

# Network communication

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### - Communication

- Characteristics of physical and protocol layers
- Models, protocols and coordination
- Communication in microservice architectures
  - Failures
- Coupling and modularity
  - Types of coupling in different architectural approaches
  - Achieving modularity at different quanta



### **Communication networks**

- Even "virtualized" networks operate in physical reality
  - Sometimes can assume locality e.g. loopback speeds (pods)
  - Distribution and decentralization may hide physical aspects
    - Translation: System might be placed to straddle a buggy or oversubscribed router/switch
  - Physical latencies: speed of light & electric signals, processing delays in switches and routers; retransmits



- ~ 40 000 km \* ½ / 300 000 km/s / 2/3 = 100 ms
- + amplification delay+ routing delay



### **Communication networks**

- Assumption of "infinite network capacity" in cloud may fail
  - Loss of 50% of network capacity in a datacenter (backhoe)
  - Limits at virtual machines and physical servers (AWS ENA 25 Gbps)
  - → even if you cannot saturate an IaaS PoP, you can closer to your service
- Further delays and failures from protocols and OS
  - Number of concurrent TCP connections (OS)
  - Bugs in protocol implementations (usually non-OS)



# **Network protocols**

### - TCP almost universal, but

- In some situations UDP may be more suitable (<u>within</u> a service)
- SCTP tends to keep popping up (alpha in K8S v1.12)

- IPv6

- Getting more common, but still mostly user-side requirement
- Deep magic
  - TCP slow start algorithm originally for congestion control
  - Anycast, multicast and broadcast (if you control subnets and/or routers)
  - VPNs and tunnels sometimes for integration (island hopping)



# **Application protocols**

- Almost all service interactions occur at application level protocols
  - HTTP and HTTPS primary (QUIC in the future?)
    - *HTTP(S)* used to transport other application level protocols
    - *SOAP, REST, ...*
  - gRPC, Thrift, AQMP, etc.
- Operate on top of TCP
  - Sometime work around TCP issues (such as slow start, with Keep-Alive connections)
  - TCP is connection-oriented: connect  $\rightarrow$  transmit  $\rightarrow$  close
  - Usually client-server, e.g. specific listener <u>address</u> and <u>port</u>



# **Communication models**

#### - Synchronous response

- Request-response pattern
- Reply expected immediately (after processing)
- Asynchronous response
  - Processing started by request
  - Immediate response provides a handle or identifier
  - Response methods
    - Polling by client (known endpoint or part of response)
    - Callback from server (agreed-upon endpoint or part of request)
    - Response publish (message queue, pubsub, blackboard, ...)

### - Message-passing

- Request itself asynchronous



## Synchronous request





### **Asynchronous communication models**





# Asynchronous requests vs. asynchronous applications

- Asynchronous communication is between two parties
- Asynchronous applications are self-contained
  - Avoid blocking at thread level: some other method of waiting
  - Event loops, I/O selectors, continuations, etc. low-level mechanisms
  - Erlang, AKKA, asyncio programming language level constructs
- A synchronous application can make asynchronous requests
- An asynchronous application can make synchronous requests



# Failures



### **Failures in distributed systems**

### - Rule of thumb:

- Everything fails all the time (randomly, when least expected)
- See <u>Network is reliable</u> paper (hint: it is not)

#### - Microservice architectures fail more

- More components, more computers, more connections, more changes, more of everything
- Risks of correlated failures can be either higher or lower than for monolithic systems
- See first lecture slide how number of components affects reliability



### Synchronous request





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# But wait, it gets worse!





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Just to understand that they lurk everywhere.

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### **Asynchronous communication models**





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# **Addressing network failures**

- Failing terribly is better than hanging indefinitely
  - At least you can see them in logs / monitoring
- All low-level socket operations can fail
  - This includes close() ... in Java, even it can cause an exception
- All network operations should have timeouts
  - Abstractions may try to hide the network (Java RMI)
- Long-lived quiescent connections are subject to random network dropouts
  - Stateful firewalls
  - TCP and protocol-level keepalives (ping, echo, ...)



# Addressing protocol failures

- Specification:
  - GET /resource
  - 200 OK with application/json or 503 Service Unavailable
- Code:

```
- resp = conn.get("/resource")
if resp.code == 200:
    j = json.loads(resp.body)
...
elif resp.code == 503:
```

- Transparent proxy
  - May return 504 Gateway timeout
  - Might randomly respond with text/html advert page
- We'll cover architectural approaches to handling network and remote failures later in the course



### **Brewer's theorem (aka CAP)**

- The CAP theorem (later proven) states that for <u>distributed</u> <u>systems</u>, out of
  - Consistency
  - Availability
  - Partition-tolerance

it is possible to achieve only CP or AP all the time

- <u>See this</u> for later description of the theorem (with critique)



### **Brewer's theorem's consequences**

- Hard partitions are generally rare
  - Most of the time it is possible to achieve both consistency and availability
- However, partitions do still occur
  - Then you need to choose between availability and consistency
  - "Eventually consistent" mechanisms choose availability
- In large enough systems, something fails all the time
- Consideration in services which is critical?



### **Brewer's theorem's consequences**

- Just accept that



it is not possible to get ACID guarantees in a distributed system

&

#### all microservice architectures are distributed systems.

