

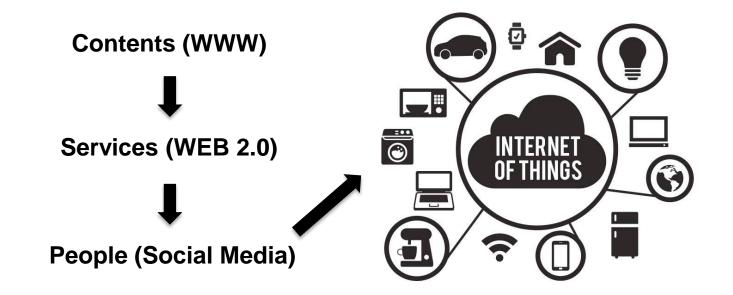
## Internet of Things (IoT) Security: Threats, Challenges & Recommendations

Samuel Marchal

## What is the Internet of Things ?

#### Latest evolution of the Internet:

• machine to machine / device to device communications



## Definition

#### Wikipedia:

• "Network of physical objects, devices, vehicles, buildings and other items embedded with electronics, software, sensors, and network connectivity that enables these objects to collect and exchange data."

#### Gartner:

• "Network of physical objects that contain embedded technology to communicate and sense or interact with their internal states or the external environment."

## IoT-based DDoS attack are the most powerful

#### State-of-the-Art DDoS attack: 20 Gbps

#### January 2016 (new World Hackers)

- IoT-based DDoS against *bbc.co.uk*
- 600 Gbps (2x previous record)

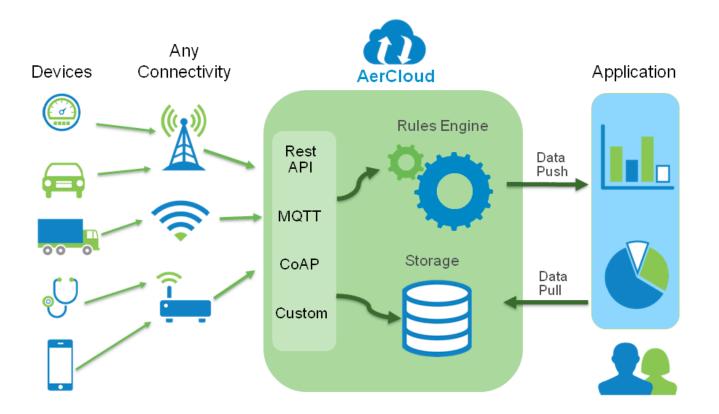
#### September 2016 (Mirai botnet)

- DDoS against OVH web host
- 1.5 Tbps
- 150,000 IoT devices (IP cameras / DVRs) used

#### **October 2016 (Mirai botnet)**

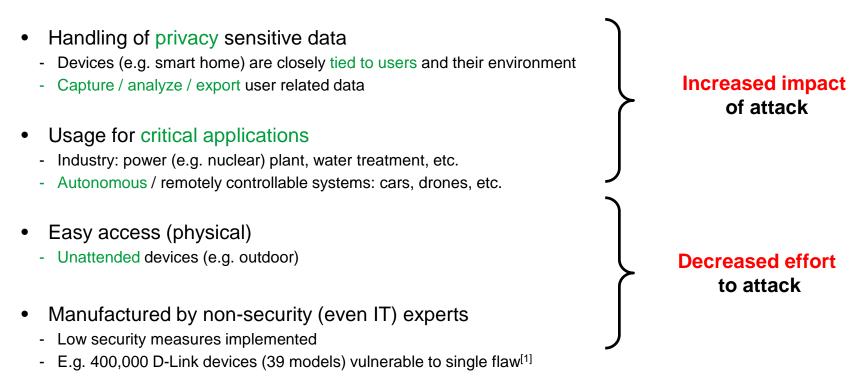
- DDoS against Dyn: dynamic DNS service
- 10s of millions of IoT devices (IP address)
- Disrupts service of Twitter, Spotify, Reddit, SoundCloud, PayPal, etc.

## IoT ecosystem



**Threats & Challenges** 

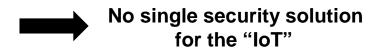
## **IoT specific threats**



[1] http://blog.senr.io/blog/400000-publicly-available-iot-devices-vulnerable-to-single-flaw

## **IoT heterogeneity**

- Entity heterogeneity (3 tiers)<sup>[1]</sup>:
  - High-end devices (laptop, smartphone, tablets)
  - Low-end devices (sensors, actuators)
  - Passive entities (barcode, QR-code, RFID)
- Communication heterogeneity:
  - Wired communications (ethernet)
  - Long range wireless communications: WiFi / 3G / 4G
  - Short range wireless communications: Bluetooth (LE) / Zigbee / 6LoWPAN



[1] Corvington and Carskadden "Threat Implications of the Internet of Things", 2013

## Node Constraints <sup>[1]</sup>

- Maximum code complexity (ROM/Flash)
- Size of state buffers (RAM)
- Amount of computation feasible in a period of time (processing capabilities)
- Available power
- User interface and accessibility in deployment (set keys, update software)

[1] RFC7228 "Terminology for Constrained-Node Networks" (https://tools.ietf.org/html/rfc7228)

## **Network Constraints**<sup>[1]</sup>

- Low achievable throughput
- High packet loss
- Asymmetric link characteristics
- Penalties for using large packets (e.g. high packet loss due to link layer fragmentation)
- Reachability over time (wake-up and sleeping time of devices)
- Lack of advance services (e.g. IP multicast)

[1] RFC7228 "Terminology for Constrained-Node Networks" (https://tools.ietf.org/html/rfc7228)

# Designing secure IoT systems

## Securing the IoT

- System (access control, authentication, etc.)
- Application
- Mobile
- Cloud
- Network (communications)

## **Recommended approaches**

- 1. Threat analysis (e.g. RFC 3552)
- 2. Follow security recommendation (e.g. NIST, IETF, etc.)
- 3. Learn from attacks
- 4. Follow design patterns

## 1. Threat Analysis

### • Define assumptions

- E.g. "the attacker has complete control of the communication channel"
- Explore scenarios
  - Active vs. passive attacker
  - On-path vs. off-path



- Fulfill requirements:
  - Authentication
  - Authorization
  - Traffic Security (confidentiality, data-origin, integrity, availability)
  - Non-repudiation (optional)

## **Threat Analysis: Limitations**

#### Gives theoretic security requirements to meet

#### But: leaves room for interpretation in implementation

- Which layer to apply security protection to ?
- Which existing security frameworks to use ?

#### **Complex to perform**

- Difficult to be comprehensive
  - E.g. consideration of vulnerable devices used as attack vector

## 2. Security recommendations (e.g. NIST, IETF)

- Key management requirements <sup>[1]</sup>
- Key length recommendations<sup>[2]</sup>
- Randomness requirements [3]
- Avoid possibility of pervasive monitoring
- Protocol or domain specific recommendation (crypto algorithm, WLAN security, use of TLS / DTLS, etc.)

[1] RFC 4107 "Guidelines for Cryptographic Key Management" (<u>https://tools.ietf.org/html/rfc4107</u>)
[2] RFC 4492 "Elliptic Curve Cryptography Cipher Suites for TLS" (<u>https://tools.ietf.org/html/rfc4492</u>)
[3] RFC 4086 "Randomness Requirements for Security" (<u>https://tools.ietf.org/html/rfc4086</u>)

## 3. Learn form Attacks

#### Selected attacks to illustrate common problems:

- Inadequate software update mechanism
- Missing Key Management
- Insecure configuration files and default passwords
- Missing communication security
- Physical attacks

## Inadequate software update mechanism

#### **Example: Huawei Home gateway**

- Embedded web server (released 2002) with buffer overflow vulnerability
- Fix released in 2005 by web server company
- Vulnerability still exploited <sup>[1]</sup> (2015)

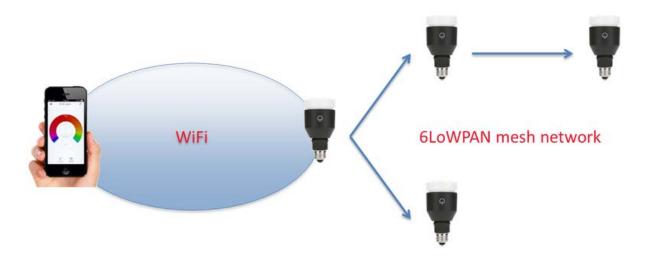


[1] http://www.computerworld.com/article/2860843/vulnerability-in-embedded-web-server-exposes-millions-of-routers-to-hacking.html

## **Missing Key Management Problem**

### Example: LIFX <sup>[1]</sup> - Internet connected light bulb

• AES key shared among all devices to simplify key management



[1] http://www.contextis.com/resources/blog/hacking-internet-connected-light-bulbs/

# Insecure Configuration Files and Default Passwords

Example: Foscam, Linksys, Panasonic surveillance / baby monitoring cameras

- Default passwords or insecure default settings
- Similar problems on LED bulbs, etc.

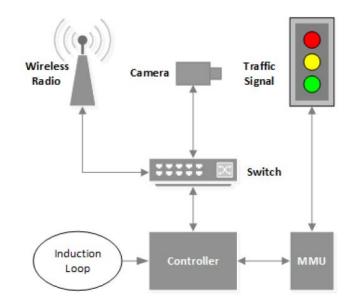


[1] <u>http://www.networkworld.com/article/2844283/microsoft-subnet/peeping-into-73-000-unsecured-security-cameras-thanks-to-default-passwords.html</u>

## **Missing Communication Security**

## Example: Traffic infrastructure <sup>[1]</sup>

- Unencrypted wireless communications
- Default Username and Passwords (published online by manufacturer)
- Controller settings can be configured remotely
- FTP connection to write configuration files
- Physical attacks

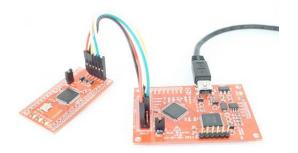


[1] Ghena, et al. Green "Lights Forever: Analyzing the Security of Traffic Infrastructure"

## **Physical Attacks**

## Example: extract keys, configuration data, firmware images

- Use of debug / test interfaces
- Sniffing on inter-bus communication interfaces be comprehensive
  - Serial Peripheral Interface (SPI)
  - Inter-Integrated Circuit (I<sup>2</sup>C)
- Key extraction within a trusted execution environment
  - Power analysis
  - Fault injection (glitching) attacks



## Intermediate Summary (methods 1/2/3)

- 90% of the threats are common among all protocols
- Most (exploited) security vulnerabilities are basic
- Many exploits of IoT systems (particularly in the consumer space) were hoaxs. But this is changing: Mirai botnet, Reaper botnet

## 4. Communications Design Patterns

- Device-to-Device
- Device via Smart Phone to Cloud
- Device via Gateway to Cloud
- P2P Communication in Local Network
- Device-to-Cloud

## **Device-to-Device Communication**

#### **Characteristics:**

- Direct communications between devices
- Link layer communication protocol (often no IP)

#### Security:

- Direct association between devices: pairing
- Channel security provided at the link layer

#### Standardization:

- RFID, 6LowPAN, ZigBee
- Bluetooth Low Energy (LE)<sup>[1]</sup>





[1] https://www.bluetooth.com/what-is-bluetooth-technology/bluetooth-technology-basics/low-energy

## **Device via Smart Phone to Cloud**

#### **Characteristics:**

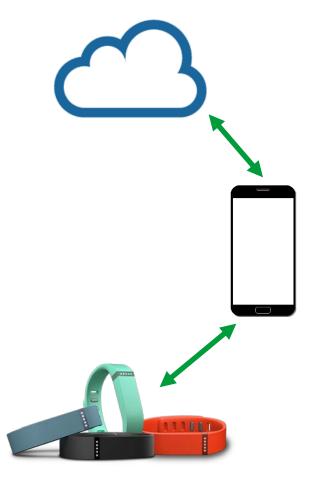
- Extension of device-to-device communication
- Device interacts with smart phone and cloud service

#### Security:

- D2D security
- Smart phone app / Web security
- Better provide end-to-end security (usually not the case)

#### Standardization:

• Bluetooth LE, NFC



## **Device via Gateway to Cloud**

#### **Characteristics:**

- Devices communicate with cloud services via a network gateway
- Apps/websites allow user-friendly, remote access/monitoring

#### Security:

- D2D security / No end-to-end security
- Authentication / pairing / key management
- Example: EAP, PANA, AAA, etc.

#### Standardization:

• IEEE 802.15.4, WiFi, Bluetooth LE

monitoring	$\mathcal{C}$	3
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## **P2P Communication in Local Network**

#### **Characteristics:**

- Variant of "device via gateway to cloud" with local-only operation.
- · Discovery of nodes to communicate with

#### Security:

- Authentication of nodes
- Trust management system
- Low need for protecting communications?

#### Standardization:

- Universal Plug and Play (UPnP) + UPnP-UP
- DNS Service Discovery
- Bonjour (Apple)



## **Device-to-Cloud**

#### **Characteristics:**

- Direct communication with cloud service
- Pre-configured for specific cloud service only
- Always-on reachability required
- Radio technology and IP-connectivity required

#### Security:

- Network access + cloud authentication
- End-to-end security

### Standardization:

• WiFi

## **Good Practices (recommendation)**

- Always encrypt avoid pervasive monitoring
- Follow encryption key length recommendation (112-bit symmetric key equivalent)
- Support automatic key management
- Automatic software update mechanism
- Communication channel security (e.g. DTLS/TLS)
- Authentication and authorization solution
- Reduce physical attack surface:
  - Crypto implementations that consider side channel attacks
  - Disabled debug facilities before launching product
  - Hardware-based crypto support
  - Memory protection unit (MPU) integration

# Research in IoT security: 1. IoT Sentinel

## **Security of IoT**

Many deployed IoT devices have security vulnerabilities

#### Patching of devices is challenging

- Patches often not available
- Missing facilities for automatic patching (software updates)

#### Increase in IoT malware and botnets that target consumer smart home IoT devices

• Mirai, Persirai, Reaper, Hajime, etc.

## **Securing IoT**

#### Securing IoT is challenging:

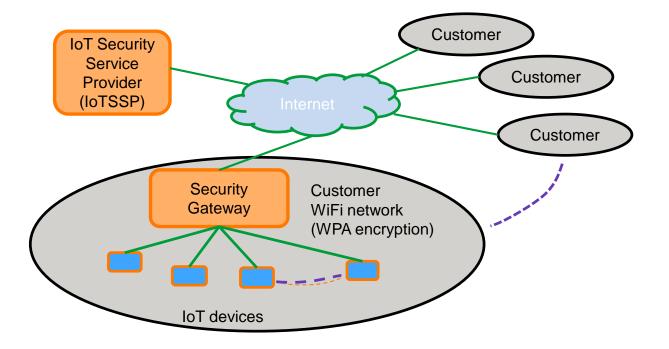
- Resource limitations
- Heterogeneity of IoT devices
- Scarcity of communications / etc.
- Daily release of new devices → dynamic threat landscape

Need for a *brownfield* security solution accommodating the fact that some devices are and will be *vulnerable* 

## **IoT Sentinel**

Need for an autonomous and adaptive solution for detecting compromised IoT devices DIOT

## System model



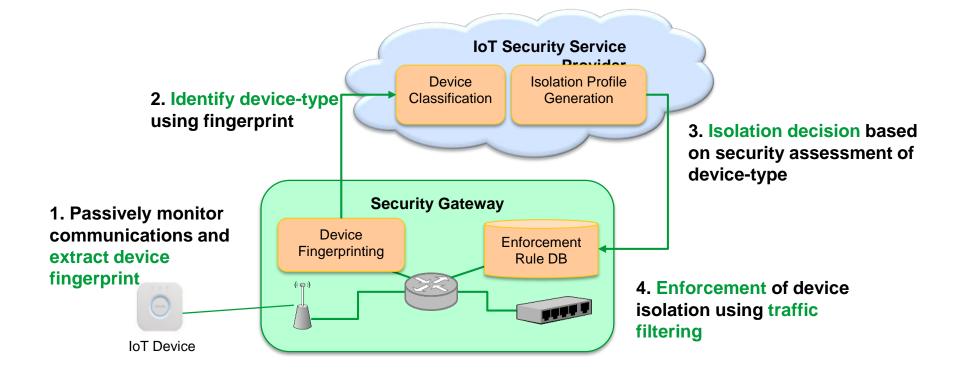
## IoT Sentinel: Automated Device-Type Identification for Security Enforcement in IoT



## **IoT Sentinel Principles**

- 1. Identify the type (make / model / SW version) of a new device when introduced to the network
- 2. Assign enforcement rules for device based on type (isolate known vulnerable types)
- 3. Constrain communications of vulnerable devices
  - Avoid vulnerable devices to be compromised
  - Contain infection

#### **IoT Sentinel system model**





# DioT: A Self-learning System for Detecting Compromised IoT Devices

#### **DÏoT: Addressing IoT challenges**

#### **Resource limitations**

Security monitoring delegated to network gateway

#### Heterogeneity of IoT devices

• Device-type specific detection model

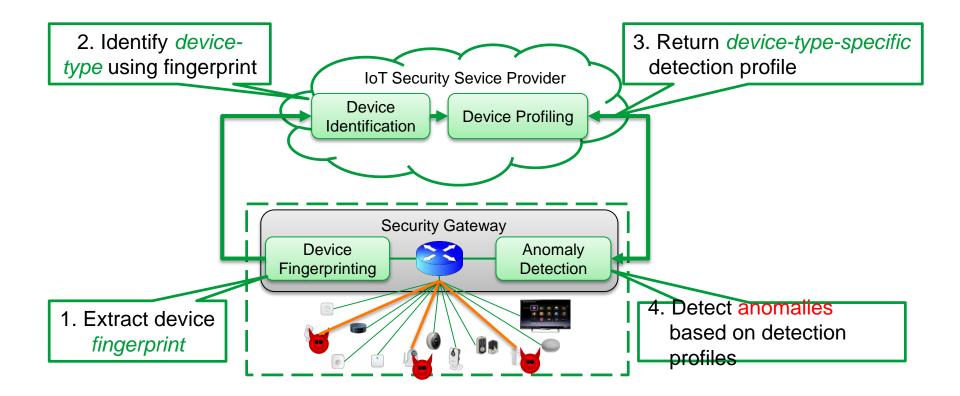
#### Scarcity of communications

- Device-type identification using periodic background traffic (always present and predictable)
- Fine grained language model of communications (GRU RNN) trained with little data

#### **Dynamic threat landscape**

- Adaptive: DIoT learns device-types and detection model as new IoT devices are deployed
- Fully autonomous / self-defined

#### **D**<sup><sup><sup>i</sup></sup></sup> D<sup><sup>i</sup></sup> OT system model





## Internet of Things (IoT) Security: Threats, Challenges & Recommendations

Samuel Marchal

## Attestation in the Internet of Things (IoT)

**Mobile Systems Security course** 

Lachlan Gunn

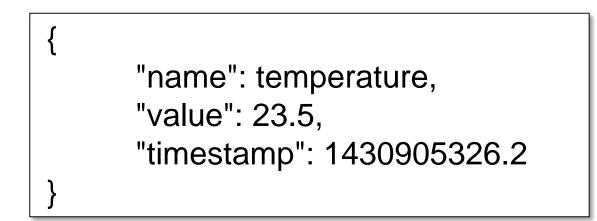
Aalto University lachlan.gunn@aalto.fi (Contributors: N. Asokan, Lucas Davi, Andrew Paverd)

#### You will learn about:

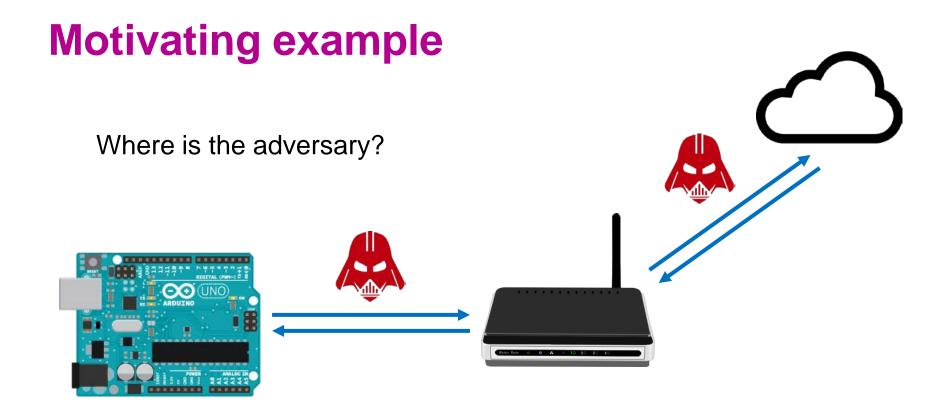
- Why do we need (remote) attestation in IoT?
- What technologies can we use for this?

### **Motivating example**

You receive the following message:



What does it mean?



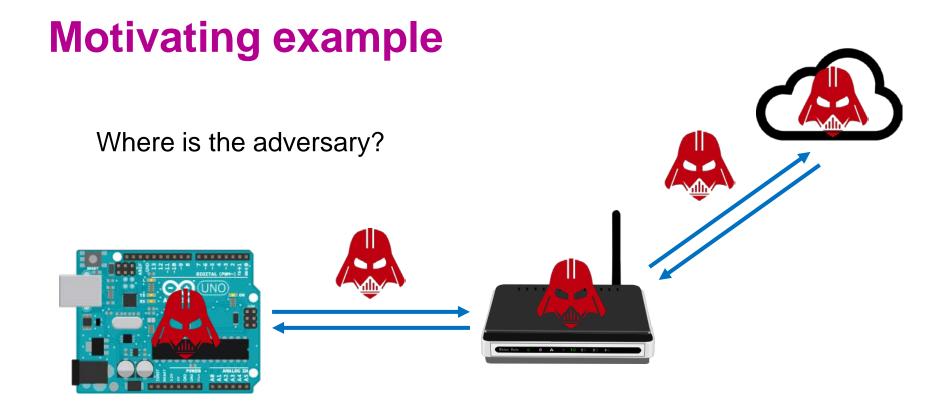
Network adversary: read, modify, falsify communication

### **Motivating example**

You receive the following message over an authenticated, integrity-protected communication channel:

{
 "name": temperature,
 "value": 23.5,
 "timestamp": 1430905326.2
}

What does it mean?



**Network adversary:** read, modify, falsify communication **Malware:** extract secrets, change state, modify behaviour

### **IoT malware: Stuxnet**

BUSINESS DULTURE DESIGN DEAR SCIENCE SECURITY TRANSPORT RIM ZETTER SECURITY 11.03.14 6:30 AM AN UNPRECEDENTED LOOK AT STUXNET, THE WORLD'S FIRST DIGITAL WEAPON	# WIR	ED A	n Unprecedented	d Look at Stuxnet	, the World's Firs.		SUBSCRIBE 🖌
AN UNPRECEDENTED LOOK AT STUXNET, THE	BUSINESS	CULTURE	DESIGN	GEAR	SCIENCE	SECURITY	TRANSPORTATIO
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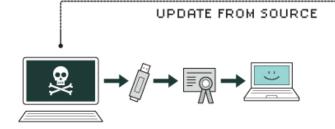
#### The Real Story of Stuxnet

How Kaspersky Lab tracked down the malware that stymied Iran's nuclear-fuel enrichment program

By David Kushner Posted 26 Feb 2013 | 14:00 GMT



#### IoT malware: Stuxnet



#### 1. infection

Stuxnet enters a system via a USB stick and proceeds to infect all machines running Microsoft Windows. By brandishing a digital certificate that seems to show that it comes from a reliable company, the worm is able to evade automated-detection systems.



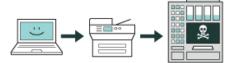
#### 2. search

Stuxnet then checks whether a given machine is part of the targeted industrial control system made by Siemens. Such systems are deployed in Iran to run high-speed centrifuges that help to enrich nuclear fuel.



#### 3. update

If the system isn't a target, Stuxnet does nothing; if it is, the worm attempts to access the Internet and download a more recent version of itself.



#### 4. compromise

The worm then compromises the target system's logic controllers, exploiting "zero day" vulnerabilitiessoftware weaknesses that haven't been identified by security experts.



#### 5. control

In the beginning, Stuxnet spies on the operations of the targeted system. Then it uses the information it has gathered to take control of the centrifuges, making them spin themselves to failure.



#### 6. deceive and destroy

Meanwhile, it provides false feedback to outside controllers, ensuring that they won't know what's going wrong until it's too late to do anything about it.

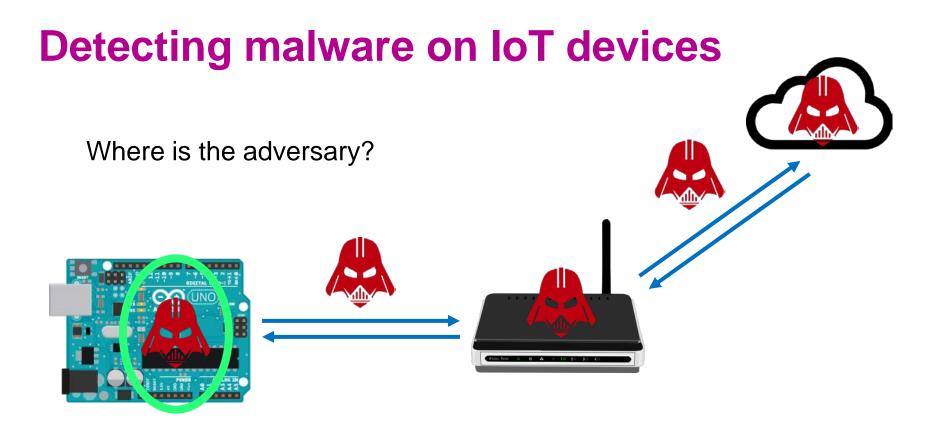
Illustration: L-Dopa http://spectrum.ieee.org/telecom/security/the-real-story-of-stuxnet

### **IoT** malware

Malware is not a new threat, but IoT...

- Broadens the attack surface
  - cost-constrained and/or resource-constrained devices
  - many more interconnected devices
- Amplifies the impact
  - access to detailed personal information
  - control of physical environment

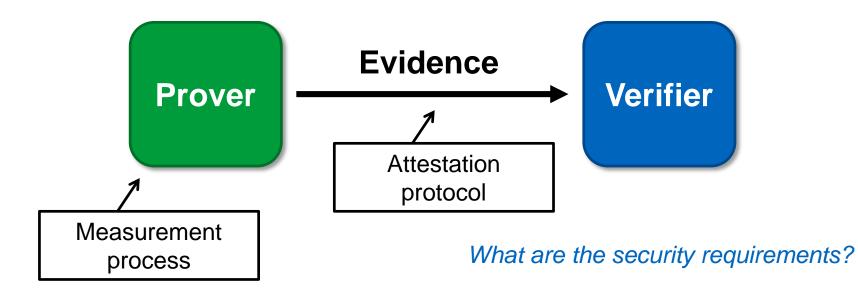




**Network adversary:** read, modify, falsify communication **Malware:** extract secrets, change state, modify behaviour

#### **Attestation**

• An interaction between two parties through which the *verifier* ascertains the current state and/or behaviour of the *prover*.



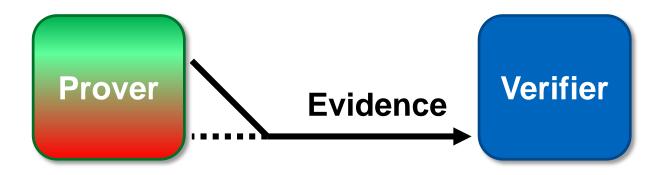
### **Attestation Requirements**

- 1. Authenticity
  - representation of the real state of the system



### **Attestation Requirements**

- 1. Authenticity
  - representation of the real state of the system
- 2. "Timeliness"
  - representation of the current state



### **TPM Attestation (Lecture 4)**



- TPM Platform Configuration Registers (PCRs)
  - store cryptographic hash
  - cannot be over-written by software, only extended

 $PCR_{1} = 0x00$ 

```
tpm_extend( PCR<sub>1</sub> hash<sub>1</sub>)
```

```
PCR_1 = hash( 0x00 || hash_1 )
```

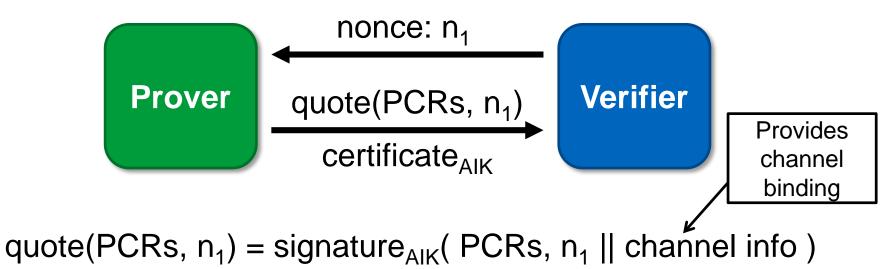
```
tpm_extend( PCR<sub>1</sub> hash<sub>2</sub>)
```

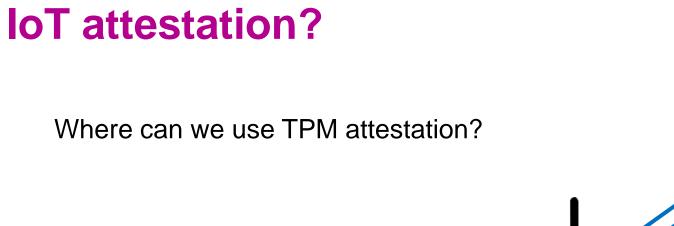
```
PCR_1 = hash(hash(0x00 || hash_1) || hash_2)
```

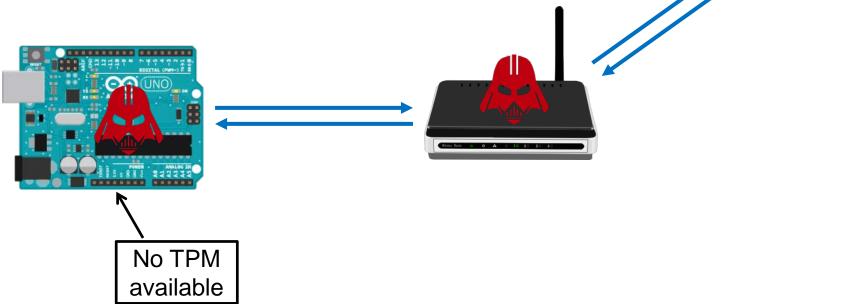
### **TPM Attestation (Lecture 4)**



- TPM Quote
  - PCR values signed by TPM-bound Attestation Identity Key (AIK)
  - includes nonce to ensure timeliness







### **TPM Attestation "Costs"**

- Additional hardware
  - takes up space
  - uses power
  - increases hardware cost (TPM + integration)
- Additional software
  - requires driver and software library

#### **TPM Attestation Limitations**

- Covers only the initial loading of software
- Deals with only one prover and one verifier
- "Decision of trustworthiness" does not scale
   measurements change with every software update

## **Attestation of Things (AoT)**

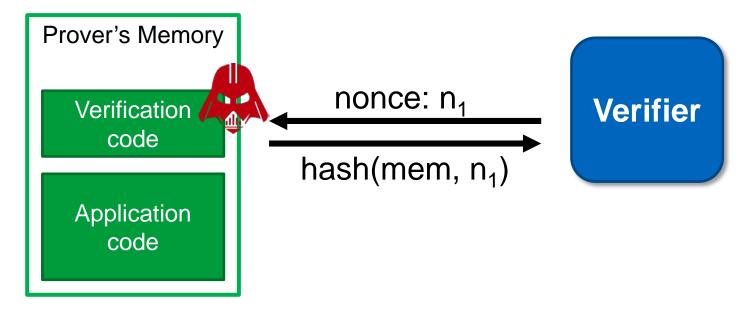
- Software-based attestation
- Hybrid attestation
- Scalability of attestation
- Run-time attestation

## **Attestation of Things (AoT)**

- Software-based attestation
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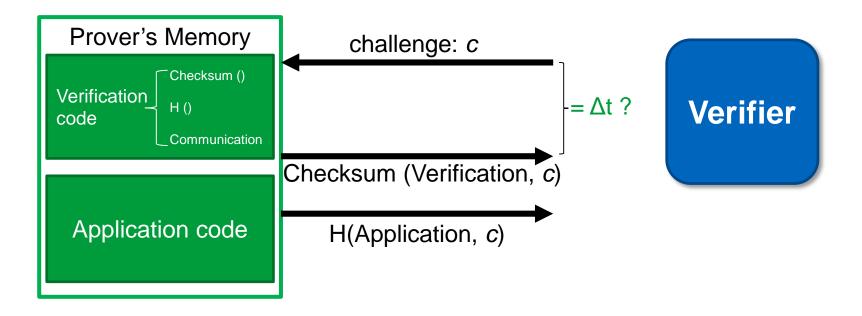
#### **Software-based Attestation**

- Assumes no hardware features to support attestation
  - No secrets on prover (e.g. no AIK)
  - Cannot guarantee specific code being run



#### **Software-based Attestation**

- Pioneer system
  - compute time-optimal checksum of verification code



A. Seshadri, M. Luk, E. Shi, A. Perrig, L. van Doorn, and P. Khosla. **Pioneer: Verifying integrity and guaranteeing execution of code on legacy platforms**. SOSP '05 Authenticity?

#### **Software-based Attestation: summary**

- Limitations of timing side channels
  - verifier must know exact hardware configuration
  - difficult to prove time-optimality
  - limited to "one-hop" networks
  - requires authenticated channel (e.g. physical connection)

## **Attestation of Things (AoT)**

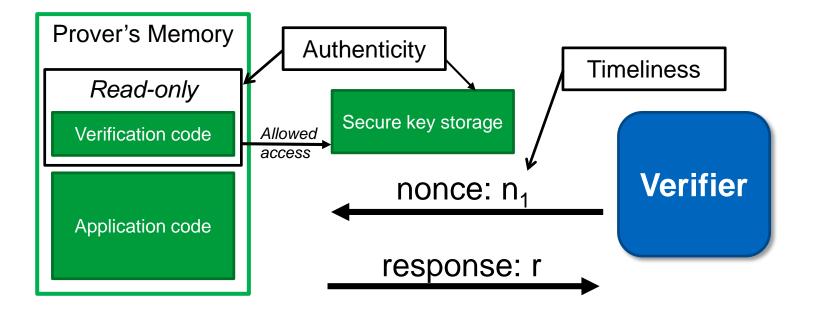
- Software-based attestation
- Hybrid attestation
- Scalability of attestation
- Run-time attestation

### **Hybrid Attestation**

- Minimal trust anchors
  - small changes to hardware
  - "hardware/software co-design"

### **Hybrid Attestation: SMART**

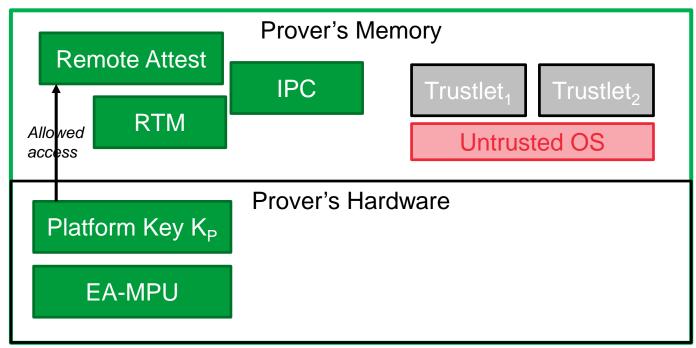
Read-only Verification code, secure key storage and atomicity of execution of Verification code



K El Defrawy, A. Francillon, D. Perito, and G. Tsudik. **SMART:** Secure and Minimal Architecture for (Establishing a Dynamic) Root of Trust. NDSS '12

### Hybrid Attestation: TrustLite & TyTAN

- Execution-Aware Memory Protection Unit (EA-MPU)
  - access control based on memory request target and origin



- P. Koeberl, et al. TrustLite: A Security Architecture for Tiny Embedded Devices. EuroSys '14
- F. Brasser, et al, TyTAN: Tiny Trust Anchor for Tiny Devices, DAC '15

### **Hybrid Attestation: summary**

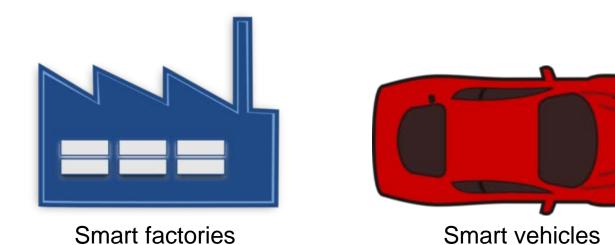
- Advantages of hybrid approaches
  - can be used across a network / over an untrusted channel
  - Verifier need not know prover's exact hardware configuration
- Drawbacks
  - Needs additional hardware support
  - But minimal MCU trust anchors soon available commercially
    - TrustZone-M (ARM Cortex-M 23/33)

## **Attestation of Things (AoT)**

- Software-based attestation
- Hybrid attestation
- Scalability of attestation
- Run-time attestation

#### **Scalability of Attestation**

- Attestation protocols usually assume a single prover
  - but IoT scenarios may involve groups of (many) provers



### **Scalability of Attestation**

- Device *swarms* 
  - dynamic topology: nodes move within swarm
  - dynamic membership: nodes join and leave the swarm



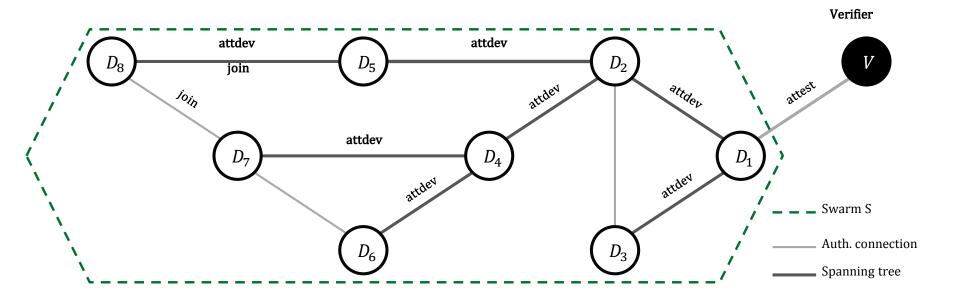
Smart factories



## **Scalability of Attestation**

- SEDA: Scalable Embedded Device Attestation
  - more efficient than attesting each node individually
  - can use any type of measurement process
  - assumes homogeneous provers

#### **Scalability of Attestation**



N. Asokan, F. Brasser, A. Ibrahim, A.-R. Sadeghi, M. Schunter, G. Tsudik, and C. Wachsmann. **SEDA:** Scalable Embedded Device Attestation. CCS '15

## **Scalability of Attestation: summary**

How to extend SEDA to

- support highly dynamic swarms?
- be resilient to physical compromise of some devices?

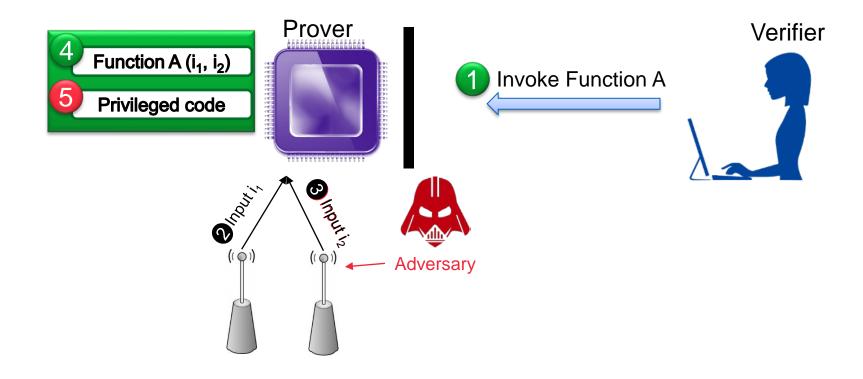
## **Attestation of Things (AoT)**

- Software-based attestation
- Hybrid attestation
- Scalability of attestation
- Run-time attestation

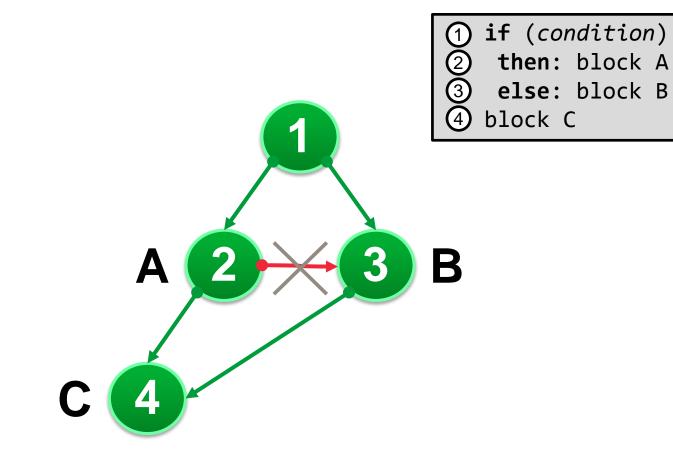
## Why run-time attestation?

- Traditional attestation measures binaries at load time
- Cannot capture run-time attacks
  - return-oriented programming
  - control data attacks

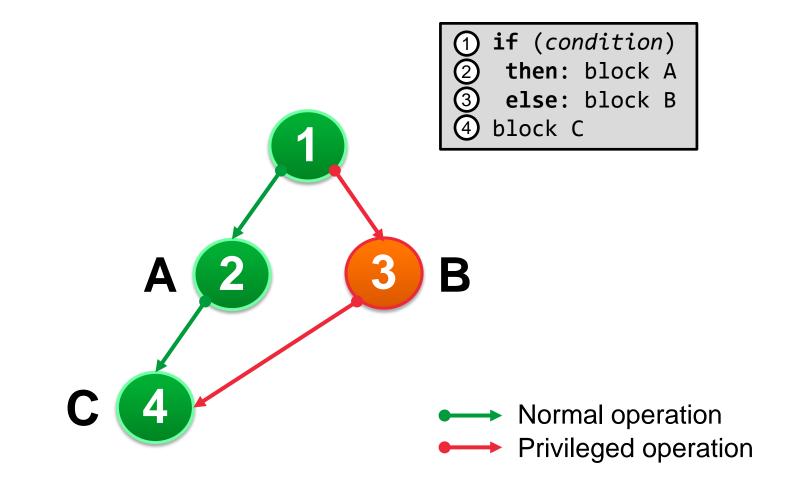
#### **Run-time attacks**

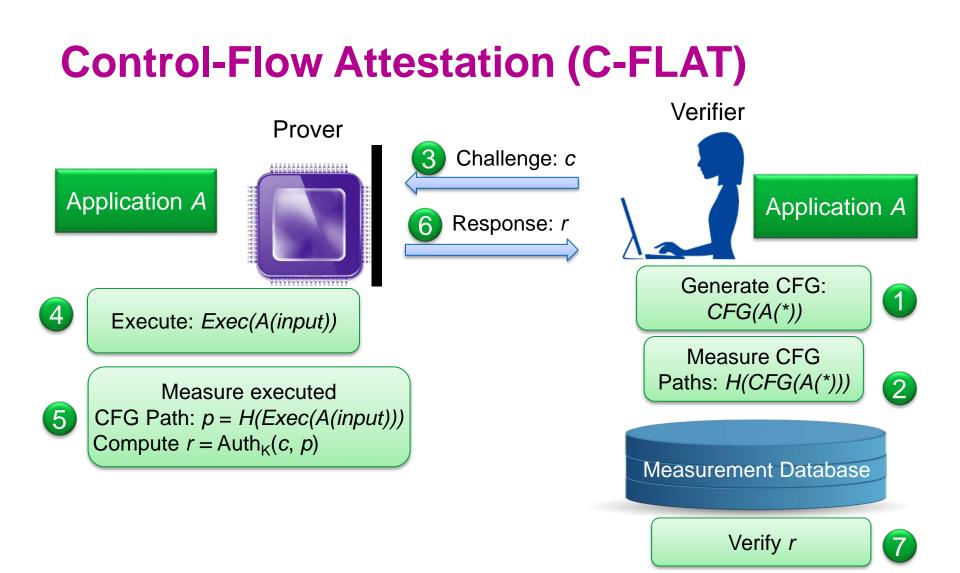


## **Control flow integrity (CFI)**



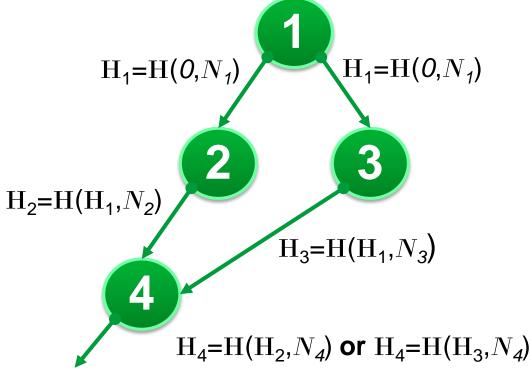
## **Run-time attacks without violating CFI**





## **C-FLAT: High-Level Idea**

• Cumulative Hash Value:  $H_j=H(H_i, N)$ , where  $H_i$  previous hash result and N is the current node

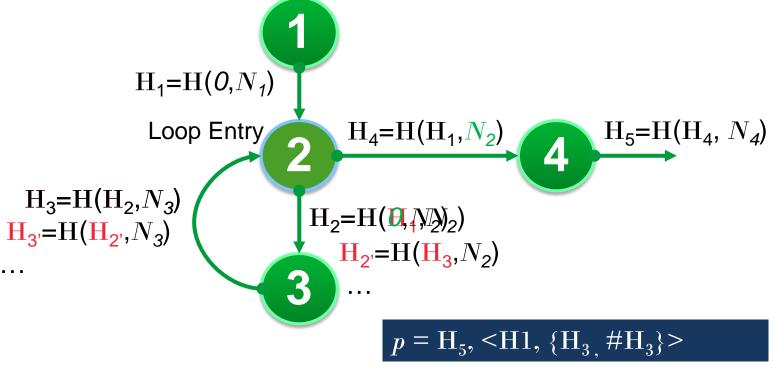




## **Handling Loops**

 ${\rm H}_{\rm X}$  different for each loop iteration

- Different loop paths/iterations  $\rightarrow$  many valid hash values
  - Our approach: treat loops as separate sub-graphs



## **Proof-of-Concept Implementation**

- Bare-metal prototype on Raspberry Pi 2
  - Single-purpose program instrumented using binary-rewriting
  - Runtime Monitor written in ARM assembler
  - Measurement Engine isolated in TrustZone-A Secure World

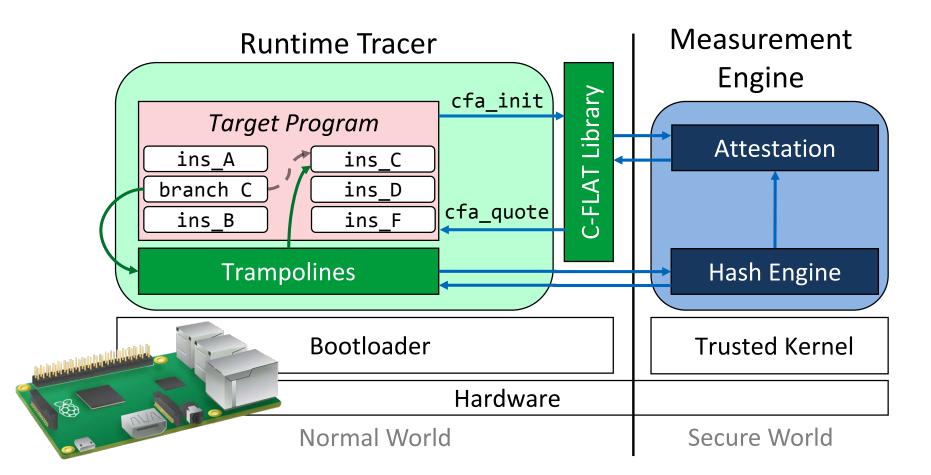


cfa\_quote: 7c 16 d6 51 20 a2 a0 c7 90 f5 ef 04 0c 2e ba bc
loop[000]: 78 22 5b 62 92 41 ca 02 7b ff 29 57 c6 6f 9b a2
path[000]: 2f a5 8c dc 1b 35 41 29 ab dd 35 5c f2 69 08 37 (1)
loop[001]: d6 90 9e a0 8c ae 90 84 9e 66 09 f8 a6 7b 52 04
path[000]: 92 fb d1 e8 90 cb 02 e5 6c f2 65 8c 86 72 0e d3 (2)

loop[006]: 05 e3 92 40 95 ef 7b 46 13 7d 6e 8b 05 be bf 41 path[000]: 67 c6 5e d4 18 13 02 bc 4a 5d 60 a0 16 85 f4 ed (9) path[001]: 78 19 af 09 0f d5 64 f4 39 b4 7a 0d 97 57 77 8c (2)

Source: https://github.com/control-flow-attestation/c-flat/blob/master/samples/syringe/syringe-auth.txt

## **Proof-of-Concept Implementation**

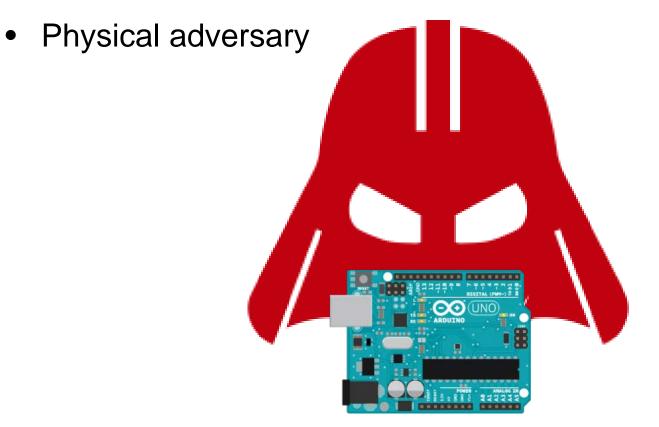


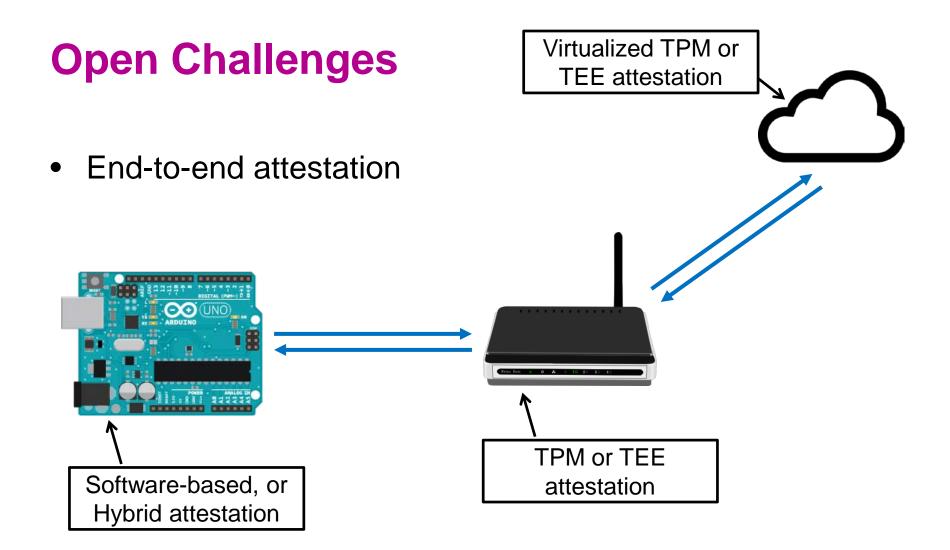
## **Run-time attestation: summary**

- How can we scale control flow attestation?
  - Better ways to aggregate measurements?
  - Faster/simpler purpose-built hash functions?
  - Purpose-built hardware support?
- What next?
  - Attestation of properties rather than measurements?
    - from attestation to checking compliance with a (dynamic) policy?

Ghada Dessouky, Shaza Zeitouni, Thomas Nyman, Andrew Paverd, Lucas Davi, Patrick Koeberl, N. Asokan, Ahmad-Reza Sadeghi, **LO-FAT: Low-Overhead Control Flow Attestation in Hardware**, DAC '17.

## **Open Challenges**





## **Open Challenges**

• Device heterogeneity



## **Did you learn about:**

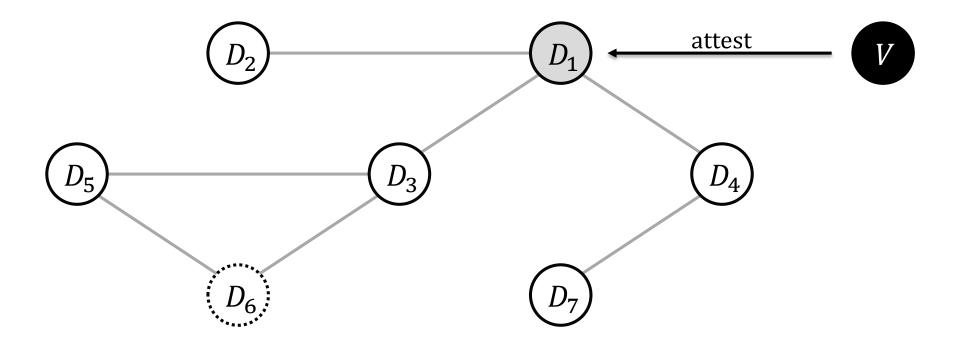
- Why we need attestation in IoT
- Some recent research towards achieving this

# Attestation in the Internet of Things (IoT)

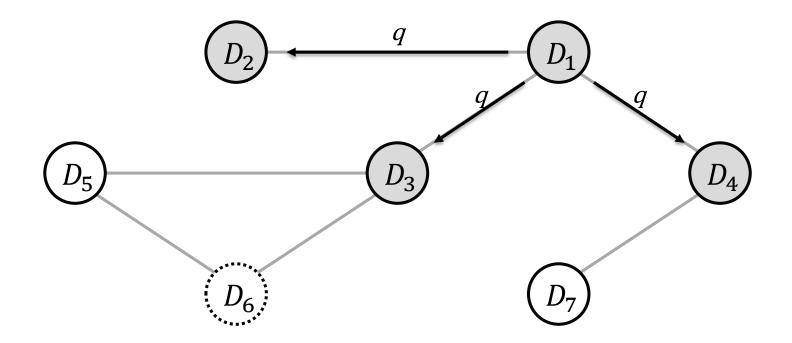
**Mobile Systems Security course** 

Lachlan Gunn

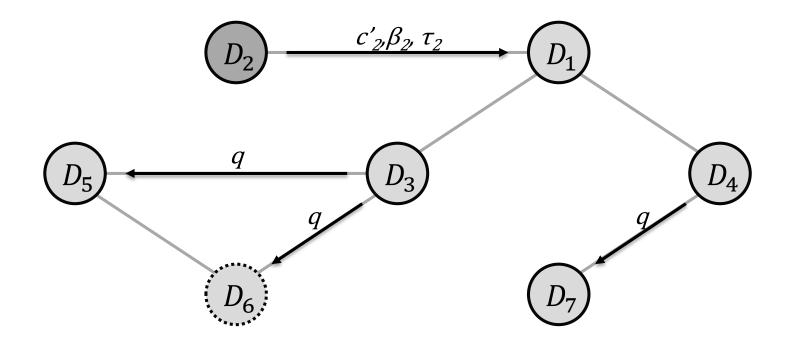
Aalto University lachlan.gunn@aalto.fi (Contributors: N. Asokan, Lucas Davi, Andrew Paverd)



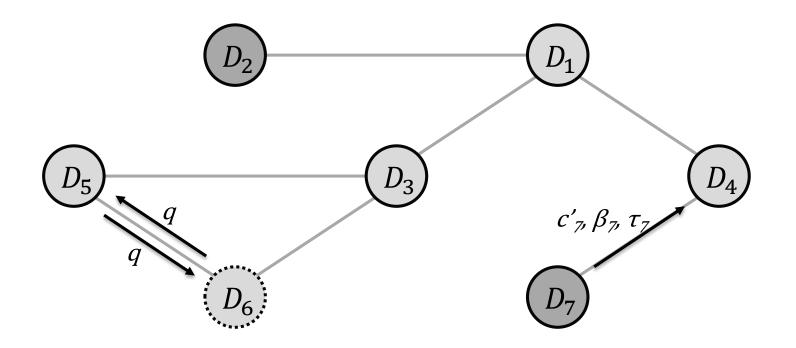
N. Asokan, F. Brasser, A. Ibrahim, A.-R. Sadeghi, M. Schunter, G. Tsudik, and C. Wachsmann. **SEDA:** Scalable Embedded Device Attestation. CCS '15



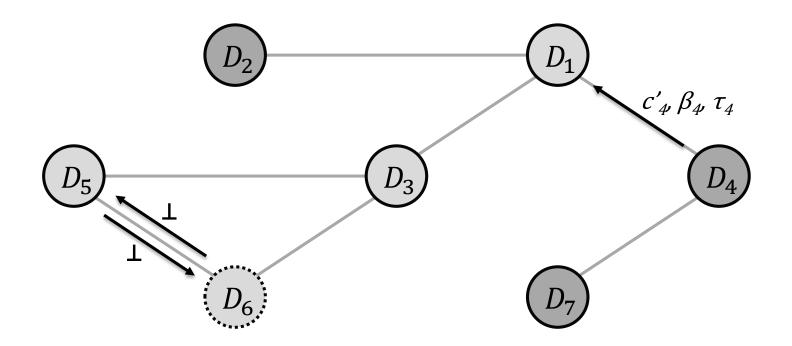
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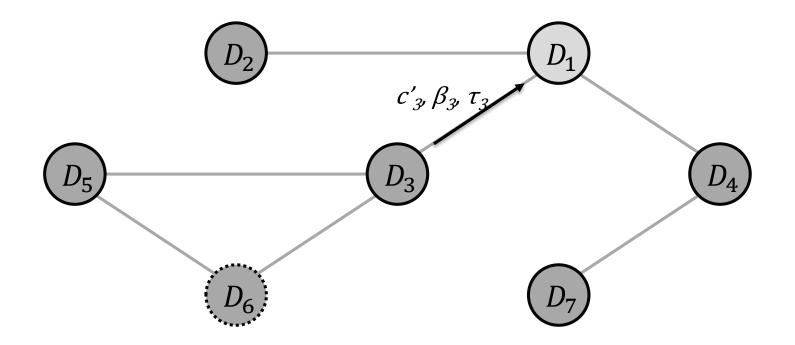


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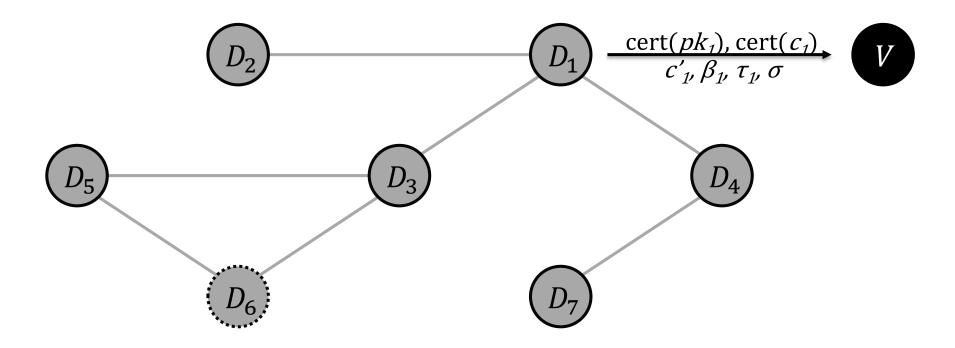
## **SEDA (in detail)** $D_1$ $c'_{5'} \beta_{5'} \tau_5$ $D_5$ $D_3$ $D_4$ $c'_{\theta} \beta_{\theta} \tau_{\theta}$ $D_6$

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97



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## **Abuse of Attestation**

- Attestation protocols usually consider a benign verifier
- But adversary could impersonate verifier to abuse attestation
- E.g., computing a MAC over all memory in a typical microcontroller could take ~ 750 ms
  - could be used for denial of service (DoS) attack

## **Abuse of Attestation**

- Verifier must be authenticated to prover
  - but asymmetric crypto is computationally expensive
- Prover must detect replays of previous requests
  - can use nonces
  - can use counters
  - can use timestamps

## **Abuse of Attestation**

- Nonces
  - require integrity-protected storage for previous nonces
- Counters
  - require minimal integrity-protected storage for counter
- Timestamps
  - require trusted synchronized clock at prover side

Use Execution-Aware Memory Access Control (EA-MAC) to protect counters and timers.

<u>F. Brasser, A.-R Sadeghi, K. B. Rasmussen, and G.</u> <u>Tsudik. Remote Attestation for Low-End Embedded</u> **Devices: the Prover's Perspective**. DAC '16