# Week IV

### **Overview**

- Matchings;
- Maximal vs Maximum;
- Matching and Flow.

#### **Definitions**

- $M \subset E$  is called *matching* if all  $e \in M$  are pairwise disjoint, i.e., if the endpoints are all different
- M ⊂ E is a maximum matching in G if M is a matching with highest cardinality, i.e.,

 $|M'| \le |M|$  for all matchings M'

# **Definition (Maximal Matching)**

**No** other edge can be added to this match;

# **Definition (Maximum Matching)**

A match with the largest possible number of edges;

#### **Definition (Maximum vs Maximal)**

A maximum matching is the maximal matching with the maximum number of edges.

Maximum Matching	Maximu
Algorithm 1: MAXIMUM MATCHING	
<b>Input:</b> undirected graph $G = (V, E)$	
<b>Output:</b> maximum matching <i>M</i>	max
$1 \text{ set } M = \emptyset$	
<sup>2</sup> while there exists <i>M</i> -augmenting path in <i>G</i> do	s.t.
3 choose <i>M</i> -augmenting path <i>P</i>	
4 set $M = (M \setminus E(P)) \cup (E(P) \setminus M)$	
5 return <i>M</i>	
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Maximum Matching - Bipartite	
Algorithm 2: MAXIMUM MATCHING BIPARTITE	
Graphs	
<b>Input:</b> undirected bipartite graph $G = (V, E)$	
<b>Output:</b> maximum matching M	
1 set $M = \emptyset$	
2 construct G'	
3 while there exists s-t-path in G' do	
4 choose <i>s</i> - <i>t</i> -path <i>P</i>	
s set $M = (M \setminus E(P)) \cup (E(P) \setminus M)$	

- 6 update G'
- 7 return M

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# um Matching ILP





