

COMBINATORIAL OPTIMIZATION

Dias, Schiewea, Pattanaki *Instructor: Pattanaki*

Introduction to Combinatorics

§ Week I §

Problem 1: All representations

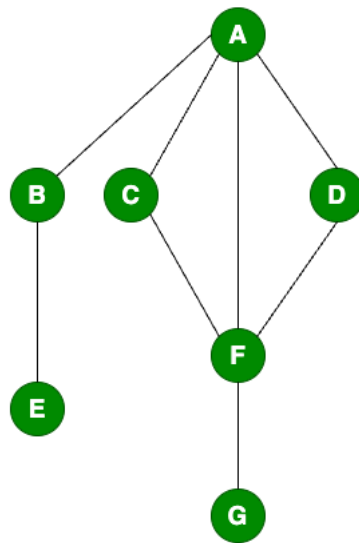


Figure 1: Example of an undirected graph.

Represent the graph above using **incidence matrix** and **list adjacency**.

Solution:

Incidence matrix:

$$\begin{pmatrix} & A & B & C & D & E & F & G \\ A & 0 & 1 & 1 & 1 & 0 & 1 & 0 \\ B & 1 & 0 & 0 & 0 & 1 & 0 & 0 \\ C & 1 & 0 & 0 & 0 & 0 & 1 & 0 \\ D & 1 & 0 & 0 & 0 & 0 & 1 & 0 \\ E & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ F & 1 & 0 & 1 & 1 & 0 & 0 & 1 \\ G & 0 & 0 & 0 & 0 & 0 & 1 & 0 \end{pmatrix}$$

Adjacency list:

$A : \{B, C, D, F\}$
 $B : \{A, E\}$
 $C : \{A, F\}$
 $D : \{A, F\}$
 $E : \{B\}$
 $F : \{A, C, D, G\}$
 $G : \{F\}$

Problem 2: DFS by Hand

Considering the graph from the previous problem, write up all the steps required to complete the **DFS** algorithm.

Solution:

We have a selection of nodes to follow to build into another data structure S :

- Starting from node A , we have four options: B , C , D and F . Then, add A in S ($S = \{A\}$);
- Picking B and put into S ($S = \{A, B\}$);
- From B , the next option to go is E which we would add to S ($S = \{A, B, E\}$);
- E is a dead-end, so it is returning to B , which does not have any other alternative to go, so it goes back to A ;
- From A , it goes to C ($S = \{A, B, E, C\}$) then goes to F ($S = \{A, B, E, C, F\}$);
- From F , it goes to G , which is a dead-end, so it returns to F ($S = \{A, B, E, C, F, G\}$);
- From F , the only available route is D resulting in $S = \{A, B, E, C, F, G, D\}$

Note that there is no single solution, and different DFS are possible.

Problem 3: Escape Room

Using any graph representation of your choice, express the following **maze** as a **graph**:

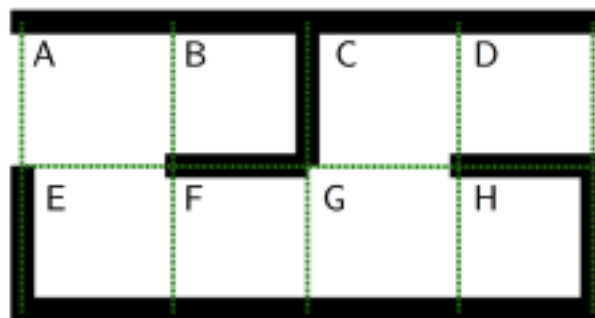


Figure 2: Search for the keys

Solution:

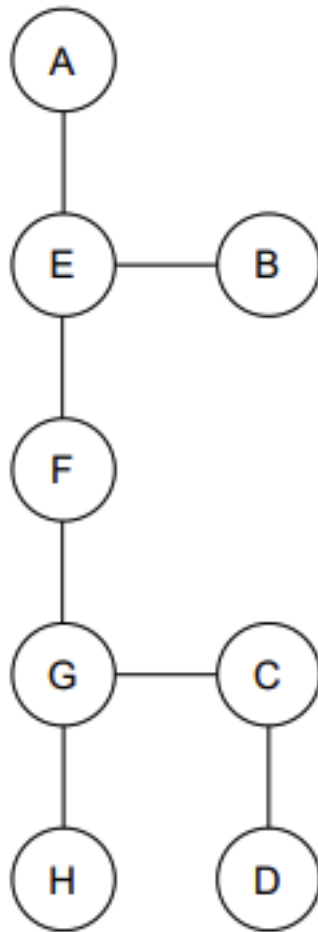


Figure 3: Graph from maze

Using **visual representation**:

Using **incidence matrix**:

$$\begin{pmatrix} & A & B & C & D & E & F & G & H \\ A & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\ B & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ C & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 \\ D & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ E & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ F & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 \\ G & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 \\ H & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \end{pmatrix} \text{ and}$$

Using **adjacency list**:

$A : \{B, E\}$
 $B : \{A\}$
 $C : \{D, G\}$
 $D : \{C\}$
 $E : \{A, F\}$
 $F : \{E, G\}$
 $G : \{C, H\}$
 $H : \{G\}$

Problem 4: Escaping

Using DFS, find at least **two different** routes to reach **room H**.

Solution:

First, using DFS:

1. Starting for A , we can go to B ($S = \{A, B\}$);
2. B is a dead-end; then we return to A and take E instead ($S = \{A, B, E\}$);
3. From E , it goes to G , where it has two options: first, let us take C ($S = \{A, B, E, F, G, C\}$);
4. From C , it goes to D and then reaches a dead-end ($S = \{A, B, E, F, G, C, D\}$);
5. Returns to G , the only option available is H ($S = \{A, B, E, F, G, C, D, H\}$).

The two ways to reach H could be from A ($A \rightarrow E \rightarrow F \rightarrow G \rightarrow H$) and from D ($D \rightarrow C \rightarrow G \rightarrow H$).

Problem 5: DFS vs Reachability

In a sufficient large graph, a **DFS algorithm** showed that such a graph is disconnected. How would you **justified** that?

Solution: If the size of *DFS* is the same as the size of the original graph G , all nodes in the DFS are connected in the original graph.

Problem 6: Building blocks organically

The hydrocarbons known as alkanes have chemical formula C_pH_{2p+2} , alkanes where C and H represent atoms of carbon and hydrogen, respectively. Graphs can represent alkane molecules. Draw a figure of a methane C_1H_4 molecule. How many “different” C_3H_8 molecules are there?

Solution:

If you know the chemical properties of alkanes, many different molecules are available. However, those graphs would be connected and have the same DFS.

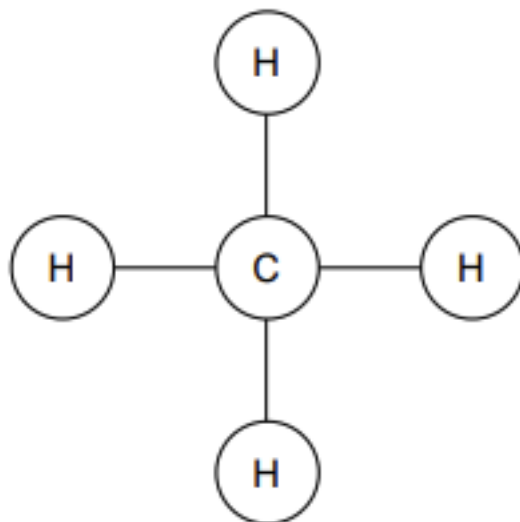


Figure 4: Methane

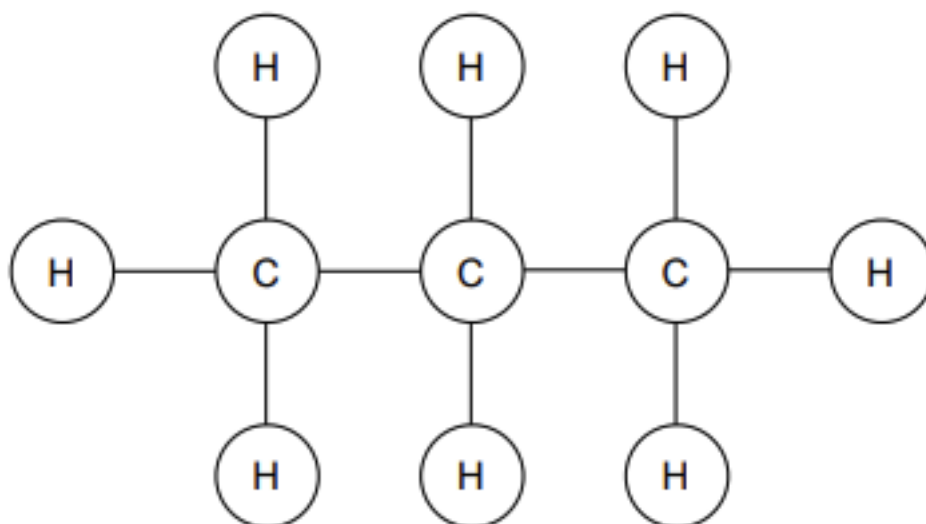


Figure 5: Propane