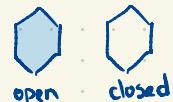
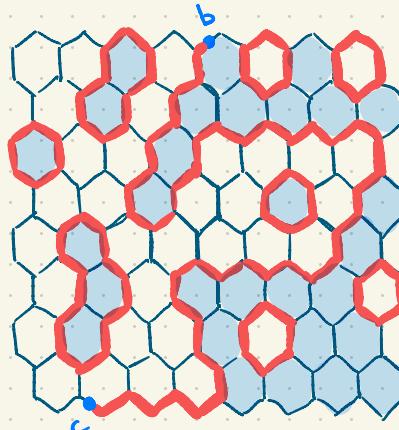


① a

- let $\phi: \Omega_G \rightarrow \Omega_{G,bc}^{\text{loop}}$

as $e \in \phi(\omega)$

iff $\begin{cases} \omega \text{ differs either side of } e \\ \omega = 1 \text{ on the } ! \text{ face by } e \\ \omega = 0 \text{ on the } ! \text{ face by } e \end{cases}$



$e \in \text{bulk}(G)$

$e \in (b,c) \subset \partial G$

$e \in (c,b) \subset \partial G$

- $\phi(\omega)$ has a self-avoiding path $b \leftrightarrow c$ given by
 $\Gamma =$ interface between open clusters touching (c,b)
 and closed clusters touching (b,c)

indeed, any path from (b,c) to (c,b) crosses an odd number of edges $e \in \phi(\omega)$ by defn of ϕ

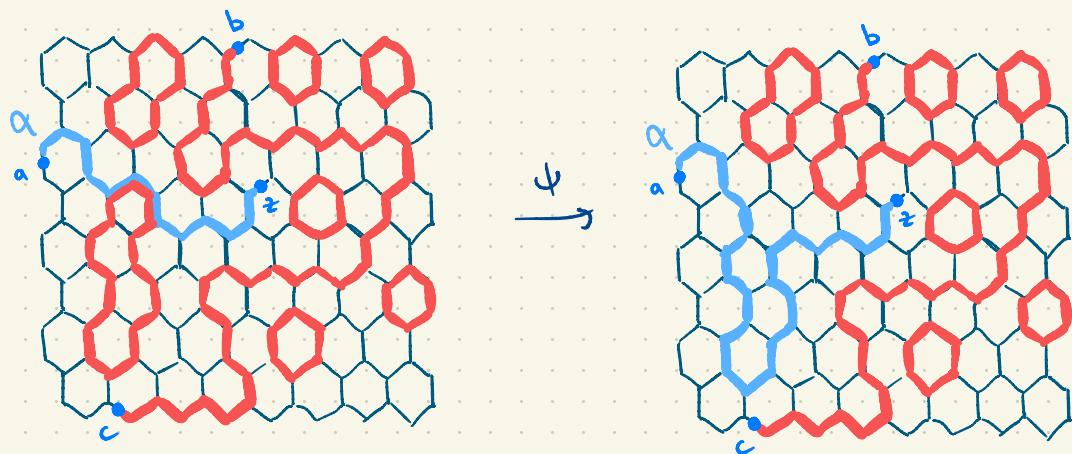
(b) if γ_0 disconnects a from z , then there's no self-avoiding path $a \leftrightarrow z$ which is disjoint from γ_0 , so in this case,

$$\mathcal{N}_{G,az,bc}^{\text{loop}}[\gamma_0] = \emptyset.$$

in the case that γ_0 does not disconnect a from z , let α be some fixed path $a \leftrightarrow z$ avoiding γ_0 . then the map

$$\psi: \mathcal{N}_{G,bc}^{\text{loop}}[\gamma_0] \rightarrow \mathcal{N}_{G,az,bc}^{\text{loop}}[\gamma_0]$$

given by $\psi(\gamma) = \gamma \Delta \alpha$ is a bijection with inverse $\psi^{-1} = \psi$. ↑ symmetric difference



c) using the above,

$$F_a(z) = \frac{|\Omega_{q,a,z,b,c}^{\text{loop}}|}{|\Omega_q|}$$

Part (b)

$$= \frac{1}{|\Omega_q|} \sum_{w_0, b \in c} |\Omega_{q,b,c}^{\text{loop}}[w_0]|$$

doesn't disconnect
a and z

Part (a)

$$= \frac{1}{|\Omega_q|} \sum_{w_0, b \in c} |\{w \in \Omega_q : \Gamma(w) = w_0\}|$$

doesn't disconnect
a and z

$$= \frac{1}{|\Omega_q|} |\{w \in \Omega_q : \Gamma(w) \text{ doesn't disconnect } a \text{ and } z\}|$$
$$= P_{\frac{1}{2}} [\Gamma(w) \text{ doesn't disconnect } a \text{ and } z]$$

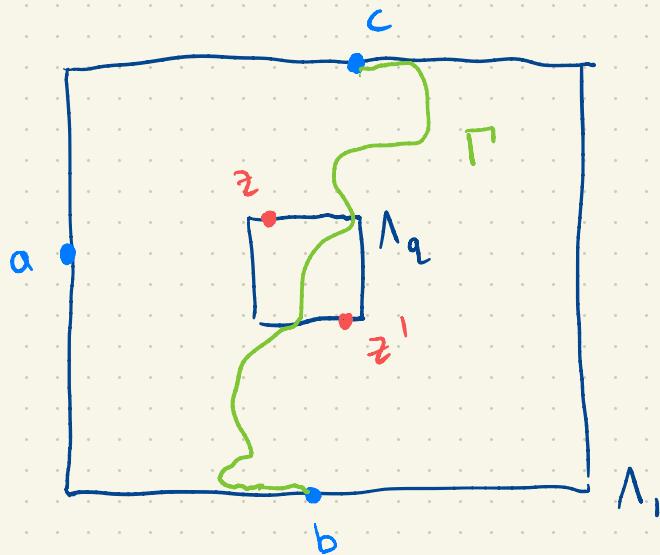
(d) by above,

$$|F_a(z) - F_a(z')| = \left| P_{\frac{1}{2}} \left[\Gamma \text{ disconnects } z \text{ from } a \right] \right| \\ - \left| P_{\frac{1}{2}} \left[\Gamma \text{ disconnects } z' \text{ from } a \right] \right|$$

$$\text{for sets } A, B, \left| |A| - |B| \right| = \left| |A \setminus B| - |B \setminus A| \right| \\ \leq \left| (A \setminus B) \cup (B \setminus A) \right| \\ = |A \Delta B|$$

so above is $\leq P_{\frac{1}{2}} \left[\Gamma \text{ disconnects exactly one of } z, z' \text{ from } a \right]$

e



Γ separates exactly one of z, z' from a

$\Rightarrow \Gamma$ intersects Λ_q

$\Rightarrow \exists$ open cluster $\Lambda_q \leftrightarrow \Lambda_1$

$\Rightarrow |F_a(z) - F_a(z')| \leq c \cdot q^\alpha \leq c \cdot |z-z'|^\alpha$.

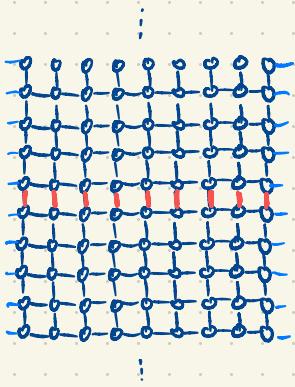
- ② a) • colour the faces of \mathbb{Z}^2 black or white according to a chessboard pattern, ie:
face f is black iff its lower left corner (x,y) has $x+y$ even.

- then the black faces form a copy of \mathbb{Z}^2 , where the black faces are the vertices & are connected by an edge iff they share a corner. call this graph \mathbb{L} .
- the white faces similarly are a copy of \mathbb{Z}^2 , and in fact it's the dual of \mathbb{L} , \mathbb{L}^* .
- a mirror is an edge of \mathbb{L} if it joins two black faces & an edge of \mathbb{L}^* o/w.
- when $p=1$, every vertex of \mathbb{Z}^2 is assigned a mirror, with its orientation chosen independently & uniformly
so each vertex is assigned either an edge of \mathbb{L} or an edge of \mathbb{L}^* .
- let $w =$ the collection of edges of \mathbb{L} assigned.

it is straightforward that $w \in P_{\frac{1}{2}}$, and so
the edges of \mathbb{L}^* assigned are w^* , with $w^* \in P_{\frac{1}{2}}$ too.

- when $p=1$, the loops in the model of mirrors separate primal and dual clusters in w, w^* ;
that is, each loop (properly oriented) has a cluster
of w on its left and a cluster of w^* on its right.
- $\exists \infty$ loop, then $\exists \infty$ primal cluster on its
left & an ∞ dual cluster on its right.
so $P_1[0 \leftrightarrow \infty] \leq P_{\frac{1}{2}}[\infty \text{ cluster in } w, \infty \text{ cluster in } w^*] = 0$
by zhang's argument

(b) consider the $2n+1$ vertical edges of $\mathbb{Z}_{2n+1} \times \mathbb{Z}$
 $\{(x, 0), (x, 1) : 1 \leq x \leq 2n+1\}$. (shown in red below).



every finite loop (and so path connecting $+\infty \leftrightarrow +\infty$ or $-\infty \leftrightarrow -\infty$) must cross an even number of these edges

\Rightarrow the total number on finite loops (and these ∞ paths) must be even

$\Rightarrow \exists$ at least one on an infinite path top \leftrightarrow bottom.

(c) one of the red edges e_0 above is incident to 0 .

by translation invariance, $P[e \leftrightarrow \infty] = P[e' \leftrightarrow \infty]$ for all e, e' red edges. so

$$\begin{aligned} 1 &= P_p[\text{at least one } e \leftrightarrow \infty] = P_p[\bigcup_{e \text{ red}} e \leftrightarrow \infty] \\ &\leq \underset{\text{union bound}}{\sum} P_p[e_0 \leftrightarrow \infty] \leq 2n+1 P_p[0 \leftrightarrow \infty] \end{aligned}$$

now $0 \leftrightarrow \infty \Rightarrow 0 \leftrightarrow \partial \Lambda_n$ so

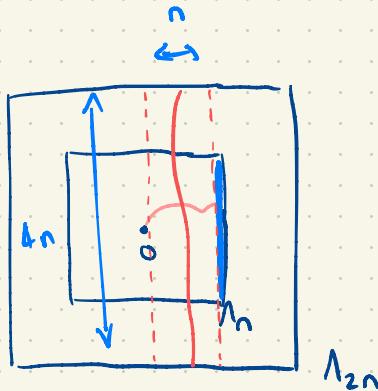
$$P_{\frac{1}{2}}[0 \leftrightarrow \partial \Lambda_n \text{ in } \mathbb{Z}_{2n+1} \times \mathbb{Z}] \geq \frac{1}{2n+1}$$

- (d) the event $0 \leftrightarrow \partial \Lambda_n$ is dependent only on the configuration strictly inside Λ_n . since the config inside is independent of outside, we have that it is distributed the same as that inside Λ_n on \mathbb{Z}^2 .

hence

$$\begin{aligned} P_{\frac{1}{2}}[0 \leftrightarrow \partial \Lambda_n \text{ in } \mathbb{Z}_{2n+1} \times \mathbb{Z}] \\ = P_{\frac{1}{2}}[0 \leftrightarrow \partial \Lambda_n \text{ in } \mathbb{Z}^2] \\ \geq \frac{1}{2n+1}. \end{aligned}$$

3. (a)

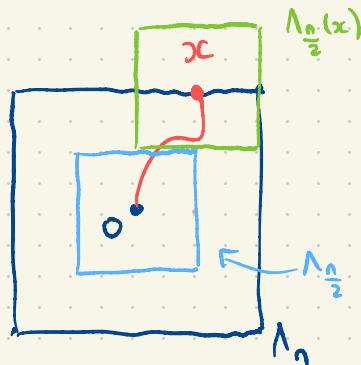


$$\text{it is clear that } P_{\frac{1}{2}}[0 \leftrightarrow \text{right side of } \partial A_n] \\ \geq \frac{1}{4} P_{\frac{1}{2}}[0 \leftrightarrow \partial A_n]$$

let A_n be the event that $[0, n] \times \{-2n\} \leftrightarrow [0, n] \times \{2n\}$
in the box $[0, n] \times [-2n, 2n]$
(seen in red above) then by RSW, $\exists c > 0$ st. $\forall n \geq 1$
 $P_{\frac{1}{2}}[A_n] \geq c.$

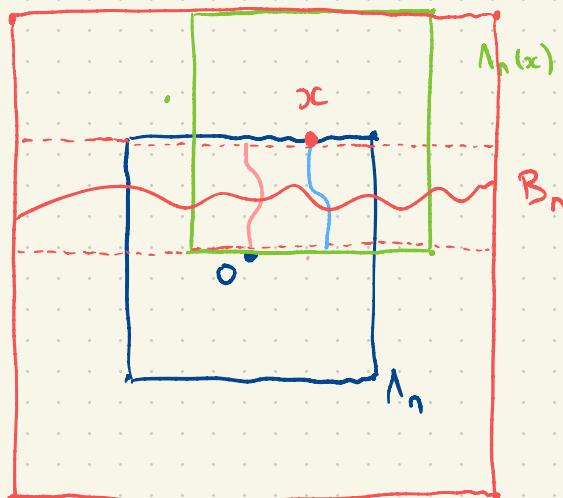
$$\text{now } P_{\frac{1}{2}}[0 \leftrightarrow A_{2n}] \geq P_{\frac{1}{2}}[A_n \cap \{0 \leftrightarrow \text{right side of } \partial A_n\}] \\ \stackrel{\text{FKG}}{\geq} P_{\frac{1}{2}}[A] P_{\frac{1}{2}}[0 \leftrightarrow \text{right side of } \partial A_n] \\ \geq \frac{c}{4} P_{\frac{1}{2}}[0 \leftrightarrow A_n].$$

b



$$\begin{aligned} P_{\frac{1}{2}}[O \leftrightarrow x] &\leq P_{\frac{1}{2}}[\{O \leftrightarrow \partial \Lambda_{\frac{n}{2}}\} \cap \{x \leftrightarrow \partial \Lambda_{\frac{n}{2}}(x)\}] \\ &= P_{\frac{1}{2}}[\{O \leftrightarrow \partial \Lambda_{\frac{n}{2}}\}]^2 \stackrel{\text{independence}}{\leq} C \cdot P_{\frac{1}{2}}[O \leftrightarrow \partial \Lambda_n]^2 \end{aligned}$$

part (a)

wlog assume $x \in \text{top of } \partial \Lambda_n$.

$$\{O \leftrightarrow \text{top}(\partial \Lambda_n)\} \cap \{x \leftrightarrow \text{bottom}(\partial \Lambda_n(x))\} \cap B_n \subset \{O \leftrightarrow x\}$$

where $B_n = \{2n\} \times [0, n] \leftrightarrow \{2n\} \times \{0, n\}$ in $[-2n, 2n] \times [0, n]$
 (shown in red).

then $P_{\frac{1}{2}}[O \leftrightarrow x]$

$$\geq P_{\frac{1}{2}}[\{O \leftrightarrow \text{top}(\partial A_n)\} \cap \{x \leftrightarrow \text{bottom}(\partial A_n(x))\} \cap B_n]$$

$$\stackrel{\text{FKG}}{\geq} P_{\frac{1}{2}}[O \leftrightarrow \text{top}(\partial A_n)] P_{\frac{1}{2}}[x \leftrightarrow \text{bottom}(\partial A_n(x))] P_{\frac{1}{2}}[B_n]$$

$$\stackrel{\text{RSW}}{\geq} \frac{1}{16} P_{\frac{1}{2}}[O \leftrightarrow \partial A_n]^2 \cdot c.$$

$$N > 2n$$

(c) let $1 \leq n \leq N$. show

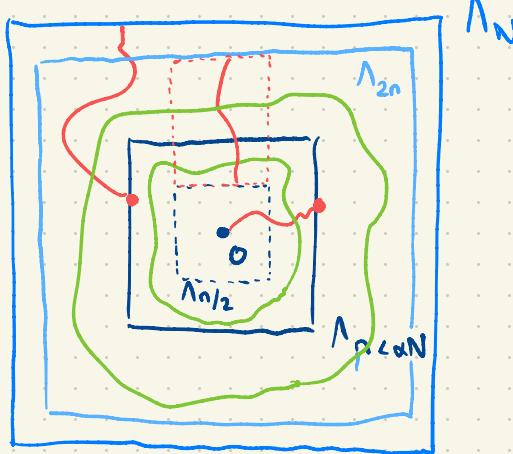
$$P_{\frac{1}{2}}[O \leftrightarrow N] \leq P_{\frac{1}{2}}[O \leftrightarrow \partial A_n] P_{\frac{1}{2}}[\partial A_n \leftrightarrow \partial A_N]$$
$$\leq c \cdot P_{\frac{1}{2}}[O \leftrightarrow N]$$



$$\{O \leftrightarrow \partial A_N\} \subset \{O \leftrightarrow \partial A_n\} \cap \{\partial A_n \leftrightarrow \partial A_N\}$$

$$\text{well, } P_{\frac{1}{2}}[O \leftrightarrow N] \leq P_{\frac{1}{2}}[O \leftrightarrow \partial A_n, \partial A_n \leftrightarrow \partial A_N]$$

$$\stackrel{\text{independence}}{=} P_{\frac{1}{2}}[O \leftrightarrow \partial A_n] P_{\frac{1}{2}}[\partial A_n \leftrightarrow \partial A_N]$$



$$P_{\frac{1}{2}}[0 \leftrightarrow \partial \Lambda_N] \geq$$

$$P_{\frac{1}{2}}[\{\emptyset \leftrightarrow \partial \Lambda_n\} \cap \{\partial \Lambda_n \leftrightarrow \partial \Lambda_N\} \cap \{\text{circuit in } \Lambda_n \setminus \Lambda_{\frac{n}{2}}\} \\ \cap \{\text{circuit in } \Lambda_{2n} \setminus \Lambda_n\} \\ \cap \{\text{vertical crossing in } [-\frac{n}{2}, \frac{n}{2}] \times [\frac{n}{2}, \min\{2n, N\}]\}]$$

FKG

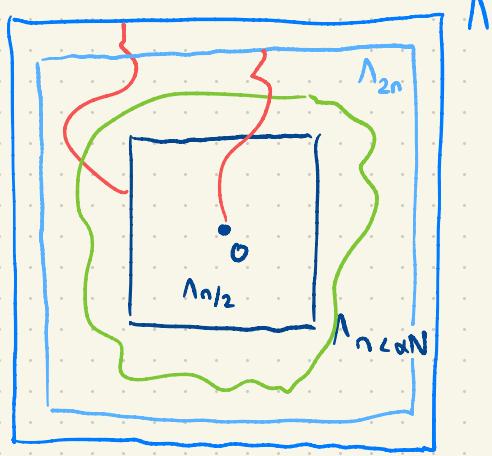
$$\geq P_{\frac{1}{2}}[0 \leftrightarrow \partial \Lambda_n] P_{\frac{1}{2}}[\partial \Lambda_n \leftrightarrow \partial \Lambda_N]$$

$$P_{\frac{1}{2}}[\text{circuit in } \Lambda_n \setminus \Lambda_{\frac{n}{2}}]$$

$$P_{\frac{1}{2}}[\text{circuit in } \Lambda_{2n} \setminus \Lambda_n]$$

$$P_{\frac{1}{2}}[\text{vertical crossing in } [-\frac{n}{2}, \frac{n}{2}] \times [\frac{n}{2}, 2n]]$$

RSW $\geq P_{\frac{1}{2}}[0 \leftrightarrow \partial \Lambda_n] P_{\frac{1}{2}}[\partial \Lambda_n \leftrightarrow \partial \Lambda_N] \cdot c_1^2 \cdot c_2$



shorter proof
by tuomas

$$P_{\frac{1}{2}}[O \leftrightarrow \partial \Lambda_N] \geq P_{\frac{1}{2}}[O \leftrightarrow \partial \Lambda_{2n}, \Lambda_n \leftrightarrow \partial \Lambda_N, \\ \text{arcuit in } \Lambda_{2n} \setminus \Lambda_n]$$

$$\stackrel{\text{FKG}}{\geq} P_{\frac{1}{2}}[O \leftrightarrow \partial \Lambda_{2n}] P_{\frac{1}{2}}[\Lambda_n \leftrightarrow \partial \Lambda_N] P_{\frac{1}{2}}[\text{arcuit in } \Lambda_{2n} \setminus \Lambda_n]$$

RSW and part (a)

$$\geq c_1 \cdot c_2 \cdot P_{\frac{1}{2}}[O \leftrightarrow \partial \Lambda_n] P_{\frac{1}{2}}[\Lambda_n \leftrightarrow \partial \Lambda_N]$$