

## Overview

- Enumeration;
- Dynamic Programming.

## Enumeration

**Enumeration** is the process that **list** of all the solutions and **test** each one individually.

## Dynamic Programming

**Dynamic programming** is a solution method by breaking them into a collection of simpler **sub-problems**, solving them once and storing their solutions.

It requires two properties:

- Recursion;
- Suboptimal Structure.

## Knapsack Problem

```

1  $m_{0,weight} \leftarrow 0$ 
2 for each cell  $m_{item,weight}$  do
3    $m_{item,weight} \leftarrow m_{item-1,weight}$  if
      $weight_{item} > weight$ 
4    $m_{item,weight} \leftarrow$ 
      $\max(m_{item-1,weight}, m_{item-1,w-weight_{item}} + v_i)$  if
      $weight_{item} \leq weight$ 
    
```

## Tower of Hanoi

### Algorithm 1: HANOI(BFS)

**Input:** Disk, Source, Destination, Auxiliary Rod

```

1 if  $Disk == 1$  then
2    $\text{move Disk from Source to Destination};$ 
3 else
4    $Hanoi(Disk-1, Source, Auxiliary Rod,$ 
      $Destination)$ 
5    $\text{move Disk from Source to Destination};$ 
6    $Hanoi(Disk-1, Auxiliary Rod, Destination,$ 
      $Source)$ 
    
```

## Bellman-Ford Algorithm

### Algorithm 2: BELLMAN-FORD'S ALGORITHM

**Input:** undirected, connected graph  $G$ , weights  
 $c: E(G) \rightarrow \mathbb{R}$ , nodes  $V$ , source  $s$

```

1  $d_v$  : distance to reach node  $v$ 
2  $p_v$  : node predecessor to node  $v$ 
3  $Q \leftarrow \emptyset$  : set of "unknown distance" nodes.
4 for each node  $v$  in  $V$  do
5    $d_v \leftarrow \infty$ 
6    $p_v \leftarrow FALSE$ 
7  $d_s \leftarrow 0$ 
8 for each node  $v$  in  $V$  do
9   for each edge  $(u, v) \in E$  do
10     $temp-dist \leftarrow d_u + c_{uv}$ 
11    if  $temp-dist < d_v$  then
12       $d_v \leftarrow temp-dist$ 
13       $p_v \leftarrow u$ 
14   for each edge  $(u, v) \in G$  do
15     if  $d_u + c_{uv} < d_v$  then
16       return Error: Negative Cycle Exist
17 return  $d_v, p_v$ 
    
```