of demand of customer  $i$  satisfied from location  $j$

Minimize

$$\sum_{j \in J} F_j y_j + \sum_{i \in I} \sum_{j \in J} c_{ij} x_{ij}$$

s.t.

$$\sum_{j \in J} x_{ij} = D_i \quad \forall i \in I$$

$$\sum_{i \in I} x_{ij} \leq K_j y_j \quad \forall j \in J$$

$$x_{ij} \geq 0 \quad \forall i \in I, \forall j \in J$$

$$y_j = 0 \text{ or } 1 \quad \forall j \in J$$

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Minimize the total **fixed (facility opening) costs** and **transportation costs**



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**All demands of all customers  
 must be met**

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$$x_{ij} \geq 0 \quad \forall i \in I, \forall j \in J$$

$$y_j = 0 \text{ or } 1 \quad \forall j \in J$$

- If facility  $j$  is not opened, **no demand can be served from there**. (If  $y_j = 0$ , then  $\sum_{i \in I} x_{ij} = 0$ ).
- If facility  $j$  is opened, **the maximum amount it can supply is bounded by its capacity**. (if  $y_j = 1$ , then  $\sum_{i \in I} x_{ij} \leq K_j$ )

# Facility Location Problem

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$$\sum_{i \in I} x_{ij} \leq K_j y_j \quad \forall j \in J$$

$$\begin{aligned} x_{ij} &\geq 0 && \forall i \in I, \forall j \in J \\ y_j &= 0 \text{ or } 1 && \forall j \in J \end{aligned}$$

Valid range constraints

# Facility Location Problem (Uncapacitated)

- If facilities have infinite (or unrestrictively large) capacities, the constraint

$$\sum_{i \in I} x_{ij} \leq K_j y_j \quad \forall j \in J$$

needs to be adjusted.

- Note that completely removing this constraint is not correct!
  - We still need to ensure that if facility  $j$  is not opened, **no demand can be served from there** (if  $y_j = 0$ , then  $\sum_{i \in I} x_{ij} = 0$ ).
- The following constraint achieves this (with  $M$  being a sufficiently large number):

$$\sum_{i \in I} x_{ij} \leq M y_j \quad \forall j \in J$$

- If  $y_j = 0$ , then  $\sum_{i \in I} x_{ij} = 0$
- If  $y_j = 1$ , then  $\sum_{i \in I} x_{ij} \leq M$  (not restrictive, since  $M$  is a large number)



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# Network Design Problem

# Different Network Designs



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- 1 fresh food distribution center
- 1 non-perishable goods distribution center
- 4 regional distribution centers



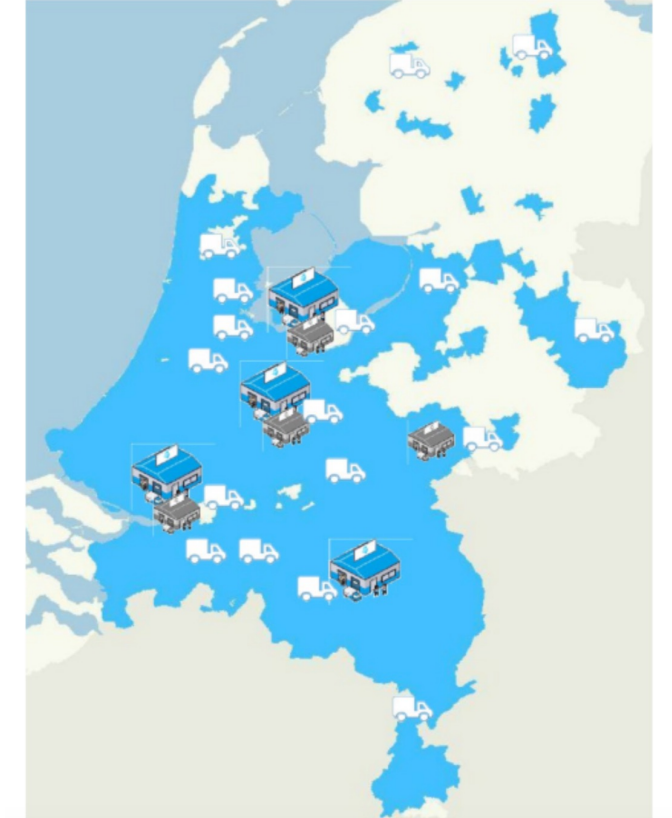
# Different Network Designs



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- 4 home shopping centers
- 4 national centers (non-perishable, frozen fresh, bake-off)
- 17 hubs





# General Approach to Network Design

- **Given**
  - A set of facilities, and
  - demand/supply quantities of these facilities
- **Find**
  - The routes to be operated
  - The features (frequency, number of intermediate stops, etc.) of the routes to be operated
  - The traffic assignment along the routes
  - The operating rules at each facility
  - Possibly, the relocation of empty vehicles and containers

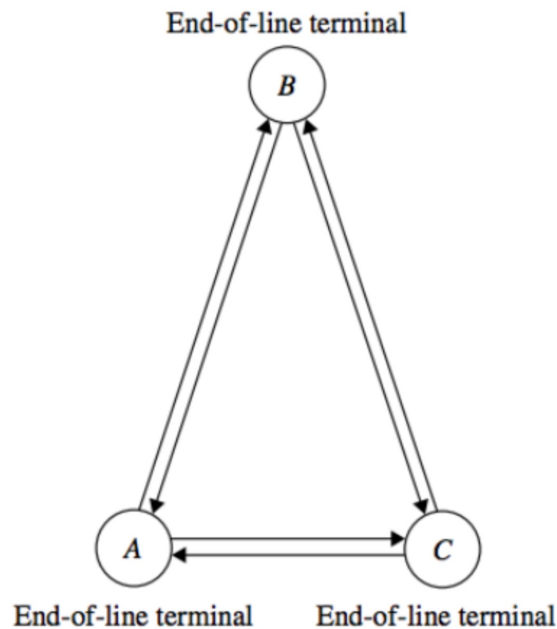
# Network Design Problem



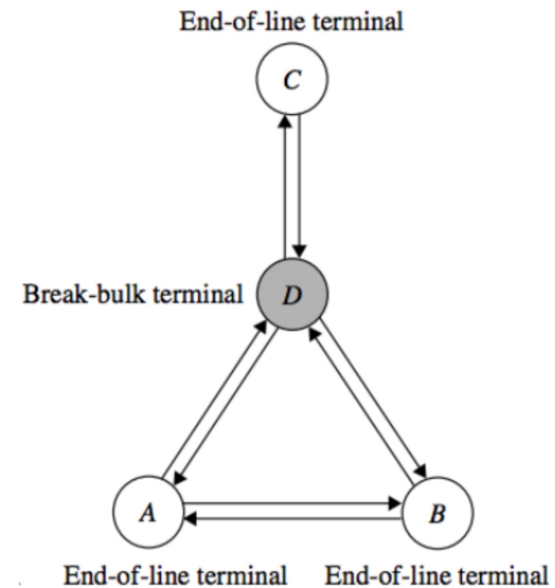
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- Consider two alternative service networks



Quicker shipments but higher operating cost



Freight consolidation: slower shipments but lower operating cost

# Basic Network Design Problem

- **Single commodity**
- **Known demand/supply quantities of each facility (node)**
- **Known costs for each link between facilities (arc)**
  - Fixed cost for opening an arc
  - (Variable cost) transportation cost over an arc
- **Find**
  - Whether or not to open each arc
  - The volume of goods transported on each arc

# Step 1: Definition of Parameters

- **Sets**
  - $D$ : Set of demand nodes
  - $S$ : Set of supply nodes
  - $N$ : Set of all nodes ( $D \cup S$ )
- **Parameters**
  - $D_i$ : demand quantity of node  $i$
  - $S_i$ : supply quantity of node  $i$
  - $u_{ij}$ : capacity of arc  $(i, j)$
  - $f_{ij}$ : fixed cost of opening arc  $(i, j)$
  - $c_{ij}$ : cost of transporting **one unit of product** on arc  $(i, j)$

# Step 1: Definition of Variables

- **Step 1a: What are the variables?**

- $y_{ij} = \begin{cases} 1, & \text{if arc } (i, j) \text{ is opened} \\ 0, & \text{otherwise} \end{cases}$
- $x_{ij}$ : quantity transported on arc  $(i, j)$

- **Step 1b: Indicate the valid range of all variables**

- $y_{ij} \in \{0,1\}, \forall i \in N, \forall j \in N$  (binary:  $y_{ij}$  values are 0 or 1 for all  $i$  and  $j$ )
- $x_{ij} \geq 0, \forall i \in N, \forall j \in N$  ( $x_{ij}$  non-negative values for all  $i$  and  $j$ )

# Step 2: Define Objective

- **Step 2a: What do you want to achieve?**

- Minimize total cost, consisting of fixed cost (cost to open a link) and variable costs (per unit transportation costs)

- **Step 2b: Express mathematically**

- Fixed costs

$$\sum_{i \in N} \sum_{j \in N} f_{ij} y_{ij}$$

- Variable costs

$$\sum_{i \in N} \sum_{j \in N} c_{ij} x_{ij}$$

- Objective function:

$$\text{Minimize } \sum_{i \in N} \sum_{j \in N} f_{ij} y_{ij} + \sum_{i \in N} \sum_{j \in N} c_{ij} x_{ij}$$

# Step 3: Formulating Constraints

- **Balance constraints:**
  - Net flow out of a supply node is equal to the supply quantity
  - Net flow into a demand node is equal to the demand quantity

Total volume of  
 products that go  
 out of node  $i$

$$\sum_{j \in N} x_{ij} - \sum_{j \in N} x_{ji} = \begin{cases} S_i & \text{if } i \in S \\ -D_i & \text{if } i \in D \end{cases}$$

Total volume of  
 products that  
 come into node  $i$

$\forall i \in N$

Constraint defined for every  
 node  $i$  in the set of nodes  $N$

# Step 3: Formulating Constraints

- **Constraints linking the two variables,  $x_{ij}$  and  $y_{ij}$** 
  - If  $y_{ij} = 0$ ,  $x_{ij}$  must be zero as well
  - If  $y_{ij} = 1$ ,  $x_{ij}$  can take any value between 0 and the arc capacity ( $u_{ij}$ )

$$x_{ij} \leq u_{ij}y_{ij}, \quad \forall i \in N, \forall j \in N$$



# Complete Formulation



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$$\text{minimize } \sum_{i \in N} \sum_{j \in N} f_{ij} y_{ij} + \sum_{i \in N} \sum_{j \in N} c_{ij} x_{ij}$$

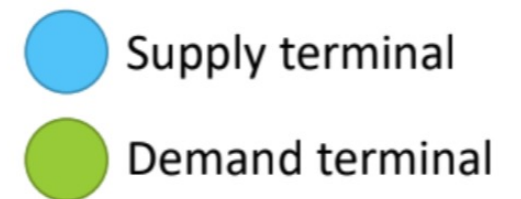
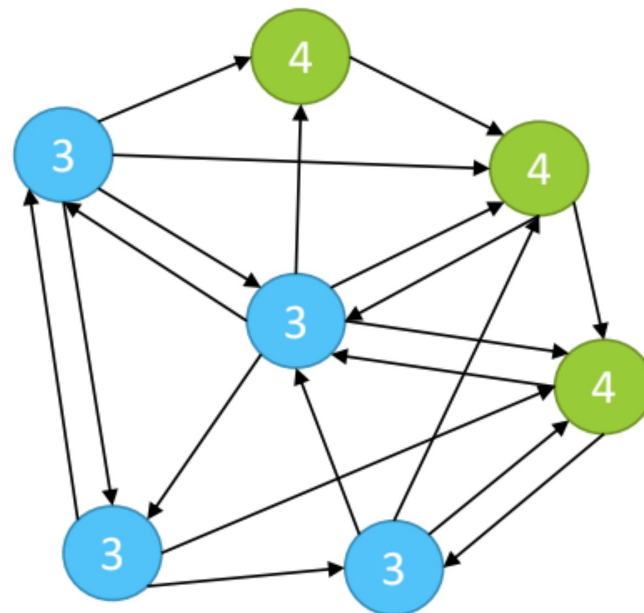
s.t.

$$\sum_{j \in N} x_{ij} - \sum_{j \in N} x_{ji} = \begin{cases} S_i & \text{if } i \in S \\ -D_i & \text{if } i \in D \end{cases} \quad \forall i \in N$$
$$x_{ij} \leq u_{ij} y_{ij}, \quad \forall i \in N, \forall j \in N$$
$$y_{ij} \in \{0,1\}, \quad \forall i \in N, \forall j \in N$$
$$x_{ij} \geq 0, \quad \forall i \in N, \forall j \in N$$

# Basic Network Design Problem

- **Example: 7-node network**

- Supply nodes: demanding 4 unit loads each
- Supply nodes: supplying 3 unit loads each
- There are arcs from every node to every other node (not all of them are drawn on the figure)
- Which arcs to use? How many unit loads to send on each arc?



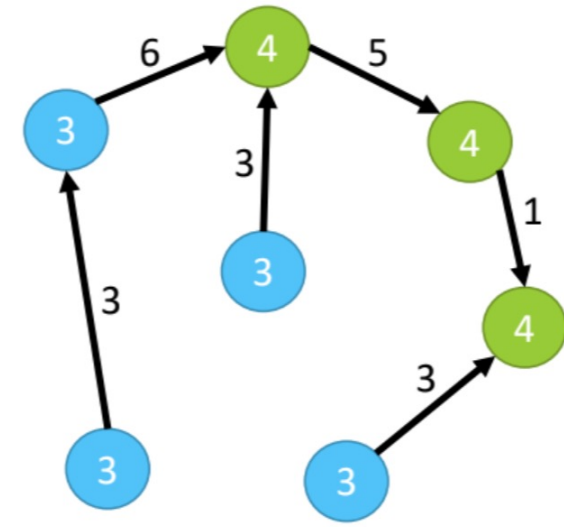
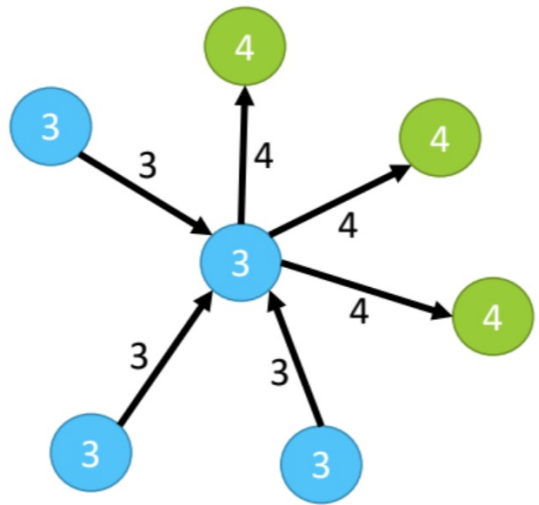
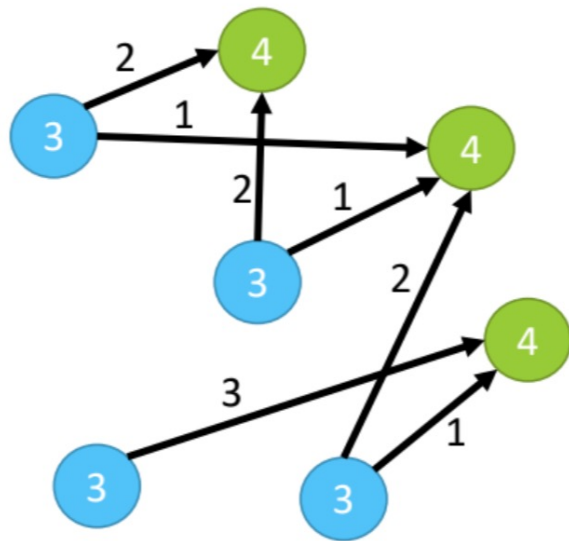
# Basic Network Design Problem



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- **Some of the possible solutions**
  - Note that there are many, many more



# Network Design Problem (Uncapacitated)

- **What if some (or all) arcs were uncapacitated)**
  - “Uncapacitated” means infinite capacity
  - Instead of  $x_{ij} \leq u_{ij}y_{ij}$ , we would get
  - $x_{ij} \leq My_{ij}$   
(where M is a sufficiently large number)

# Network Design Problem (Multi-commodity)

- **Multi-commodity network design problem variant**
- **Given demand/supply quantities of each node for each commodity, and costs**
  - Fixed cost for opening an arc
  - (Variable cost) Cost of transporting each commodity over an arc
- **Find**
  - Whether or not to open an arc
  - The volume of **each commodity** transported on each arc

# Network Design Problem (Multi-commodity)



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- **Decision variables**

- $y = \begin{cases} 1, & \text{if arc}(i, j) \text{ is opened} \\ 0, & \text{otherwise} \end{cases}$
- $x_{ij}^k$ : quantity of commodity  $k$  transported on arc  $(i, j)$

- **Parameters**

- $K$ : Set of commodities
- $D(k)$ : Set of nodes that demand commodity  $k$
- $S(k)$ : Set of nodes that supply commodity  $k$
- $N$ : Set of all nodes ( $N = \bigcup_{k \in K} [S(k) \cup D(k)]$ )
- $D_i^k$ : Demand quantity of node  $i$  for commodity  $k$
- $S_i^k$ : Supply quantity of commodity  $k$  in node  $i$
- $u_{ij}$ : Capacity of arc  $(i, j)$
- $f_{ij}$ : Fixed cost of opening arc  $(i, j)$
- $c_{ij}^k$ : Unit cost of transporting commodity  $k$  on arc  $(i, j)$

# Network Design Problem (Multi-commodity)

- Complete formulation for multi-commodity capacitated

$$\text{minimize } \sum_{i \in N} \sum_{j \in N} f_{ij} y_{ij} + \sum_{i \in N} \sum_{j \in N} \sum_{k \in K} c_{ij}^k x_{ij}^k$$

s.t.

$$\sum_{j \in N} x_{ij}^k - \sum_{j \in N} x_{ji}^k = \begin{cases} S_i & \text{if } i \in S(k) \\ -D_i & \text{if } i \in D(k) \end{cases}, \quad \forall i \in N, \forall k \in K$$

$$\sum_{k \in K} x_{ij}^k \leq u_{ij} y_{ij}, \quad \forall i \in N, \forall j \in N$$

$$\sum_{k \in K} x_{ij}^k \leq u_{ij} y_{ij}, \quad \forall i \in N, \forall j \in N$$

$$y_{ij} \in \{0,1\}, \quad \forall i \in N, \forall j \in N$$

$$x_{ij}^k \geq 0, \quad \forall i \in N, \forall j \in N, \forall k \in K$$

Thank you!

Questions?

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